

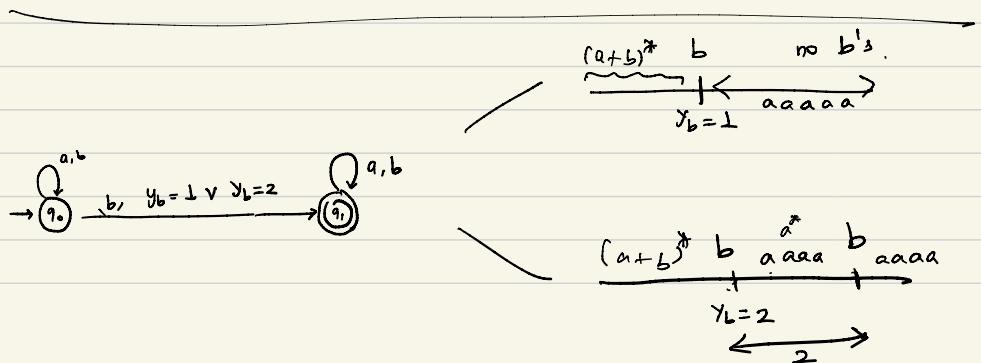
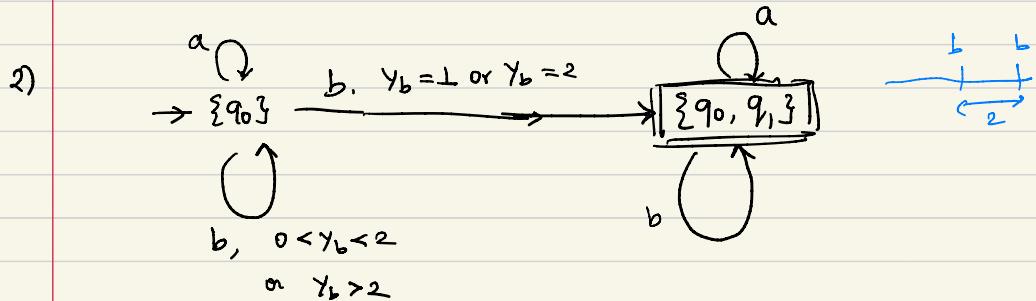
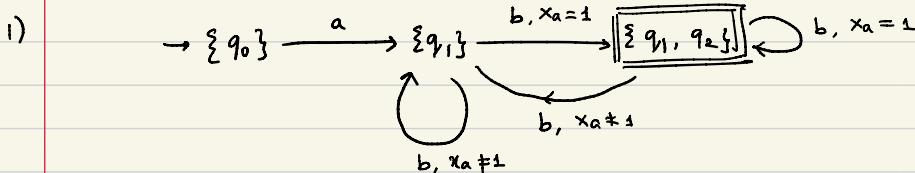
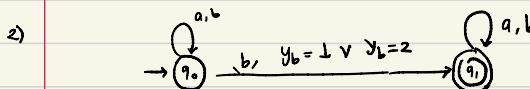
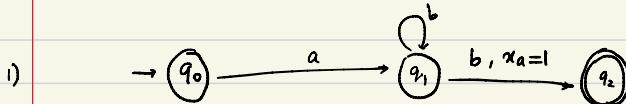
TIMED AUTOMATA

LECTURE 9

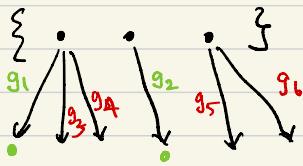
GOALS OF TODAY'S LECTURE:

- 1. Examples of determinizing E.C.A.
- 2. Expressive power among different classes of ECA and T.A.

Determinize the following ECA:

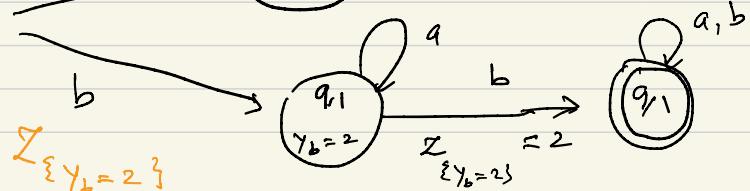
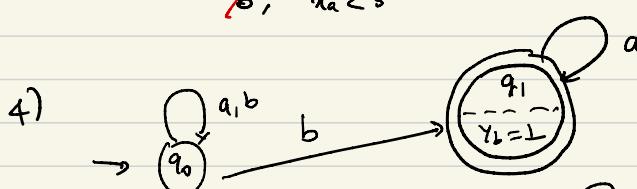
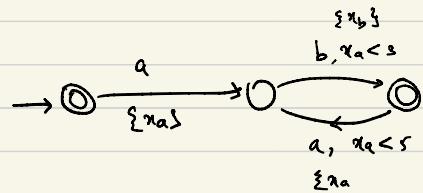
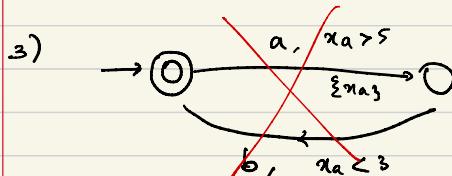
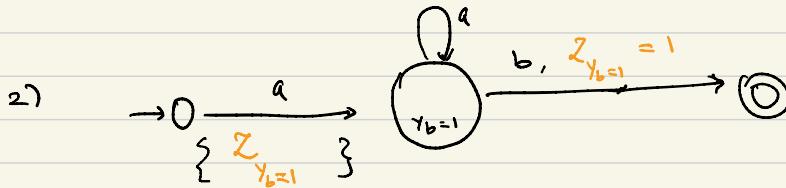
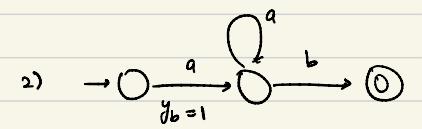
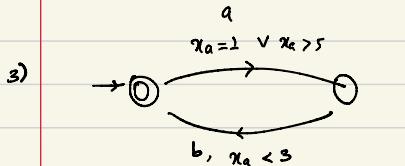
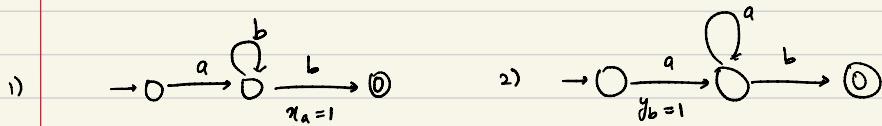


Recalling the subset idea:



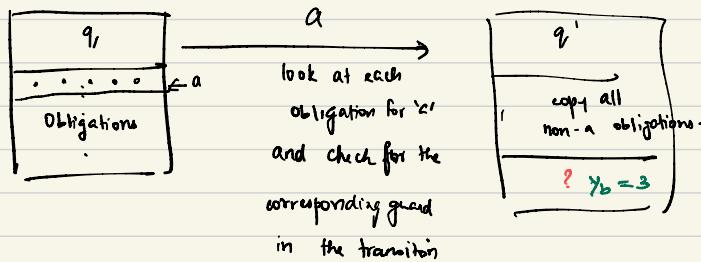
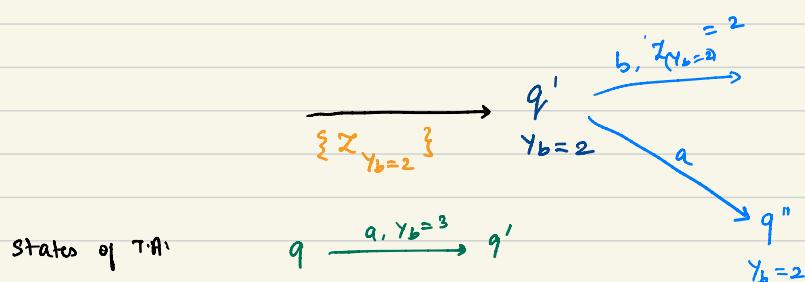
$$g_1 \wedge g_2 \wedge !g_3 \wedge !g_4 \wedge !g_5 \wedge !g_6.$$

Converting ECA to timed automata:



Converting EPA to NTA:

General idea: $\xrightarrow{y_b=2} q'$



* If obligation is $y_b = 1$, then no 'b' transition should exist from that state.

* What are accepting states?

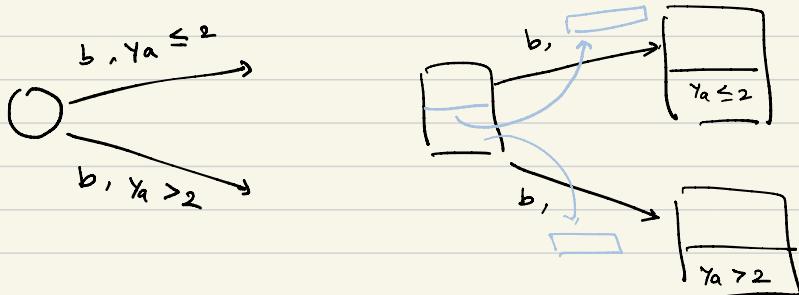


$q \in F$ (in the EPA)

only obligations are of the form $y_b = 1$

Question: In this conversion from EPA to T.A:

will a deterministic EPA be converted to a det. T.A.?



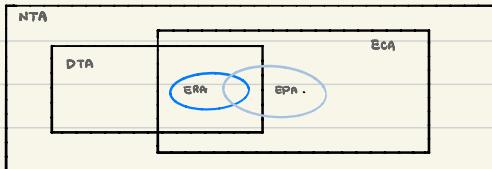
- The determinism existing in the EPA gets destroyed by this construction.

Question: Compare the languages accepted by EPA and det. T.A.

Is $EPA \subseteq DTA$? No:

Construct a language recognized by an EPA, but not by a Det. T.A.

Expressive Power of different classes:



1) $ERA \subseteq DTA$:

Without 1 guards: just adding $\{x_a\}$ in every transition on 'a' works.

- $ERA \rightarrow DTA$: ERA \rightarrow apply this transformation $\rightarrow DTA$

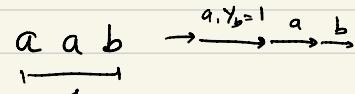
- We worry about 1 later.

2) $ERA \not\subseteq EPA$



3) $DTA \not\subseteq EPA$ ($\exists L$ in DTA which is not ERA recg.)

3) $EPA \not\subseteq ERA$



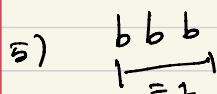
4) $EPA \not\subseteq DTA$

$\exists a \in L$ which EPA
but not by DTA

5) $DTA \not\subseteq ECA$

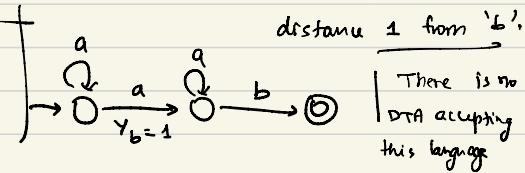
$\exists L$ in DTA
which is not
ECA recg.

6) $ECA \not\subseteq DTA$ $\exists L$ in ECA
which is not
DTA recg.



7) $a^k b$ s.t. there exist
some 'a' which is at
distance 1 from 'b'.

7) $ECA \not\subseteq EPA \cup ERA$



Summary of ECA

- Determinizable class
- We have compared expressive power

$$\begin{array}{c} \text{ECA} \\ L(B) \subseteq \boxed{L(A)} \\ \downarrow \\ \text{deterministic} \\ \downarrow \\ \text{complement} \\ \downarrow \\ \text{take intersection} \\ \text{with } L(B) \end{array}$$

$$q_0 \xrightarrow{a, \times ! = \delta} \textcircled{0}$$

