

Programming in Haskell

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LECTURE 2

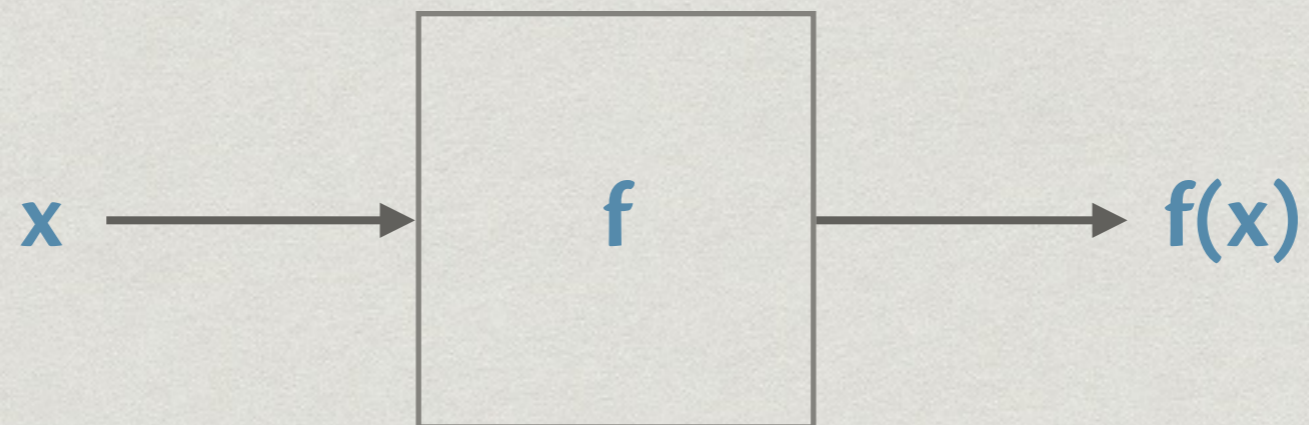
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Programs as functions

- * Functions transform inputs to outputs



- * **Program**: rules to produce output from input
- * **Computation**: process of applying the rules

Building up programs

How do we describe the rules?

- * Start with built in functions
- * Use these to build more complex functions

Building up programs ...

Suppose

- * ... we have the whole numbers, $\{0, 1, 2, \dots\}$
- * ... and the successor function, **SUCC**

$$\text{SUCC } 0 = 1$$

$$\text{SUCC } 1 = 2$$

$$\text{SUCC } 2 = 3$$

...

- * Note: we that write **succ 0**, not **succ(0)**

Building up programs ...

We can **compose** **succ** twice to build a new function

- * $\text{plusTwo } n = \text{succ } (\text{succ } n)$

If we compose **plusTwo** and **succ** we get

- * $\text{plusThree } n = \text{succ } (\text{plusTwo } n)$

Inductive/recursive definitions

- * $\text{plus } n \ 0 = n$, for every n
- * $\text{plus } n \ 1 = \text{succ } n = \text{succ } (\text{plus } n \ 0)$
- * Assume we know how to compute $\text{plus } n \ m$
- * Then, $\text{plus } n \ (\text{succ } m)$ is $\text{succ } (\text{plus } n \ m)$

Computation

- * Unravel the definition

- * `plus 7 3`

 - `= plus 7 (succ 2)`

 - `= succ (plus 7 2)`

 - `= succ (plus 7 (succ 1))`

 - `= succ (succ (plus 7 1))`

 - `= succ (succ (plus 7 (succ 0)))`

 - `= succ (succ (succ (plus 7 0)))`

 - `= succ (succ (succ 7))`

Inductive/recursive definitions

- * $\text{plus } n \ 0 = n$, for every n
- * $\text{plus } n \ 1 = \text{succ } n = \text{succ } (\text{plus } n \ 0)$
- * Assume we know how to compute $\text{plus } n \ m$
- * Then, $\text{plus } n \ (\text{succ } m)$ is $\text{succ } (\text{plus } n \ m)$

Recursive definitions ...

Multiplication is repeated addition

- * $\text{mult } n \ m$ means apply $\text{plus } n, m$ times
- * $\text{mult } n \ 0 = 0$, for every n
- * $\text{mult } n \ (\text{succ } m) = \text{plus } n \ (\text{mult } n \ m)$

Summary

- * Functional programs are rules describing how outputs are derived from inputs
- * Basic operation is function composition
- * Recursive definitions allow repeated function composition, depending on the input

Types

Functions work on values of a fixed type

- * `succ` takes a whole number as input and produces a whole number as output
- * `plus` and `mult` take two whole numbers as input and produce a whole number as output
 - * Can also define analogous functions for real numbers

Types

How about `sqrt`, the square root function?

- * Even if the input is a whole number, the output need not be—may have a fractional part
- * Number with fractional values are a different type from whole numbers
 - * In Mathematics, whole numbers are often treated as a subset of fractional or real numbers

Types

Other types

- * `capitalize 'a' = 'A',`
`capitalize 'b' = 'B',...`
- * Inputs and outputs are letters or “characters”

Functions and types

- * We will be careful to ensure that any function we define has a well defined type
 - * The function `plus` that adds two whole numbers will be different from another function `plus` that adds two fractional numbers

Functions have types

- * A function that takes inputs of type A and produces output of type B has a type $A \rightarrow B$
- * In Mathematics, we write $f: S \rightarrow T$ for a function with domain S and codomain T
- * A type is a just a set of permissible values, so this is equivalent to providing the type of f

Collections

- * It is often convenient to deal with collections of values of a given type
 - * A list of integers
 - * A sequence of characters — words or strings
 - * Pairs of numbers
- * Such collections are also types of values

Summary

- * Functions manipulate values
- * Each input and output value comes from a well defined set of possible values — a **type**
- * We will only allow functions whose type can be defined
 - * Functions themselves inherit a type
- * Collections of values also types

Haskell

- * A programming language for describing functions
- * A function description has two parts
 - * Type of inputs and outputs
 - * Rule for computing outputs from inputs
- * Example

`sqr :: Int -> Int` Type definition

`sqr x = x * x` Computation rule

Basic types

- * **Int**, Integers
 - * Operations: **+**, **-**, *****, **/** (Note: **/** produces **Float**)
 - * Functions: **div**, **mod**
- * **Float**, Floating point (“real numbers”)
- * **Char**, Characters, **'a'**, **'%'**, **'7'**, ...
- * **Bool**, Booleans, **True** and **False**

Basic types ...

- * Bool, Booleans, True and False
- * Boolean expressions
 - * Operations: &&, ||, not
 - * Relational operators to compare Int, Float, ...
 - * ==, /=, <, <=, >, >=

Defining functions

- * `xor` (Exclusive or)

- * Input two values of type `Bool`

- * Check that exactly one of them is `True`

```
xor :: Bool -> Bool -> Bool (why?)
```

```
xor b1 b2 = (b1 && (not b2)) ||  
            ((not b1) && b2)
```


Pattern matching

- * Multiple definitions, by cases

```
xor :: Bool -> Bool -> Bool
xor True  False = True
xor False True  = True
xor b1    b2    = False
```

- * Use first definition that matches, top to bottom
 - * `xor False True` matches second definition
 - * `xor True True` matches third definition

Pattern matching ...

- * When does a function call match a definition?
 - * If the argument in the definition is a constant, the value supplied in the function call must be the same constant
 - * If the argument in the definition is a variable, any value supplied in the function call matches, and is substituted for the variable (the “usual” case)

Pattern matching ...

- * Can mix constants and variables in a definition

```
or :: Bool -> Bool -> Bool
or True  b      = True
or b     True   = True
or b1    b2     = False
```

- * `or True False` matches first definition
- * `or False True` matches second definition
- * `or False False` matches third definition

Pattern matching ...

- * Another example

```
and :: Bool -> Bool -> Bool
and True  b = b
and False b = False
```

- * The second argument is used differently in the two definitions
 - * First definition: the value b determines the answer
 - * Second definition: the value b is ignored

Pattern matching ...

- * Another example

```
and :: Bool -> Bool -> Bool
and True  b = b
and False _ = False
```

- * Symbol `_` denotes a “don’t care” argument
 - * Any value matches this pattern
 - * The value is not captured, cannot be reused

Pattern matching ...

```
or :: Bool -> Bool -> Bool
or True _      = True
or _    True   = True
or _    _      = False
```

- * Can have more than one `_` in a definition

Recursive definitions

- * Base case: $f(0)$
- * Inductive step: $f(n)$ defined in terms of smaller values, $f(n-1)$, $f(n-2)$, ..., $f(0)$
- * Example: factorial
 - * $0! = 1$
 - * $n! = n \times (n-1)!$

Recursive definitions ...

- * In Haskell

```
factorial :: Int -> Int
factorial 0 = 1
factorial n = n * (factorial (n-1))
```

- * Note the bracketing in `factorial (n-1)`
 - * `factorial n-1` would be read as `(factorial n) - 1`
- * No guarantee of termination: what is `factorial (-1)`

Conditional definitions

- * Use conditional expressions to selectively enable a definition
- * For instance, “fix” `factorial` for negative inputs

```
factorial :: Int -> Int
factorial 0 = 1
factorial n
  | n < 0 = factorial (-n)
  | n > 0 = n * (factorial (n-1))
```


Conditional definitions ..

```
factorial :: Int -> Int
factorial 0 = 1
factorial n
  | n < 0 = factorial (-n)
  | n > 0 = n * (factorial (n-1))
```

- * Second definition has two parts
 - * Each part is **guarded** by conditional expression
 - * Test guards top to bottom
 - * Note the indentation

Conditional definitions ..

```
factorial :: Int -> Int
factorial 0 = 1
factorial n
  | n < 0 = factorial (-n)
  | n > 0 = n * (factorial (n-1))
```

- * Multiple definitions can have different forms
 - * Pattern matching for `factorial 0`
 - * Conditional definition for `factorial n`

Conditional definitions ...

- * Guards may overlap

```
factorial :: Int -> Int
factorial 0 = 1
factorial n
  | n < 0 = factorial (-n)
  | n > 1 = n * (factorial (n-1))
  | n > 0 = n * (factorial (n-1))
```


Conditional definitions ...

- * Guards may not cover all cases

```
factorial :: Int -> Int
factorial 0 = 1
factorial n
  | n < 0 = factorial (-n)
  | n > 1 = n * (factorial (n-1))
```

- * No match for `factorial 1`

Program error: pattern match failure: factorial 1

Conditional definitions ...

- * Replace the last guard by `otherwise`

```
factorial :: Int -> Int
factorial n
  | n == 0 = 1
  | n > 0  = n * (factorial (n-1))
  | otherwise = factorial (-n)
```

- * “Catch all” condition, always true
- * Ensures that at least one definition matches

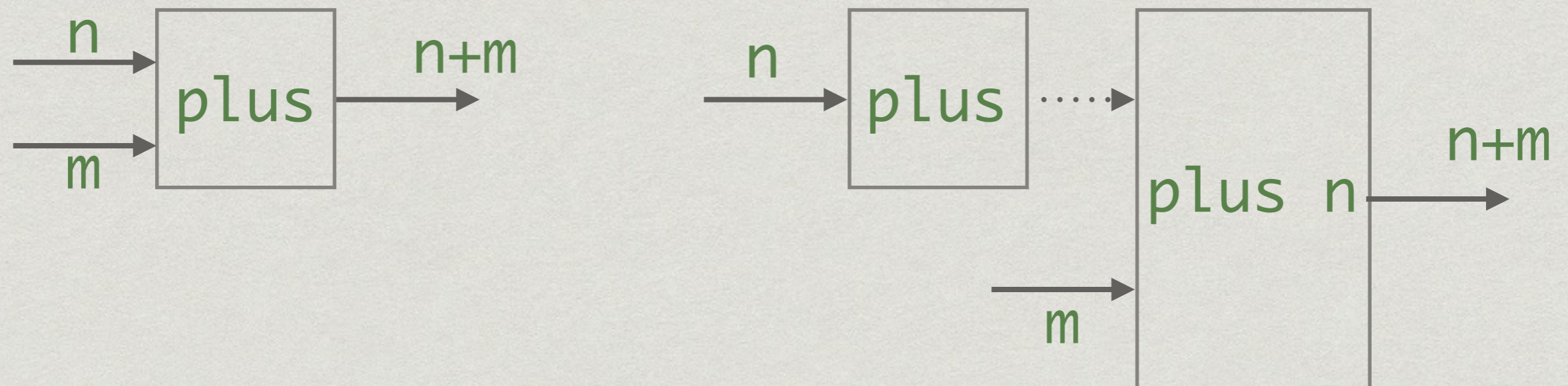
Functions with multiple inputs

Recall that we write `plus n m`, not `plus(n,m)`

- * Normally, functions come with an arity
 - * Number of arguments
- * Instead, assume all functions take only one input!
 - * This is called currying, for the logician Haskell Curry (after whom the language is also named)

Multiple inputs ...

$\text{plus}(n,m) = n + m$ $\text{plus } n \ m = n + m$

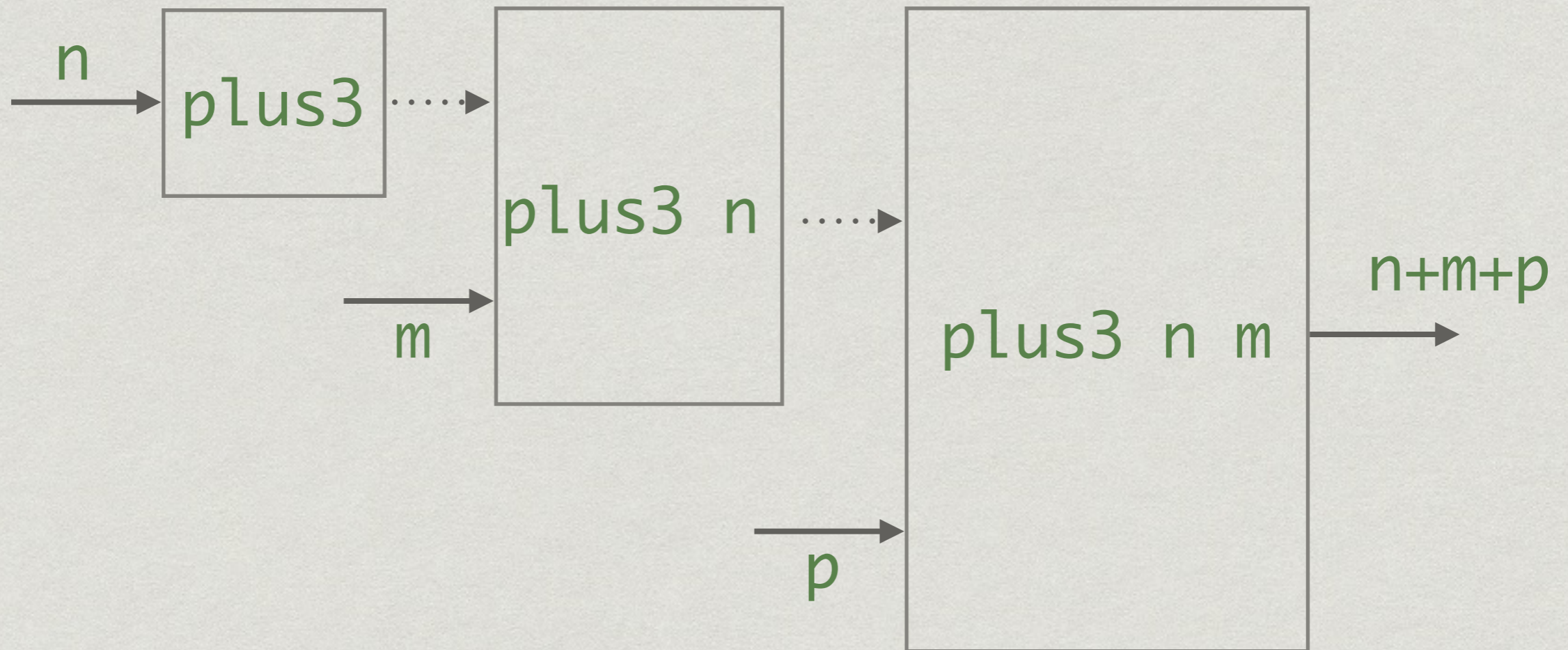


Type of plus

- * plus n: input Int, output Int, so $\text{Int} \rightarrow \text{Int}$
- * plus : input Int, output $\text{Int} \rightarrow \text{Int}$, so $\text{Int} \rightarrow (\text{Int} \rightarrow \text{Int})$

Multiple inputs ...

`plus3 n m p = n + m + p`



`plus3 : Int -> (Int -> (Int -> Int))`

Multiple inputs ...

- * Consider a function with many arguments

$$f\ x1\ x2\ \dots\ xn = y$$

- * Suppose each x_i is of type `Int`, y is of type `Bool`

- * Type of f is

$$f :: Int \rightarrow (Int \rightarrow (\dots (Int \rightarrow Bool)\dots))$$

- * Correspondingly, we should write

$$(\dots((f\ x1)\ x2)\ \dots)\ xn = y$$

Multiple inputs ...

- * Fortunately, Haskell knows this!
- * Implicit bracketing for types is from the right, so

`f :: Int -> Int -> ... -> Int -> Bool`

means

`f :: Int -> (Int -> (... -> (Int -> Bool)...))`

Multiple inputs ...

- * Likewise, function application brackets from left

- * So

$f\ x1\ x2\ \dots\ xn$

means

$(\dots((f\ x1)\ x2)\ \dots)\ xn$

- * Which is why we have to be careful to write factorial $(n-1)$ because, factorial $n-1$ means $(\text{factorial } n) - 1$

Running Haskell programs

- * Haskell interpreter ghci
 - * Interactively call builtin functions
 - * Load user-defined Haskell code from a text file

Setting up ghci

- * Download and install the Haskell Platform
 - * <https://www.haskell.org/platform/>
 - * Available for Windows, Linux, MacOS

Using ghci

- * Create a text file (extension .hs) with your Haskell function definitions
- * Run ghci at the command prompt
- * Load your Haskell code
 - * `:load myfile.hs`
- * Call functions interactively within ghci

Compiling

- * ghc is a compiler that creates a standalone executable from a .hs file
 - * ghc stands for Glasgow Haskell Compiler
 - * ghci is the associated interpreter
- * Using ghc requires some advanced concepts
 - * We will come to this later in the course

Summary

- * ghci is a user-friendly interpreter
 - * Can load and interactively execute user defined functions
- * ghc is a compiler
 - * But we need to know more Haskell before we can use it