

Programming in Haskell: Lecture 11

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Combining elements

- Sum all numbers in a list

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sum :: [Int] -> Int
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- What is the common pattern?

Combining elements

- Combining elements using v and f

```
combine :: (Int -> Int -> Int) -> Int -> [Int] -> Int
combine f v []      = v
combine f v (x:xs) = f x (combine f v xs)
```

Combining elements

- Combining elements using v and f

```
combine :: (Int -> Int -> Int) -> Int -> [Int] -> Int
combine f v []      = v
combine f v (x:xs) = f x (combine f v xs)
```

- Sum and product can be expressed as:

```
sum      = combine (+) 0
product = combine (*) 1
```

foldr

- Built-in combine is called **foldr** (fold right)

```
foldr :: (a -> b -> b) -> b -> [a] -> b
```

```
foldr f v [] = v
```

```
foldr f v (x:xs) = f x (foldr f v xs)
```

```
foldr f v [x1, x2, x3]
= f x1 (foldr f v [x2, x3])
= f x1 (f x2 (foldr f v [x3]))
= f x1 (f x2 (f x3 (foldr f v [])))
= f x1 (f x2 (f x3 v))
```

foldr

- **foldr** replaces `[]` by `v` and `:` by ``f`` in the list:

```
xs
```

```
= x1 : (x2 : (x3 : (... : xn-1 : (xn : []))))
```

```
foldr f v xs
```

```
= x1 `f` (x2 `f` (x3 `f` (... `f` xn-1 `f` (xn `f` v))))
```


foldr examples

- Sum and product

```
sum      = foldr (+) 0
```

```
product = foldr (*) 1
```

```
xs
```

```
= x1 : (x2 : (x3 : (... : xn-1 : (xn : []))))
```

```
sum xs
```

```
= x1 + (x2 + (x3 + (... + xn-1 + (xn + 0))))
```

```
product xs
```

```
= x1 * (x2 * (x3 * (... * xn-1 * (xn * 1))))
```

foldr and anonymous functions

- Can express **length** in terms of **foldr**

```
length = foldr f 0
```

```
  where
```

```
    f x n = n+1
```

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- Impedes readability sometimes

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```
length = foldr f 0
```

```
  where
```

```
    f x n = n+1
```

- Not always convenient to name such functions
- Impedes readability sometimes
- **Anonymous functions:**

```
length = foldr (\x n -> n+1) 0
```

Anonymous functions

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f = \x n -> n+1
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- $\lambda x n \rightarrow n+1$ is an **unnamed** function of two arguments that increments its second argument by 1
- Can also give it a name!
- The two definitions of **f** below are equivalent:

```
f = \x n -> n+1
f x n = n+1
```

- Anonymous functions are very convenient to use with higher order functions

More `foldr` examples

- `foldr (:) []` is equivalent to the identity function on lists

More **foldr** examples

- **foldr** (:) [] is equivalent to the identity function on lists
- $f = \mathbf{foldr} (\backslash x\ l \rightarrow l++[x]) []$

```
f [x1, x2, x3]
= (f [x2, x3]) ++ [x1]
= ((f [x3]) ++ [x2]) ++ [x1]
= (((f []) ++ [x3]) ++ [x2]) ++ [x1]
= (([] ++ [x3]) ++ [x2]) ++ [x1]
= [x3, x2, x1]
```

More `foldr` examples

- `foldr (:) []` is equivalent to the identity function on lists
- `f = foldr (\x l -> l++[x]) []`

```
f [x1, x2, x3]
= (f [x2, x3])      ++ [x1]
= ((f [x3])        ++ [x2]) ++ [x1]
= (((f [])         ++ [x3]) ++ [x2]) ++ [x1]
= (([]             ++ [x3]) ++ [x2]) ++ [x1]
= [x3, x2, x1]
```

- `f` is just **reverse**, but takes time proportional to n^2

More `foldr` examples

- `foldr (:) []` is equivalent to the identity function on lists
- `f = foldr (\x l -> l++[x]) []`

```
f [x1, x2, x3]
= (f [x2, x3]) ++ [x1]
= ((f [x3]) ++ [x2]) ++ [x1]
= (((f []) ++ [x3]) ++ [x2]) ++ [x1]
= (([] ++ [x3]) ++ [x2]) ++ [x1]
= [x3, x2, x1]
```

- `f` is just `reverse`, but takes time proportional to n^2
- `concat` is just `foldr (++) []`

foldr1

- Sometimes there is no natural value to assign to the empty list

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- Sometimes there is no natural value to assign to the empty list
- For example, finding the maximum value in a list
- Maximum is undefined for empty list
- We use `foldr1` in such cases
- Uses the last element as initial value

```
foldr1 :: (a -> a -> a) -> [a] -> a
```

```
foldr1 f [x]      = x
```

```
foldr1 f (x:xs) = f x (foldr1 f xs)
```

```
maximum = foldr1 max
```

Folding from the left

- Sometimes it is useful to fold from the left

```
foldl :: (b -> a -> b) -> b -> [a] -> b
```

```
foldl f v []      = v
```

```
foldl f v (x:xs) = foldl f (f v x) xs
```

```
foldl f v [x1,x2,...,xn-1,xn]
```

```
  = (((v `f` x1) `f` x2) ... xn-1) `f` xn
```

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- Sometimes it is useful to fold from the left

```
foldl :: (b -> a -> b) -> b -> [a] -> b
```

```
foldl f v [] = v
```

```
foldl f v (x:xs) = foldl f (f v x) xs
```

```
foldl f v [x1,x2,...,xn-1,xn]
```

```
= (((v `f` x1) `f` x2) ... xn-1) `f` xn
```

- Translate a string of digits to a number

```
strToNum :: String -> Int
```

```
strToNum = foldl (\n c -> 10*n + digitToInt c) 0
```

Folding from the left

- Let $g\ n\ c = 10 * n + \text{digitToInt } c$

Folding from the left

- Let $g\ n\ c = 10*n + \text{digitToInt } c$
- Here is how $\text{strToNum} = \text{foldl } g\ 0$ works

```
strToNum           "234"
= foldl g 0        "234"
= foldl g (g 0 '2') "34"
= foldl g (g (g 0 '2') '3') "4"
= foldl g (g (g (g 0 '2') '3') '4') ""
=
= g (g (g 0 '2') '3') '4'
= 10 * (g (g 0 '2') '3') + 4
= 10 * (10 * (g 0 '2') + 3) + 4
= 10 * (10 * (10*0 + 2) + 3) + 4
= 234
```

foldr

- Fold from the right using function f and initial value v

```
foldr :: (a -> b -> b) -> b -> [a] -> b
```

```
foldr f v [] = v
```

```
foldr f v (x:xs) = f x (foldr f v xs)
```

```
foldr f v [x1, x2, x3, ..., xn]
```

```
= f x1 (foldr f v [x2, x3, ..., xn])
```

```
= f x1 (f x2 (foldr f v [x3, ..., xn]))
```

```
= f x1 (f x2 (f x3 (foldr f v [x4, ..., xn])))
```

```
= ...
```

```
= f x1 (f x2 (f x3 (... (f xn (foldr f v [])) ...)))
```

```
= f x1 (f x2 (f x3 (... (f xn v) ...)))
```


foldr

- Fold from the right using function (+) and initial value 0

```
foldr (+) 0 [1..100]
= (+) 1 (foldr (+) 0 [2..100])
= (+) 1 ((+) 2 (foldr (+) 0 [3..100]))
= (+) 1 ((+) 2 ((+) 3 (foldr (+) 0 [4..100])))
= ...
= (+) 1 ((+) 2 ((+) 3 (... ((+) 100
                               (foldr (+) 0 [])) ...)))
= (+) 1 ((+) 2 ((+) 3 (... ((+) 100 0) ...)))
= ...
= 5050
```

foldr

- Fold from the right using function f and initial value v

```
foldr f v [x1, x2, x3, ..., xn]
= ...
= f x1 (f x2 (f x3 (... (f xn v) ...)))
```

foldr

- Fold from the right using function `f` and initial value `v`

```
foldr f v [x1, x2, x3, ..., xn]
= ...
= f x1 (f x2 (f x3 (... (f xn v) ...)))
```

- If `f` needs both inputs, it will be applied only at the end

foldr

- Fold from the right using function `f` and initial value `v`

```
foldr f v [x1, x2, x3, ..., xn]
= ...
= f x1 (f x2 (f x3 (... (f xn v) ...)))
```

- If `f` needs both inputs, it will be applied only at the end
- Need space to carry huge expressions around

foldl

- Fold from the left using function f and initial value v

```
foldl :: (b -> a -> b) -> b -> [a] -> b
```

```
foldl f v [] = v
```

```
foldl f v (x:xs) = foldl f (f v x) xs
```

```
foldl f v [x1, x2, x3, ..., xn]
```

```
= foldl f (f v x1) [x2, x3, ..., xn]
```

```
= foldl f (f (f v x1) x2) [x3, ..., xn]
```

```
= foldl f (f (f (f v x1) x2) x3) [x4, ..., xn]
```

```
= ...
```

```
= foldl f (f (... (f (f (f v x1) x2) x3) ...)) xn []
```

```
= f (... (f (f (f v x1) x2) x3) ...) xn
```

foldl

- Fold from the left using function (+) and initial value 0

```
foldl (+) 0 [1..100]
= foldl (+) ((+) 0 1) [2..100]
= foldl (+) ((+) ((+) 0 1) 2) [3..100]
= foldl (+) ((+) ((+) ((+) 0 1) 2) 3) [4..100]
= ...
= foldl (+) ((+) (... ((+) ((+) ((+) 0 1) 2) 3)
              ...)) 100)
      []
= (+) (... ((+) ((+) ((+) 0 1) 2) 3) ...) 100
```

foldl

- Fold from the left using function f and initial value v

```
foldl f v [x1, x2, x3, ..., xn]
= ...
= f (... (f (f (f v x1) x2) x3) ...) xn
```

foldl

- Fold from the left using function `f` and initial value `v`

```
foldl f v [x1, x2, x3, ..., xn]
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```

- Same problem as with `foldr`

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foldl f v [x1, x2, x3, ..., xn]
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```

- Same problem as with `foldr`
- Need space to carry huge expressions around

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- Defined in `Data.List`

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- **Eager version** of `foldl`

```
foldl' :: (b -> a -> b) -> b -> [a] -> b
foldl' f v []      = v
foldl' f v (x:xs) = y `seq` foldl' f y xs
  where y = f v x
```

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- `seq :: a -> b -> b`

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- `seq :: a -> b -> b`
- Evaluates the first argument first and then returns the second argument

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- `seq :: a -> b -> b`
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```

- `seq :: a -> b -> b`
- Evaluates the first argument first and then returns the second argument
- Useful when first argument is used in second argument
- Forces the values in `foldl'` to be evaluated as early as possible

foldl'

- Computing with `foldl'`:

```
foldl' f v [x1,x2,x3,...,xn]
= foldl' f y1 [x2,x3,...,xn] -- y1 = f v x1
= foldl' f y2 [x3,...,xn] -- y2 = f y1 x2
= foldl' f y3 [x4,...,xn] -- y3 = f y2 x3
= ...
= foldl' f yn [] -- yn = f yn-1 xn
= yn
```


foldl'

- `foldl' (+) 0`:

```
foldl' (+) 0 [1..100]
= foldl' (+) 1 [2..100]      -- 1 = (+) 0 1
= foldl' (+) 3 [3..100]     -- 3 = (+) 1 2
= foldl' (+) 6 [4..100]     -- 6 = (+) 3 3
= ...
= foldl' (+) 5050 []        -- 5050 = (+) 4950 100
= 5050
```

foldr on infinite lists

- **foldr** can be made to work on infinite lists

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- If **f** does not require the second argument, the fold can terminate
- A complicated **head**:

```
foldr (\x y -> x) 0 [1..]  
= (\x y -> x) 1 (foldr (\x y -> x) 0 [2..])  
= 1
```

foldr on infinite lists

- **foldr** can be made to work on infinite lists
- If **f** does not require the second argument, the fold can terminate
- A complicated **head**:

```
foldr (\x y -> x) 0 [1..]  
= (\x y -> x) 1 (foldr (\x y -> x) 0 [2..])  
= 1
```

- Does not work with left folds:

```
foldl' (\x y -> x) 0 [1..]  
= foldl' (\x y -> x) 0 [2..]  
= foldl' (\x y -> x) 0 [3..]  
= ...
```

Simulating `foldl` using `foldr`

- Let `step x g a = g (f a x)`

Simulating `foldl` using `foldr`

- Let `step x g a = g (f a x)`
- **Claim:** For all `g`, `xs` and `e`, `foldr step g xs e = g (foldl f e xs)`

Simulating `foldl` using `foldr`

- Let `step x g a = g (f a x)`
- **Claim:** For all `g`, `xs` and `e`, `foldr step g xs e = g (foldl f e xs)`
- **Proof:** By induction on `length xs`

```
foldr step g [] e = g e = g (foldl f e [])
```

```
foldr step g (x:xs) e
= step x (foldr step g xs) e
= foldr step g xs (f e x)
= g (foldl f (f e x) xs)
    -- (ind. hyp. applied on g, xs and (f e x))
= g (foldl f e (x:xs))
```


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- **Proof:** By induction on `length xs`

```
foldl step' g [] e = g e = g (foldr f e [])
```

```
foldl step' g (x:xs) e
= foldl step' (step' g x) xs e
= step' g x (foldr f e xs)
  -- (ind. hyp. applied on step' g x, xs and e)
= g (f x (foldr f e xs))
= g (foldr f e (x:xs))
```

Some useful functions

- `flip :: (a -> b -> c) -> b -> a -> c`

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- `foldr f v l` can be changed to `foldl (flip f) v l`

Some useful functions

- `flip :: (a -> b -> c) -> b -> a -> c`
- `flip f` behaves like `f`, but accepts the arguments in reverse order
- `flip (:) [1..10] 0 = [0..10]`
- `foldr f v l` can be changed to `foldl (flip f) v l`
- Other useful functions

```
const :: a -> b -> a
```

```
const x y = x
```

```
($) :: (a -> b) -> a -> b
```

```
($) f x = f x
```

```
($!) :: (a -> b) -> a -> b
```

```
($!) f x = x `seq` f x      -- Eager version
```


`foldl` using `foldr`, again

- For finite lists:

```
foldl f = flip (foldr step id)
  where step x g a = g (f a x)
```

```
flip (foldr step id) e xs
= foldr step id xs e
= id (foldl f e xs)
= foldl f e xs
```

foldr using **foldl**, again

- For finite lists:

```
foldr f = flip (foldl step' id)  
  where step' g x a = g (f x a)
```

```
flip (foldl step' id) e xs  
= foldl step' id xs e  
= id (foldr f e xs)  
= foldr f e xs
```