Unit-6: Model-checking ω -regular properties

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NPTEL-course

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Module 2:

ω -regular expressions to NBA

$\Sigma = \{ a, b \}$

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Example 1: Infinite word consisting only of a
                                      a^{\omega}
                Example 2: Infinite words containing only a or only b \ a^{\omega} + b^{\omega}
         { aaaaaaaaaaaaaaa..., bbbbbbbbbbbbbb...}
                                      аа\Sigma^*аа \cdot b^{\omega}
 Example 3: a word in aa\Sigma^*aa followed by only b-s
 Example 4: Infinite words where b occurs only finitely often (a + b)^* \cdot a^{\omega}
(a^*b)^{\omega}
 Example 5: Infinite words where b occurs infinitely often
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ω -regular expressions

$$G = E_1 \cdot F_1^{\omega} + E_2 \cdot F_2^{\omega} + \cdots + E_n \cdot F_n^{\omega}$$

$$E_1, \ldots, E_n, F_1, \ldots, F_n$$
 are regular expressions and $\epsilon \notin L(F_i)$ for all $1 \le i \le n$

$$L(F^{\omega}) = \{ w_1 w_2 w_3 \dots \mid \text{each } w_i \in L(F) \}$$

More examples

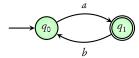
- $(a+b)^{\omega}$ set of all infinite words
- $a(a+b)^{\omega}$ infinite words starting with an a
- $(a + bc + c)^{\omega}$ words where every b is immediately followed by c
- $(a+b)^*c(a+b)^{\omega}$ words with a single occurrence of c
- $((a+b)^*c)^{\omega}$ words where *c* occurs infinitely often

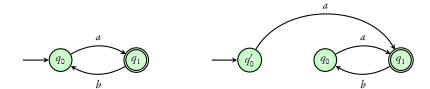
ω -regular expressions

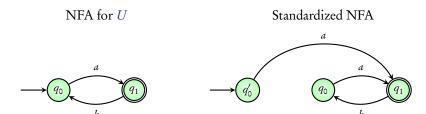
$$G = E_1 \cdot F_1^{\omega} + E_2 \cdot F_2^{\omega} + \cdots + E_n \cdot F_n^{\omega}$$

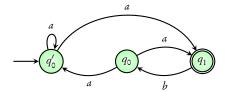
Goal: Convert ω -regular expression to NBA

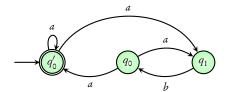
Part 1: Given regular expression U, find NBA for U^{ω}

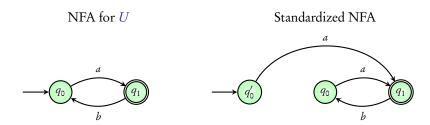


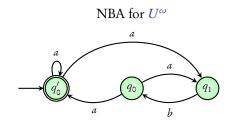


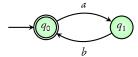


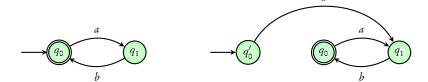


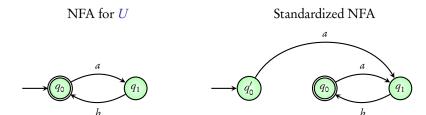




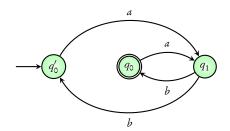




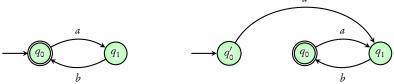


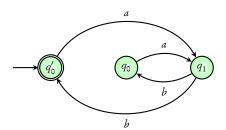


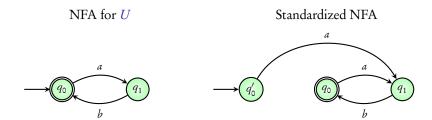
NFA for U Standardized NFA $\begin{array}{c} a \\ \hline q_0 \\ \hline q_1 \end{array}$ Standardized NFA

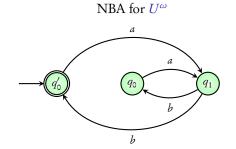


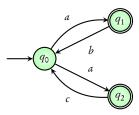
Standardized NFA

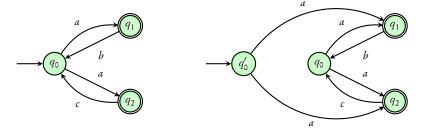


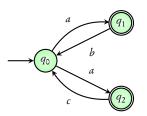


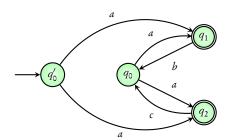


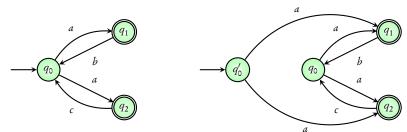


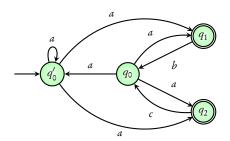


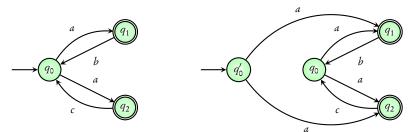


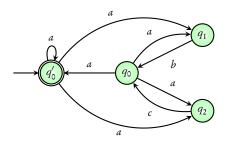




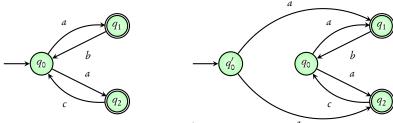




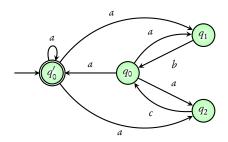


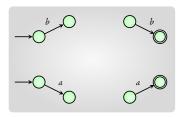


NFA for *U*

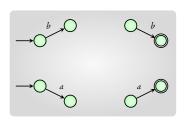


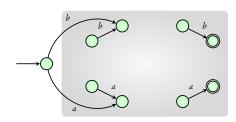
NBA for U^{ω}



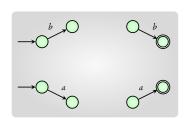


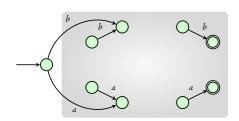
Standardized NFA for *U*



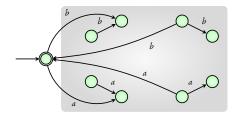


Standardized NFA for *U*





NBA for U^{ω}



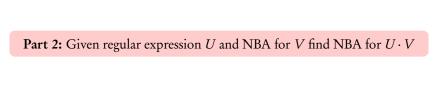
ω -regular expressions

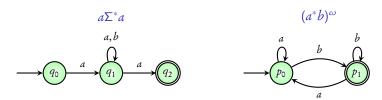
$$G = E_1 \cdot F_1^{\omega} + E_2 \cdot F_2^{\omega} + \cdots + E_n \cdot F_n^{\omega}$$

Goal: Convert ω -regular expression to NBA

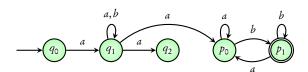
Part 1: Given regular expression U, find NBA for U^{ω}

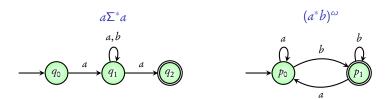
Done!



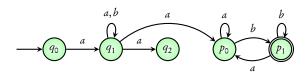


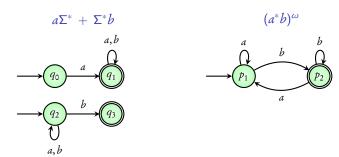


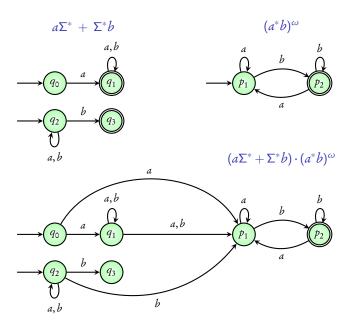




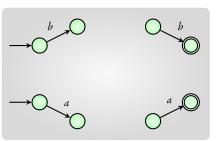


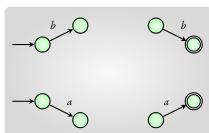




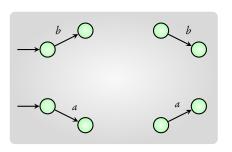


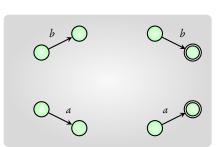
U



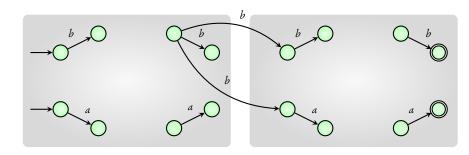


U V

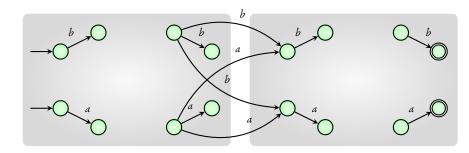




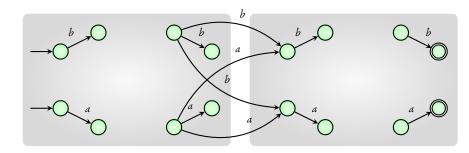
U



U V



$U \cdot V$



ω -regular expressions

$$G = E_1 \cdot F_1^{\omega} + E_2 \cdot F_2^{\omega} + \cdots + E_n \cdot F_n^{\omega}$$

Goal: Convert ω -regular expression to NBA

Part 1: Given regular expression U, find NBA for U^{ω}

Part 2: Given regular expression U and NBA for V find NBA for $U \cdot V$

Done!

Part 3: Given NBA for U and NBA for V find NBA for U + V

Part 3: Given NBA for U and NBA for V find NBA for U + V

Union of NBA already seen in Unit 5

Part 1: Given regular expression U, find NBA for U^{ω}

Part 2: Given regular expression U and NBA for V find NBA for $U \cdot V$

Part 3: Given NBA for U and NBA for V find NBA for U + V

Theorem

Every ω -regular expression can be converted to an NBA