

1. Consider the algorithm Büchi-win in Figure 3.1 (Page 79) of the text. Design a game  $G$  for which Büchi-win makes at least 2 recursive calls.
2. Design a game  $G$  with  $O(n)$  vertices for which Büchi-win makes at least  $n$  recursive calls.
3. Consider the following algorithm for a mixed Büchi and co-Büchi objective: two disjoint sets  $T$  and  $S$  are specified; Player 1 wins plays that visit  $T$  infinitely often and  $S$  only finitely often. Design an algorithm to compute winning regions of each player in this mixed objective game.
4. Consider the mixed Büchi and co-Büchi objective mentioned in the previous question. Give an example of a game graph on which the algorithm given below does not terminate.

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1 algorithm mixed-win( $G$ )
2   if  $reach_1(T) = V$  and  $S = \emptyset$ 
3   then  $(W_0, W_1) = (\emptyset, V)$ 
4   else
5      $W'_0 = V \setminus reach_1(T)$ 
6      $G' = V \setminus reach_0(W'_0)$ 
7      $(W''_0, W''_1) = \text{mixed-win}(G')$ 
8      $(W_0, W_1) = (V \setminus W''_1, W''_1)$ 
9   endif
10  return  $(W_0, W_1)$ 
```