## Unit-11: Binary Decision Diagrams (BDDs)

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## Module 2: <br> Ordered BDDs

$f\left(x_{1}, x_{2}, x_{3}, x_{4}\right)= \begin{cases}1 & \text { if an even number of variables is } 1 \\ 0 & \text { otherwise }\end{cases}$


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Ordered BDD for $f$ with order $\left[x_{1}, x_{2}, x_{3}, x_{4}\right]$

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f\left(x_{1}, x_{2}, x_{3}, x_{4}\right)= \begin{cases}1 & \text { if an even number of variables is } 1 \\ 0 & \text { otherwise }\end{cases}
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Reduced Ordered BDD for $f$ with order $\left[x_{1}, x_{2}, x_{3}, x_{4}\right]$

$$
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Ordered BDD for $f$ with order $\left[x_{3}, x_{1}, x_{4}, x_{2}\right]$

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f\left(x_{1}, x_{2}, x_{3}, x_{4}\right)= \begin{cases}1 & \text { if an even number of variables is } 1 \\ 0 & \text { otherwise }\end{cases}
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Reduced Ordered BDD for $f$ with order $\left[x_{3}, x_{1}, x_{4}, x_{2}\right]$

$$
f\left(x_{1}, x_{2}, x_{3}, x_{4}\right)= \begin{cases}1 & \text { if an even number of variables is } 1 \\ 0 & \text { otherwise }\end{cases}
$$


$f$ is not sensitive to ordering

Reduced Ordered BDD for $f$ with order $\left[x_{3}, x_{1}, x_{4}, x_{2}\right]$

$$
g\left(x_{1}, x_{2}, x_{3}, x_{4}, x_{5}, x_{6}\right)=\left(x_{1}+x_{2}\right) \cdot\left(x_{3}+x_{4}\right) \cdot\left(x_{5}+x_{6}\right)
$$



Reduced Ordered BDD (ROBDD) for $g$ with order $\left[x_{1}, x_{2}, x_{3}, x_{4}, x_{5}, x_{6}\right]$


ROBDD for $g$ with order $\left[x_{1}, x_{3}, x_{5}, x_{2}, x_{4}, x_{6}\right]$


ROBDD for $g$ with order $\left[x_{1}, x_{3}, x_{5}, x_{2}, x_{4}, x_{6}\right]$

## Ordered BDDs

- BDDs with a specified ordering of variables
- For a given ordering, the reduced OBDD is unique
- Size of OBDD depends on the chosen ordering
- In practice, heuristics exist to find good orderings


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## Coming next: Operations on OBDDs

## Algorithm to reduce an OBDD





















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- If 0 -child and 1 -child of $n$ have same label, set label of $n$ to be that label


## Algorithm to reduce OBDD

- Leaves: Label all 0 leaves with \#0 and all 1 leaves with \#1
- Intermediate node $n$
- If 0 -child and 1 -child of $n$ have same label, set label of $n$ to be that label
- If there is another node $m$ such that $m$ has the same variable $x_{i}$ and the children of $n$ and $m$ have same label, then set label of $n$ to be the label of $m$


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Reference: Logic in Computer Science, $2^{\text {nd }}$ edition, by Huth and Ryan Section 6.2.1

## Coming next: Algorithm for $\mathrm{OBDD}_{1}+\mathrm{OBDD}_{2}$



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$\left(R_{1}, S_{1}\right)$
(


$$
{ }_{\left(R_{2}, S_{2}\right)}^{\overbrace{}^{(2)}},{ }^{\left(R_{1}, S_{1}\right)}
$$
















$\bigcirc$


















Reduce the resulting OBDD

## Algorithm for $\mathrm{OBDD}_{1}+\mathrm{OBDD}_{2}$

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$$
\operatorname{apply}(+, r, s)
$$

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- If both $r$ and $s$ are terminals, create a terminal node $r+s$
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- If $r$ is $x_{i}$ node and $s$ is a terminal or an $x_{j}$ node with $j>i$, create an $x_{i}$ node with:


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- If $s$ is $x_{i}$ node and $r$ is a terminal: similar to Case 3


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- If $s$ is $x_{i}$ node and $r$ is a terminal: similar to Case 3
$\mathrm{OBDD}_{1}+\mathrm{OBDD}_{2}: \operatorname{apply}\left(+\right.$, root $_{1}$, root $\left._{2}\right)$ and then reduce


## Operations on BDDs

- OR: $\operatorname{apply}\left(+, \operatorname{root}_{1}, \operatorname{root}_{2}\right)$
- AND: $\operatorname{apply}\left(\cdot, \operatorname{root}_{1}\right.$, root $\left._{2}\right)$
- XOR: $\operatorname{apply}\left(\mathrm{XOR}, \operatorname{root}_{1}, \operatorname{root}_{2}\right)$
- NOT: Use the fact that $\bar{f}=f$ XOR 1


## OBDDs

## Reduction algorithm

Operations on OBDDs

