# Unit-11: Binary Decision Diagrams (BDDs)

B. Srivathsan

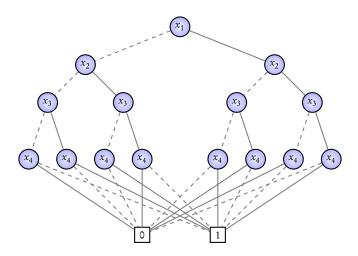
Chennai Mathematical Institute

NPTEL-course

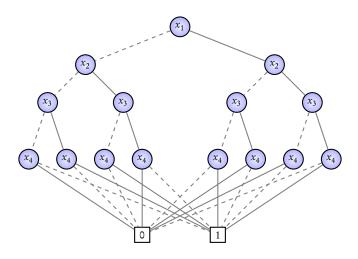
July - November 2015

# Module 2: Ordered BDDs

$$f(x_1, x_2, x_3, x_4) = \begin{cases} 1 & \text{if an even number of variables is 1} \\ 0 & \text{otherwise} \end{cases}$$

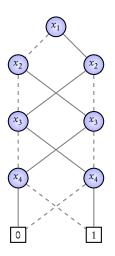


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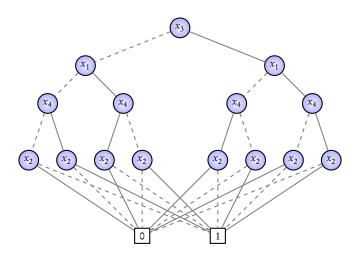
**Ordered** BDD for f with order  $[x_1, x_2, x_3, x_4]$ 

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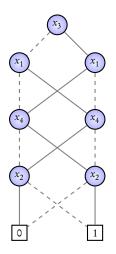
**Reduced Ordered** BDD for f with order  $[x_1, x_2, x_3, x_4]$ 

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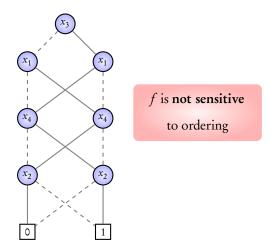
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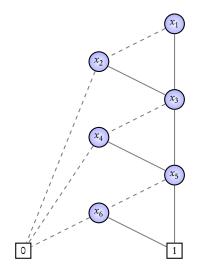
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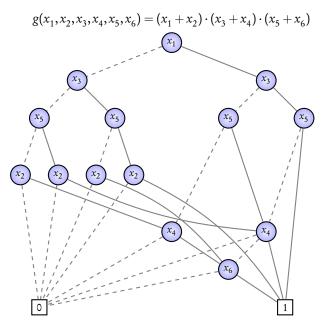


**Reduced Ordered** BDD for f with order  $[x_3, x_1, x_4, x_2]$ 

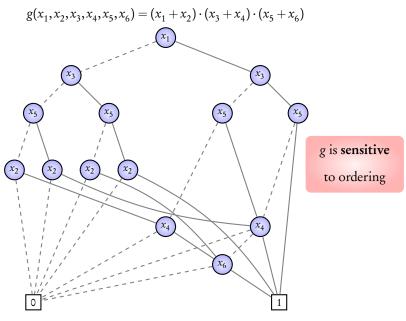
$$g(x_1, x_2, x_3, x_4, x_5, x_6) = (x_1 + x_2) \cdot (x_3 + x_4) \cdot (x_5 + x_6)$$



**Reduced Ordered BDD (ROBDD)** for g with order  $[x_1, x_2, x_3, x_4, x_5, x_6]$ 



**ROBDD** for g with order  $[x_1, x_3, x_5, x_2, x_4, x_6]$ 



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#### Ordered BDDs

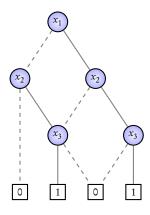
- ► BDDs with a **specified ordering** of variables
- ► For a given ordering, the reduced OBDD is **unique**
- Size of OBDD depends on the chosen ordering
- ► In practice, **heuristics** exist to find good orderings

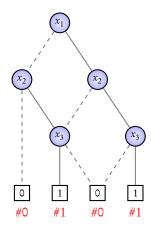
#### Ordered BDDs

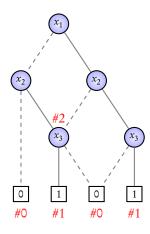
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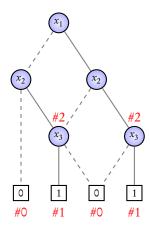
Coming next: Operations on OBDDs

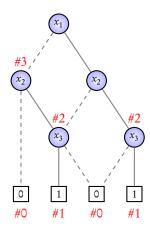
## Algorithm to reduce an OBDD

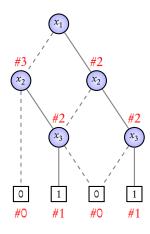


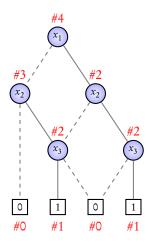


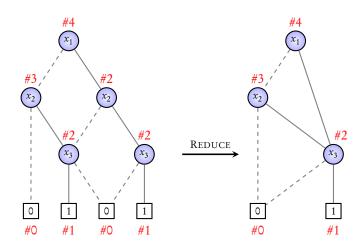


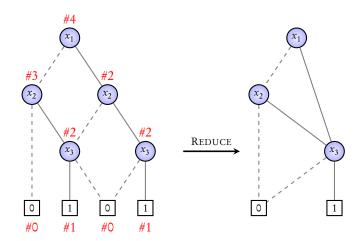


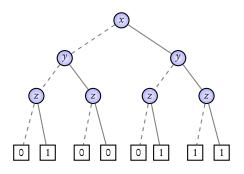


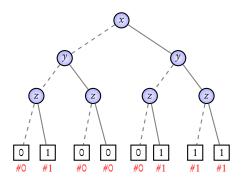


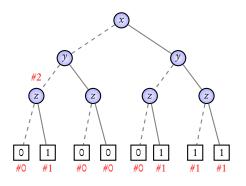


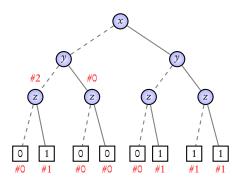


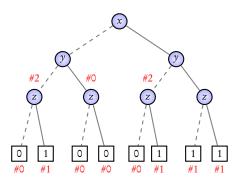


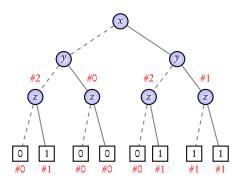


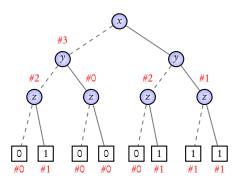


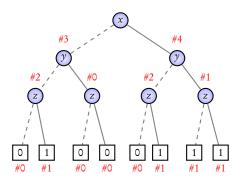


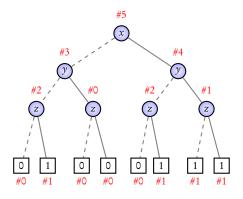


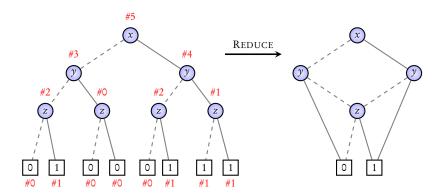












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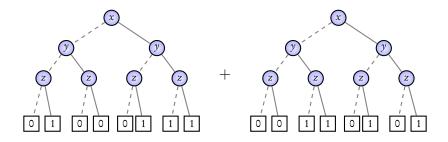
- ▶ Leaves: Label all 0 leaves with #0 and all 1 leaves with #1
- ▶ Intermediate node n
  - ► If 0-child and 1-child of *n* have same label, set label of *n* to be that label
  - ▶ If there is **another** node *m* such that *m* has the **same variable** *x*<sub>i</sub> and the **children** of *n* and *m* have **same** label, then set label of *n* to be the label of *m*

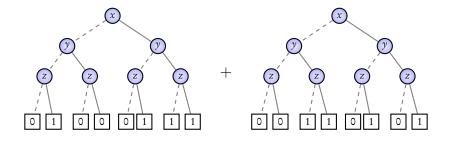
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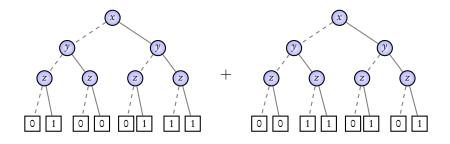
**Reference:** Logic in Computer Science, 2<sup>nd</sup> edition, by *Huth* and *Ryan*Section 6.2.1

#### Coming next: Algorithm for $OBDD_1 + OBDD_2$

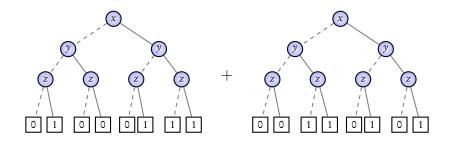


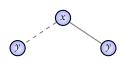


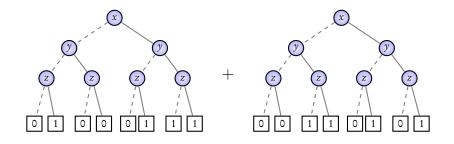


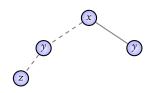


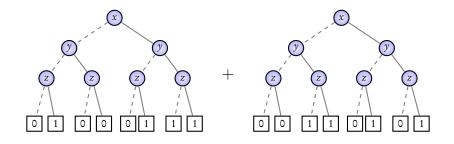


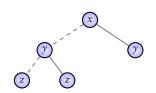


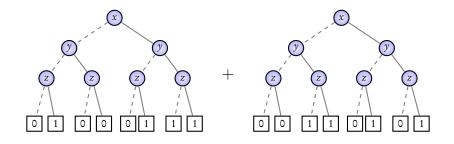


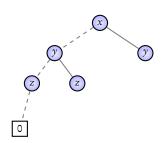


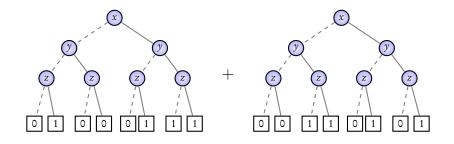


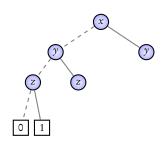


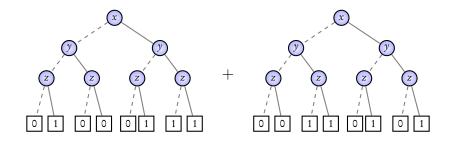


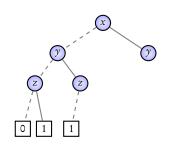


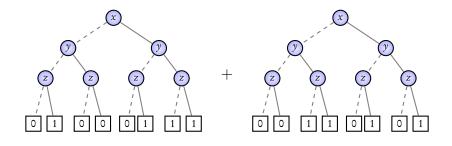


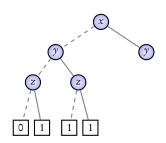


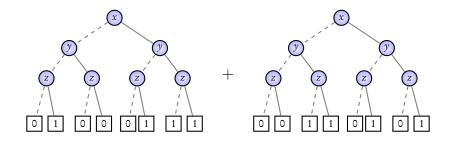


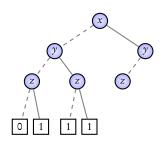


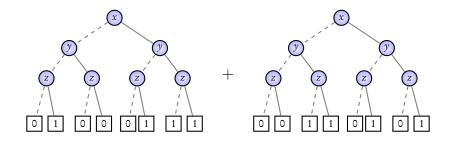


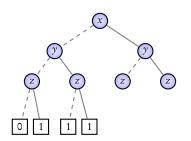


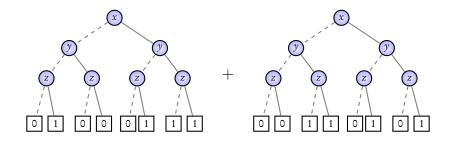


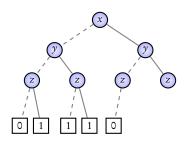


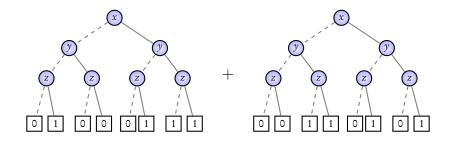


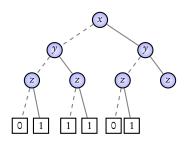


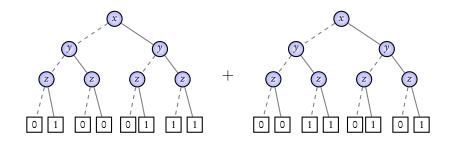


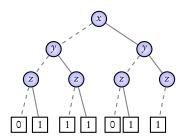


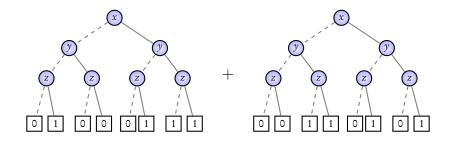


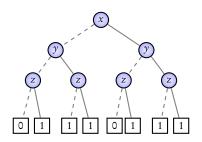


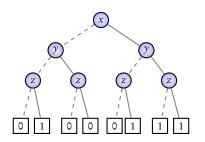


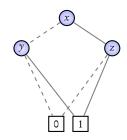


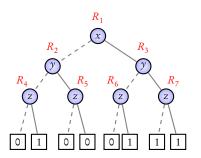


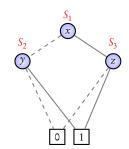


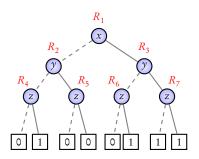


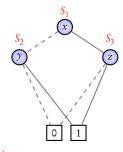




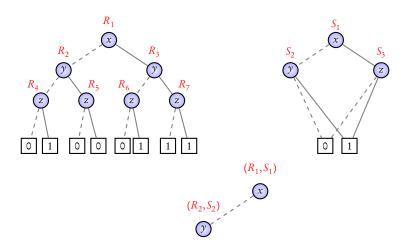


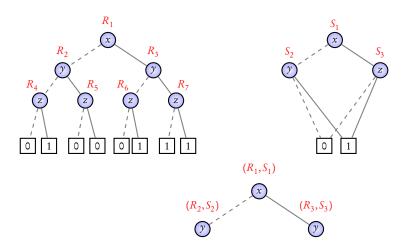


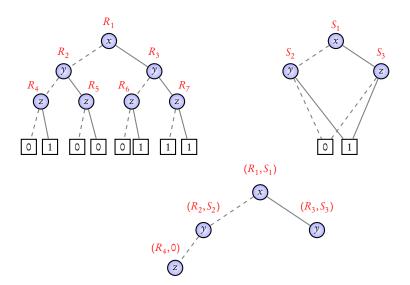


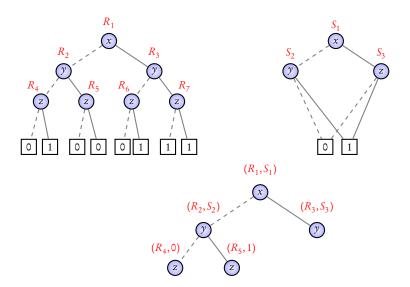


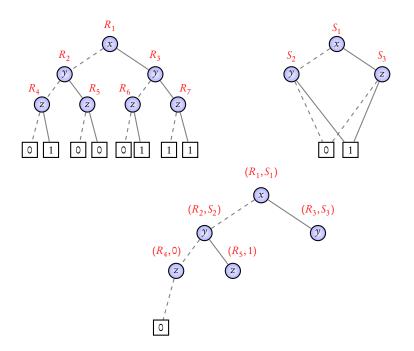


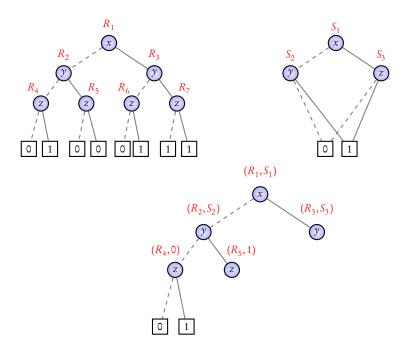


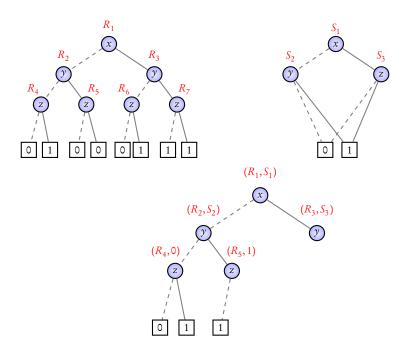


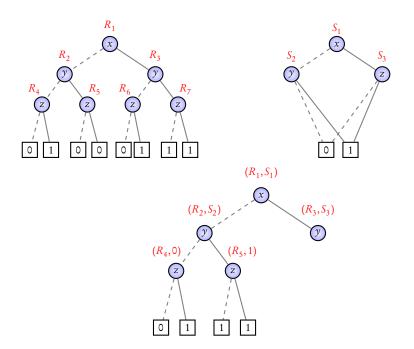


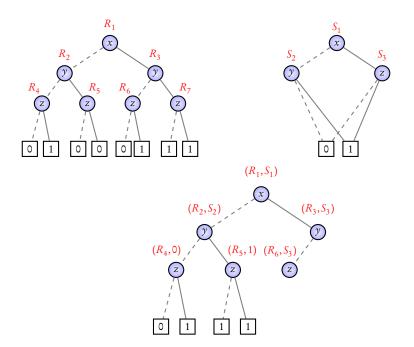


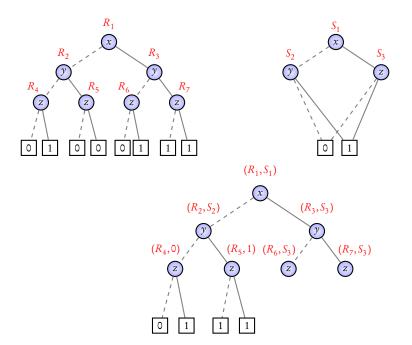


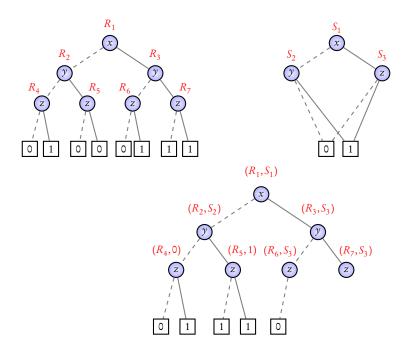


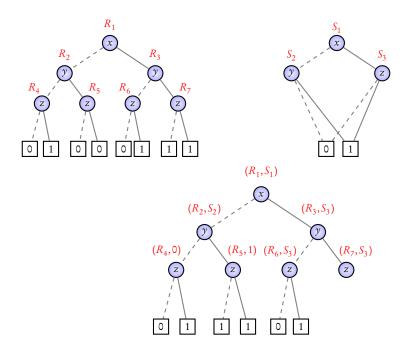


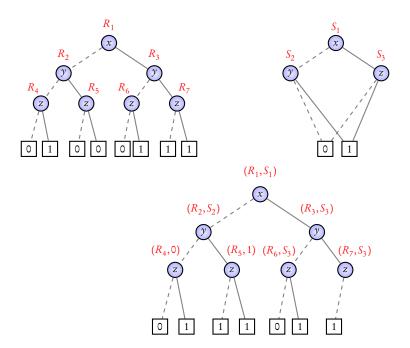


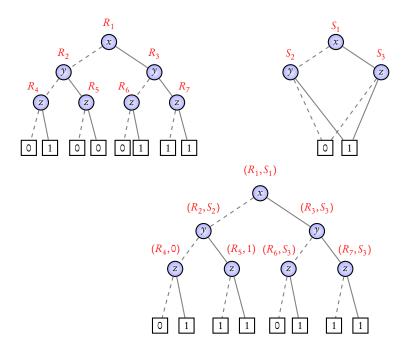


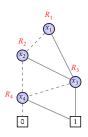


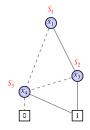


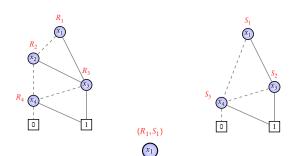


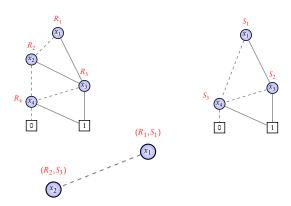


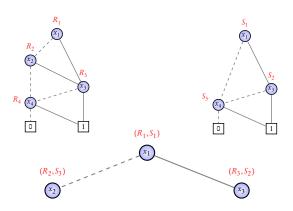


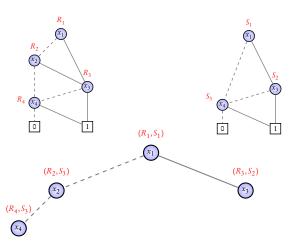


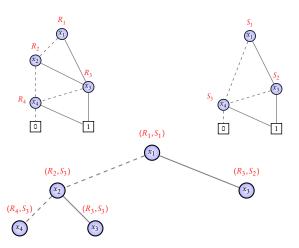


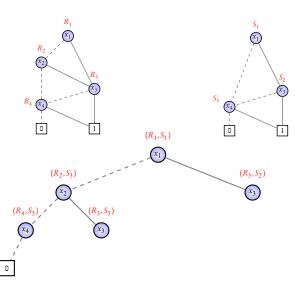


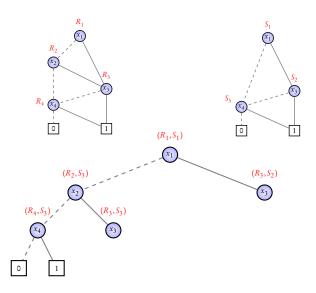


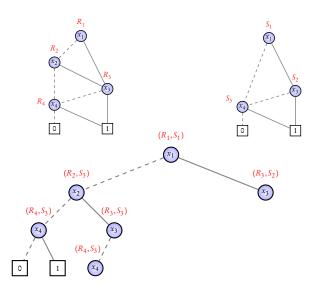


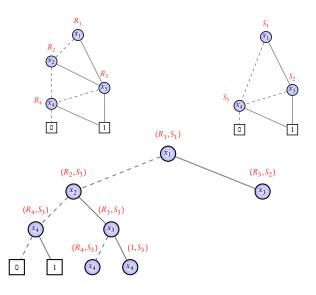


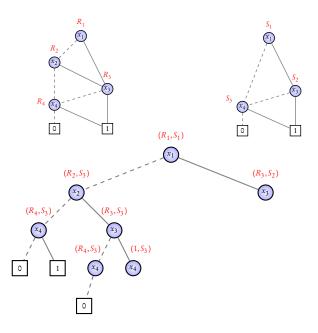


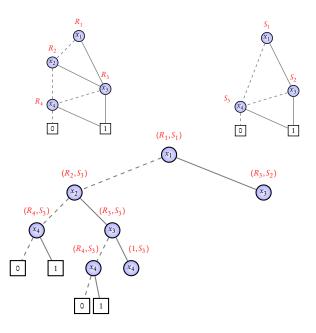


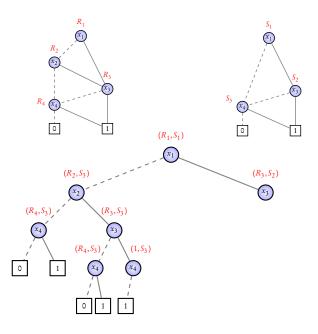


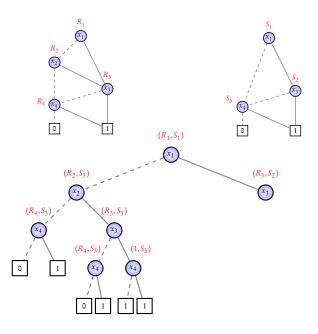


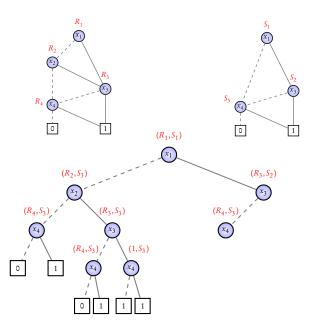


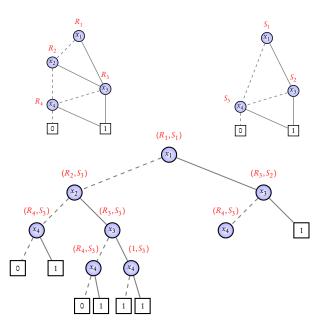


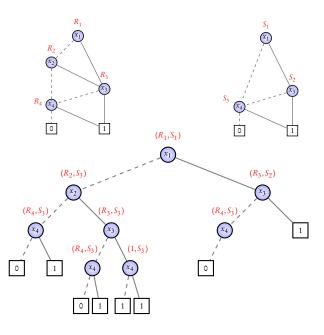


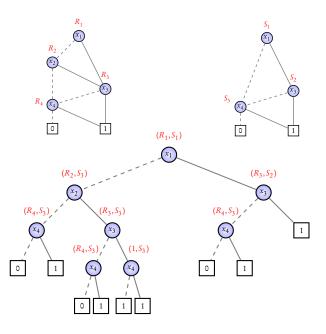


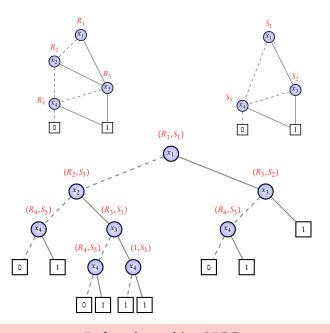












# Algorithm for OBDD<sub>1</sub> + OBDD<sub>2</sub>

# Algorithm for OBDD<sub>1</sub> + OBDD<sub>2</sub>

apply(+,r,s)

$$apply(+,r,s)$$

▶ If both r and s are terminals, create a terminal node r + s

$$apply(+,r,s)$$

- ▶ If both r and s are terminals, create a terminal node r + s
- ▶ If both r and s are  $x_i$  nodes, create an  $x_i$  node with:

$$apply(+,r,s)$$

- ▶ If both r and s are terminals, create a terminal node r + s
- ▶ If both r and s are  $x_i$  nodes, create an  $x_i$  node with:
  - ▶ left child: apply(+,left(r),left(s))

$$apply(+,r,s)$$

- ▶ If both r and s are terminals, create a terminal node r + s
- ▶ If both r and s are  $x_i$  nodes, create an  $x_i$  node with:
  - ▶ left child: apply(+,left(r),left(s))
  - ▶ right child: apply(+,right(r),right(s))

$$apply(+,r,s)$$

- ▶ If both r and s are terminals, create a terminal node r + s
- ▶ If both r and s are  $x_i$  nodes, create an  $x_i$  node with:
  - ▶ left child: apply(+,left(r),left(s))
  - ▶ right child: apply(+,right(r),right(s))
- ▶ If r is  $x_i$  node and s is a terminal or an  $x_j$  node with j > i, create an  $x_i$  node with:

$$apply(+,r,s)$$

- ▶ If both r and s are terminals, create a terminal node r + s
- ▶ If both r and s are  $x_i$  nodes, create an  $x_i$  node with:
  - left child: apply(+,left(r),left(s))
    right child: apply(+,right(r),right(s))
  - right child: appry(+,right(7),right(3))
- ▶ If r is  $x_i$  node and s is a terminal or an  $x_j$  node with j > i, create an  $x_i$  node with:
  - left child: apply(+, left(r), s)

$$apply(+,r,s)$$

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- ▶ If s is  $x_i$  node and r is a terminal: similar to Case 3

$$apply(+,r,s)$$

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- If s is  $x_i$  node and r is a terminal: similar to Case 3

 $OBDD_1 + OBDD_2$ : apply(+, root<sub>1</sub>, root<sub>2</sub>) and then reduce

#### Operations on BDDs

```
ightharpoonup OR: apply(+,root_1,root_2)
```

▶ AND: apply( $\cdot$ , root<sub>1</sub>, root<sub>2</sub>)

► XOR: apply(XOR, root<sub>1</sub>, root<sub>2</sub>)

▶ NOT: Use the fact that  $\overline{f} = f$  XOR 1

#### **OBDDs**

Reduction algorithm

Operations on OBDDs