

Programming in Haskell

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LECTURE 14

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Enumerated data types

- * The `data` keyword is used to define new types
- * `data Bool = False | True`
- * `data Day = Sun | Mon | Tue | Wed | Thu | Fri | Sat`

Data types with parameters

- * data Shape = Circle Float
| Square Float
| Rectangle Float Float
- * Circle 5.0, Square 4.0, Rectangle 3.0 4.0

Functions on data types

- * Functions can be defined using pattern matching

- * `weekend :: Day -> Bool`

- `weekend Sat = True`

- `weekend Sun = True`

- `weekend _ = False`

- * `area :: Shape -> Float`

- `area (Circle r) = pi*r*r`

- `area (Square x) = x*x`

- `area (Rectangle l w) = l*w`

- where

- `pi = 3.1415927`

Functions on data types

- * What about
`weekend2 :: Day -> Bool`
- * `weekend2 d`
 - | `(d == Sat || d == Sun) = True`
 - | otherwise = False
- * Error - No instance for `(Eq Day)` arising from a use of
`'=='`

Functions on data types

- * How about this function?
- * `nextday :: Day -> Day`
`nextday Sun = Mon`
`nextday Mon = Tue`
...
`nextday Sat = Sun`
- * Invoking `nextday Fri` in `ghci` will lead to error
- * Error - No instance for `(Show Day)` arising from a use of `'print'`

Add data types to typeclasses

- * To check equality of two values of a data type `a`, `a` must belong the type class `Eq`
- * We add `Day` to the type class `Eq` as follows
- * `data Day = Sun | Mon | ... | Sat`
 `deriving Eq`
- * Default behaviour: `Sun == Sun, Tue /= Fri, ...`
- * Now `weekday2` compiles without error

The type class `Show`

- * To make `nextday` work, we must make `Day` an instance of `Show`
- *

```
data Day = Sun | Mon | ... | Sat
  deriving (Eq, Show)
```
- * The type class `Show` consists of all data types that implement the function `show`

More derivations

- * `show` converts its input to a string which can be printed on the screen
- * Default text representation
- * `show Wed == "Wed"`
- * `data Day = Sun | Mon | ... | Sat`
 `deriving (Eq, Show, Ord)`
- * `Sun < Mon < ... < Sat`

More derivations ...

- * `data Shape = Circle Float
 | Square Float
 | Rectangle Float Float
 deriving (Eq, Ord, Show)`
- * `show (Circle 5.0) == "Circle 5.0"`
- * `Square 4.0 == Square 4.0`
`Square 4.0 /= Square 3.0`
`Circle 5.0 /= Rectangle 3.0 4.0`
- * `Square 4.0 > Circle 5.0`

Constructors

- * Square, Circle, Sun, Mon, ... are constructors

- * They are functions

Sun :: Day

Rectangle :: Float -> Float -> Shape

Circle :: Float -> Shape

Constructors ...

- * Constructors can be used just like other functions
- * `Circle 5.0 :: Shape`
- * `map Circle :: [Float] -> [Shape]`
- * `map Circle [3.0, 2.0] = [Circle 3.0, Circle 2.0]`

Records

- * `data Person = Person String Int Float String
deriving Show`
- * `guy = Person "Alpha" 21 5.8 "+914427470226"`
- * `name :: Person -> String
name (Person n _ _ _) = n`
- * `age :: Person -> Int
age (Person _ a _ _) = a`

Records ...

- * `height :: Person -> Float`
`height (Person _ _ h _) = h`
- * `phone :: Person -> Int`
`phone (Person _ _ _ p) = p`

Record syntax

- * `data Person = Person { name :: String
 , age :: Int
 , height :: Float
 , phone :: String
 } deriving Show`
- * `guy = Person {name="Alpha", age = 21,
 height = 5.8, phone = "+914427470226"}`
- * The field names are actually functions
- * `name :: Person -> String
age :: Person -> Int`

Summary

- * The keyword `data` is used to declare new data types
- * The keyword `deriving` to derive as an instance of a type class
- * Data types with parameters - `Shape`, `Person`
- * Sum type or union - `Day`, `Shape`
- * Product type or struct - `Person`

Abstract data types

- * Consider a **Stack** data type
 - * a collection of **Ints** stacked one on top of the other
 - * **push**: place an element on top of the stack
 - * **pop**: remove the topmost element of the stack
- * Behaviour similar to lists

Abstract data types

- * `data Stack = Stack [Int]`
- * The value constructor `Stack` is a function that converts a list of `Int` to a `Stack` object
- * Internal representation hidden

Abstract data types

- * `empty :: Stack`
`empty = Stack []`
- * `push :: Int -> Stack -> Stack`
`push x (Stack xs) = Stack (x:xs)`
- * `pop :: Stack -> (Int, Stack)`
`pop (Stack (x:xs)) = (x, Stack xs)`
- * `isempty :: Stack -> Bool`
`isempty (Stack []) = True`
`isempty (Stack _) = False`

Type parameters

- * Polymorphic user-defined data types
- * `data Stack a = Stack [a]`
 `deriving (Eq, Show, Ord)`
- * `empty :: Stack a`
- * `push :: Int -> Stack a -> Stack a`
- * `pop :: Stack a -> (a, Stack a)`
- * `isempty :: Stack a -> Bool`

Type parameters...

- * Suppose we want to sum all elements in a stack
- * `sumStack (Stack xs) = sum xs`
- * What is the type of `sumStack`?
- * Applicable only if the stack has numeric elements
- * `sumStack :: (Num a) => Stack a -> a`

A custom show

- * `show (Stack [1,2,3]) == "Stack [1,2,3]"`
- * deriving `Show` defines a default implementation for `show`
- * Suppose we want something mildly fancy
- * `show (Stack [1,2,3]) == "1->2->3"`

A custom show

- * One can change the default behaviour
- *

```
printElems :: (Show a) => [a] -> String
printElems [] = ""
printElems [x] = show x
printElems (x:xs) = show x ++ "->" ++
                    printElems xs
```
- *

```
instance (Show a) => Show (Stack a) where
    show (Stack l) = printElems l
```


Queues

- * Consider a `Queue` data type
 - * a collection of `Ints` arranged in a sequence
 - * `enqueue`: add an element at the end of the queue
 - * `dequeue`: remove the element at the start of the queue

Queues

- * `data Queue a = Queue [a]`
- * `empty :: Queue a`
`empty = Queue []`
- * `isEmpty :: Queue a -> Bool`
`isEmpty (Queue []) = True`
`isEmpty (Queue _) = False`

Queues

- * enqueue :: a -> Queue a -> Queue a
enqueue x (Queue xs) = Queue (xs ++ [x])
- * dequeue :: Queue a -> (a, Queue a)
dequeue (Queue (x:xs)) = (x, Queue xs)

Queues

- * Each **enqueue** on a queue of length n takes $O(n)$ time
- * Enqueueing and dequeueing n elements might take $O(n^2)$ time

Efficient queue

- * Use two lists
- * Represent q_1, q_2, \dots, q_n as
 $[q_1, q_2, \dots, q_j]$ and $[q_n, q_{n-1}, \dots, q_{j+1}]$
- * Second list is the second part of queue in reversed order
- * **enqueue** adds an element at the start of the second list
- * **dequeue** removes an element from the start of the first list

Efficient queue

- * What if we try to dequeue when the first list is empty?
- * We reverse the second list into the first, and remove the first element

Efficient queue

- * If we add n elements, we get a queue
NuQu $[]$ $[q_n, q_{n-1}, \dots, q_1]$
- * Next **dequeue** takes $O(n)$ time to reverse the list
- * After one dequeue we get NuQu $[q_2, \dots, q_n]$ $[]$
- * Next $n-1$ **dequeue** operations take $O(1)$ time

Amortized analysis

- * How many times is an element touched?
 - * Once when it is added to the second list
 - * Twice when it is moved from the second to first
 - * Once when it is removed from the first list
- * Each element is touched at most four times
- * Any sequence of n instructions involves at most n elements
- * So any sequence of n instructions takes only $O(n)$ steps

Summary

- * Abstract data types
- * We can define polymorphic user-defined data types by supplying type parameters
- * Conditional polymorphism for functions defined on such data
- * The `instance` keyword to define non-default implementation of functions
- * Efficient queues and amortized analysis