

①

1) a)

$$\underbrace{((p \wedge \neg p) \vee a_1)}_{\phi_1} \wedge \underbrace{(a_3 \vee \neg a_2)}_{\phi_2} \wedge \underbrace{(\neg a_1 \vee a_2)}_{\phi_3} \wedge \underbrace{(\neg a_3 \vee (a_4 \wedge \neg a_5))}_{\phi_4}$$

to make $\phi_1 = 1$, $a_1 = 1$ [as $(p \wedge \neg p)$ is not satisfiable]

∴ Hence

$$a_2 = 1 \quad [\text{to make } \phi_3 = 1]$$

$$\neg a_1 \vee a_2 = 0 \vee 1 = 1$$

∴ Hence

$$a_3 = 1 \quad [\text{to make } \phi_2 = 1]$$

∴ Hence

$$(a_4 \wedge \neg a_5) = 1 \quad [\text{to make } \phi_4 = 1]$$

Hence

$$a_4 = 1, a_5 = 0$$

So

$$a_1 = \text{True}$$

$$a_2 = \text{True}$$

$$a_3 = \text{True}$$

$$a_4 = \text{True}$$

$$a_5 = \text{False}$$



(2)

(3) $\Rightarrow \neg((a_1 \wedge \neg a_2) \vee a_3) \Rightarrow (a_3 \vee \neg a_4)$

$\therefore (a_1 \wedge \neg a_2) \vee a_3 = 1$

[to satisfy $\neg(p \Rightarrow q)$
we have only one
assignment
 $p=1, q=0$]

and

$(a_3 \vee \neg a_4) = 0$

[this gives $a_3 = 0$
 $a_4 = 1$]

So,

$a_1 \wedge \neg a_2 = 1$

[as $a_3 = 0$]

$\therefore a_1 = 1$

$a_2 = 0$

So

$a_1 = T$

$a_2 = \perp$

$a_3 = \perp$

$a_4 = T$



(3)

2) a) $(a \Rightarrow b) \vee (a \wedge \neg b)$

~~same as $(\neg a \vee b) \vee (a \wedge \neg b)$~~

to falsify this

we need $a \Rightarrow b = 0$

[so $a=1, b=0$]

and $(a \wedge \neg b) = 0$

[so $a=0, b=1$]

hence

$\begin{matrix} a \\ b \end{matrix} \neq \begin{matrix} \top \\ \perp \end{matrix}$ No assignment ✓

(1)

b) $(\neg b \Rightarrow (a \Rightarrow c)) \vee (a \wedge b)$

to falsify this we need

$(\neg b \Rightarrow (a \Rightarrow c)) = 0$ and $(a \wedge b) = 0$

$\alpha.$ $b \vee (a \Rightarrow c) = 0$

$\alpha.$ $b \vee (\neg a) \vee c = 0$ Both need not be 0.

[so $a=0, b=0$]

[Putting $a=0, b=0$, we ~~get~~ have no value for c to make that 0]

$a=1, b=0, c=0$

hence no assignment.

(4)

3) Not possible

Let $\phi[a, b]$ be an formula
using a, b, \neg, \oplus only

Note there are only 4 possible
assignments say a_1, a_2, a_3, a_4

Claim. $\phi[a, b]$ will always have only
even many satisfiable assignment.

Base case $i) \phi [a, b] = a$

a	b	ϕ
0	1	0
0	0	0
1	0	1
1	1	1

dr. None will satisfy
Basically even no. of satisfying assignments

ii) $\phi [a, b] = \neg a$

a	b	ϕ
0	1	1
0	0	1
1	0	0
1	1	0

iii) $\phi [a, b] = a \oplus b$

follows from the truth table

iv) $\phi [a, b] = a \oplus a$

"

Same for $\phi = b, \neg b, b \oplus b$

v) $\phi (a, b) = a \oplus \neg a$
always true

③

inductive step

assume it is true for

on $\emptyset \neq \phi[a, b]$

with $|\phi| < n$

Now for ϕ , with $|\phi| = n$

We say

i) $\phi = \neg \phi'$, where $|\phi'| = n-1$

The case analysis is confusing. [holds the statement as it is true for $\forall |\phi| \leq n-1$]

It has to be a little more structured.

ii) $\phi = \phi_1 \oplus \phi_2$

~~for $a_1 = 1$~~

let the 4 assignments be a_1, a_2, a_3, a_4

and
if $\phi_1 = 1$ for a_1, a_2
if $\phi_2 = 1$ for a_1, a_2

then $\phi_1 \oplus \phi_2 = 0$ for a_1, a_2
 $\phi_1 \oplus \phi_2 = 1$ for a_3, a_4

if $\phi_2 = 1$ for only one of a_1 or a_2
say for a_1 and for a_3 (say)
 $\phi_1 \oplus \phi_2 = 0$ for a_1, a_4
 $\phi_1 \oplus \phi_2 = 1$ for a_2, a_3

⑥

if $\phi_2 = 1$ for a_3, a_4

then $\phi_1 \oplus \phi_2 = 1$ for any assignment.

Hence our claim is true

Now

if

$$\phi[a, b] = a \wedge b \quad a \Rightarrow b$$

then it can not have
even many satisfaction.



⑧

⇔

$X \rightarrow$ finitely satisfiable

$\Rightarrow X \rightarrow$ satisfiable [proved in the class]

$\Leftarrow \exists v$, s.t. $\forall \phi \in X, v(\phi) = \neg$

Now for α , if $v(\alpha) = \top$, then $X \cup \alpha$ satisfiable

else $v(\neg\alpha) = \top$,

hence $X \cup \{\neg\alpha\}$ satisfiable

\Rightarrow ~~one~~ one of $X \cup \{\alpha\}$ or $X \cup \{\neg\alpha\}$

satisfiable

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(7)

5)

let C represents the chain of all finitely satisfiable ~~set~~ set constructed from this. let $\alpha_1, \alpha_2, \dots$ be the

~~for~~ formulas outside X

now $C_0 = X$

$$C_i = C_{i-1} \cup \{\alpha_i\} \text{ or } C_{i-1} \cup \{\neg \alpha_i\}$$

which one is satisfiable
(from 4)

Say $M = \cup C_i$

claim M is finitely satisfiable

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for any finite $S \subseteq M$

we have C_i s.t. $S \subseteq C_i$

Why is Zorn's Lemma needed?

$C_i \rightarrow$ finitely satisfiable

$\Rightarrow S \rightarrow$ is satisfiable

$\Rightarrow M$ is finitely satisfiable

Hence \exists an maximal chain from

Zorn's Lemma