

Programming in Haskell

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LECTURE 9

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Anonymous functions

- * Usual practice with functions
 - * Define functions – giving it a name
 - * Use them elsewhere
- * Sometimes it breaks the flow to follow this pattern
- * **Unnamed functions**

Anonymous functions

- * Example:

```
foldr f 0 [1..]  
  where f x y = x
```

- * Easier to say this:

```
foldr (\x y -> x) 0 [1..]
```

- * We are specifying the function we want to use without naming it
- * `\x y -> x` is a function that takes two inputs and returns the first input

Computations with foldr

* `foldr f a [x1, x2, ..., xn]`

⇒ `f x1 (foldr f a [x2, ..., xn])`

⇒ `f x1 (f x2 (foldr f a [x3, ..., xn]))`

⇒ `f x1 (f x2 (f x3 (foldr f a [x4, ..., xn])))`

⇒ ...

⇒ `f x1 (f x2 (f x3 (... (f xn (foldr f a []))...)))`

⇒ `f x1 (f x2 (f x3 (... (f xn a)...)))`

Computations with foldr

* foldr (+) 0 [1..100]

⇒ (+) 1 (foldr (+) 0 [2..100])

⇒ (+) 1 ((+) 2 (foldr (+) 0 [3..100]))

⇒ ...

⇒ (+) 1 ((+) 2 (... ((+) 100 (foldr (+) 0 []))...))

⇒ (+) 1 ((+) 2 (... ((+) 100 0)...))

⇒ ...

⇒ 5050

Computations with foldr

* `foldr f a [x1, x2, ..., xn]`

⇒ `f x1 (foldr f a [x2, ..., xn])`

⇒ ...

⇒ `f x1 (f x2 (f x3 (... (f xn a)...))`

- * If `f` needs both inputs, it will be applied only at the end
- * Need space to carry around huge expressions

Computations with foldl

* `foldl f a [x1, x2, ..., xn]`

⇒ `foldl f (f a x1) [x2, ..., xn]`

⇒ `foldl f (f (f a x1) x2) [x3, ..., xn]`

⇒ `foldl f (f (f (f a x1) x2) x3) [x4, ..., xn]`

⇒ ...

⇒ `foldl f (f ... (f (f (f a x1) x2) x3)) ... xn) []`

⇒ `f ... (f (f (f a x1) x2) x3) ... xn`

Computations with foldl

* foldl (+) 0 [1..100]

⇒ foldl (+) ((+) 0 1) [2..100]

⇒ foldl (+) ((+) ((+) 0 1) 2) [3..100]

⇒ ...

⇒ foldl (+) ((+) ... (+) ((+) 0 1) 2)... 100) []

⇒ (+) ... (+) ((+) 0 1) 2)... 100

⇒ ...

⇒ 5050

Computations with foldl

- * `foldl f a [x1, x2, ..., xn]`

 - ⇒ `foldl f (f a x1) [x2, ..., xn]`

 - ⇒ ...

 - ⇒ `f ... (f (f (f a x1) x2) x3) ... xn`

- * Same problem as with `foldr`

- * Huge expression carried around till the end

Computations with foldl'

* `foldl' f a [x1, x2, ..., xn]`

⇒ `foldl' f y1 [x2, ..., xn]`

– $y_1 = f\ a\ x_1$

⇒ `foldl' f y2 [x3, ..., xn]`

– $y_2 = f\ y_1\ x_2$

⇒ `foldl' f y3 [x4, ..., xn]`

– $y_3 = f\ y_2\ x_3$

⇒ ...

⇒ `foldl' f yn []`

– $y_n = f\ y_{(n-1)}\ x_n$

⇒ `yn`

* Eager evaluation

Computations with foldl'

* foldl' (+) 0 [1..100]

⇒ foldl' (+) 1 [2..100]

⇒ foldl' (+) 3 [3..100]

⇒ ...

⇒ foldl' 5050 []

⇒ 5050

Computations with foldl'

- * `foldl'` defined in **Data.List**
- * $\text{foldl}' f a [] = a$
 $\text{foldl}' f a (x:xs) = y \text{ `seq` foldl}' f y xs$
 where $y = f a x$
- * The `seq` function takes two arguments, evaluates the first, and returns the value of the second
- * $\text{seq} :: a \rightarrow b \rightarrow b$
- * Forces the values in `foldl'` to be computed as early as possible

foldr on infinite lists

- * `foldr` works on infinite lists sometimes when `foldl` or `foldl'` does not
- * `foldr (\x y -> x) 0 [1..]`
⇒ `(\x y -> x) 1 (foldr (\x y -> x) 1 [2..])`
⇒ 1
- * `foldl' (\x y -> x) 0 [1..]`
⇒ `foldl' (\x y -> x) 0 [2..]`
⇒ `foldl' (\x y -> x) 0 [3..]`
⇒ `foldl' (\x y -> x) 0 [4..]`
⇒ ...

foldl using foldr

- * Let $\text{step } x \ g = \lambda a \rightarrow g (f a \ x)$
- * **Claim:** For all expressions e ,
 $\text{foldr step id } xs \ e = \text{foldl f e } xs$
- * **Proof:** By induction on length of xs
 - * $(\text{foldr step id } []) \ e = \text{id } e = e = \text{foldl f e } []$
 - * $(\text{foldr step id } (x:xs)) \ e$
 - $\Rightarrow (\text{step } x (\text{foldr step id } xs)) \ e$
 - $\Rightarrow (\lambda a \rightarrow (\mathbf{\text{foldr step id } xs}) (\mathbf{f a \ x})) \ e$
 - $\Rightarrow (\lambda a \rightarrow \mathbf{\text{foldl f (f a x) xs}}) \ e$ – By induction hypothesis
 - $\Rightarrow \text{foldl f (f e x) } xs = \text{foldl f e } (x:xs)$

Useful functions

- * $\text{flip} :: (a \rightarrow b \rightarrow c) \rightarrow b \rightarrow a \rightarrow c$
- * If we have a definition $\text{foldr } f \ a \ l$ and want to change it to foldl , we do $\text{foldl } (\text{flip } f) \ a \ l$
- * $\text{const} :: a \rightarrow b \rightarrow a$
- * $\text{const } x \ y = x$
- * $(\$) :: (a \rightarrow b) \rightarrow a \rightarrow b$
 $(\$) \ f \ x = f \ x$
- * $(\$!) :: (a \rightarrow b) \rightarrow a \rightarrow b$
 $(\$!) \ f \ x = y \ \text{`seq`} \ f \ y$
 where $y = x$
behaviour
 - This is not the official definition
 - Only conveys the intended