Sep 27/28, 2022

Example: You are given 3ⁿ coive, all identical except for one which is heavier. Using a balance, prove that you can find the heavy coin in a weighings. Solution If n=1, we have only 3 coins. Pick any two and compare them on the balance. If the two coins have the same weight, then the one left ont to the heariest. Otherwise the bolance telles ve which is henrier. So one weighing is enough. he now know how to write this when n=1. Now lot us work this out for n=2. We have " covins. We can make three piles out of them of three coins early. Rick any two piles and compare them on the balance. If they balance, the pile left out is herriest. If not, the balance tells us which one is heavier. In either care, after one weighing me have a pile of three coins, identical, except for one coin which is hearing. We have already seen (from the n=1 case) that one more overighing and we can identify the heavy coin. How would you do a pile of 3'= 27 coins! The idea is none clear. Split it into 3 piles of 9 coins each. One weighing tells you which pile of 9=32 coins has the heavy coin. The previous case tells you that 2 further weighings and you have identified the heavy coin. In total, three weighings are mongh when n=3. In general, if n 22, and you have worked out how to identify the heavy coin for a pile of 3n-1 coins in n-1 maighings, then for a pile of 3" coins, you should break this pile into 3 piles of 3n-1 coins carl. In me weighing you can write out which pile of 3ⁿ⁻¹ cosins has the heavy coin. And since you have worked not the n-1 care, in another n-1 weighings you are done.

Example : Find an upper bound for the sequence

$$\sqrt{2}$$
, $\sqrt{2 \pm \sqrt{2}}$, $\sqrt{2 \pm \sqrt{2 \pm \sqrt{2}}}$, $\sqrt{2 \pm \sqrt{2 \pm \sqrt{2}}}$, $\sqrt{2 \pm \sqrt{2}}$?
Solution:
Solution:
Since $2 \le 4$, $\sqrt{2} \le \sqrt{4} = 2$.
What about $\sqrt{2 \pm \sqrt{2}}$?
We have $2 \pm \sqrt{2} \le 2 \pm 2 = 4$, and hence
 $\sqrt{2 \pm \sqrt{2}} \le 2 \pm 2 = 4$, and hence
 $\sqrt{2 \pm \sqrt{2}} \le 2 \pm 2 = 4$
by that about $\sqrt{2 \pm \sqrt{2}} < 2$.
We have
 $2 \pm \sqrt{2 \pm \sqrt{2}} \le 2 \pm 2 \le 4$
and hence
 $\sqrt{2 \pm \sqrt{2 \pm \sqrt{2}}} \le 54 = 2$.
Solutionse on this usay, it hashs very likely that 2
18 an upper bound for the sequence. Hence to the basic
idea. Let $8n$ be the other term of the sequence. Then
 $4n \pm \sqrt{2 \pm \sqrt{2}} = 1 \pm 2 \pm 5 \pm 5 \pm 2$.
So once calculations above, we first showed $8_1 \le 2$. Then
uncle the fast that $8_1 \le 2$ to show that $8_2 \le 2$. And then
we used the fast that $8_1 \le 2$ to show that $8_2 \le 2$.
So general, suppre we have shown that $8_{15} \le 2$.
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So general shows that $8_{2} \le 2 + 3 \pm 3 \pm 2$.
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The principle of mathematical in derition
Let Si, Se, So, ... be a sequence of statements (e.g. equation)
such that
(bace case) Si is true
(Inductive step)
$$S_k \Rightarrow S_{k+1}$$
 for all $k \in k$.
Then, Sn is true for all $n \in k$.
The picture of this dominado above should help with
understanding the principle.
Example: Shows that $\sum_{k=1}^{n} i = n(m_k)$ $n \in N$.
Real by induction: Let Sn be the assortion above.
The have case, when $n=1$, is true because
 $\sum_{k=1}^{n} i = 1 = \frac{1(m_k)}{1}$ Si is true
let use three that S_{k+1} is true.
Let use three that S_{k+1} is true.
 $\sum_{k=1}^{n} i = (\sum_{k=1}^{n} i) + (k+1)$ (since S_k is true)
 $= \frac{k(k+1) + 2(k+1)}{2}$
This shows that S_{k+2} is true, is true to N be the state S_k is true)
 $\sum_{k=1}^{n} i = (\sum_{k=1}^{n} i) + (k+1)$ (since S_k is true)
 $\sum_{k=1}^{n} \sum_{k=1}^{n} \sum_{k=1}^{$

Example: Show that $1+3+5+...+(2n-1)=n^2$, ne N. Solution: $S_n: 1+3+5+\dots+(2n-1)=n^2$ Have to show Shis the for all nEIN. Set n=1, The L.S. J.Sn=1 & the R.S. A Sn=1. So Si is true (base case) Suppose Spe is true for some k >1. L.S. of Skal unit Pictoral prof 1+3+5+...+ (2k-1)+ (2k+1) 1+3 +5 +7 = 1+ 3+5+...+ (2k-1) + (2k+1) = k² + (2k+1) (since Sp ie tome). $= (k+l)^2$ The general $= R.S. \int S_{k+1}$ • • case can • • also be proved • Thus by drawing Sk => Sk+1 kew. an nyn senare and • • using the same trick By induction, In is time for all n. // Example: Let n be a positive integer. Show that any 2ⁿx2ⁿ chessboard with one square removed can be tiled noing using L-shaped pieces, where these pieces cover three squares at a time. - L-shaped Tile covering 3 squares. Square removed Write out the solution yourself. The solution was discussed in class. Wint: If you know how to tile a 2"x 2"-' chess board with a square removed, can you use that knowledge to tile a

2ⁿ x 2ⁿ chessboard with a square removed?

Sometimes induction is NOT strong enough.

| Example: Define a, az, az, | by the recursive velations |
|--------------------------------------|----------------------------|
| $\alpha_1 = 3, \alpha_2 = 5$ | $a_1 = 3$ |
| $a_n = 2a_{n-1} - a_{n-2},$ | $n_3 3$. $q_2 = 5$ |
| Let us compute az, az, az. | az = 7 |
| $a_3 = 2a_2 - a_1$ | ay = 9 |
| = 2(5) - 3 | $a_{\rm s} = 11$ |
| = 7 | |
| $a_4 = 2a_3 - a_2$ | Pattern suggest that |
| = 2(7) - 5 | $a_n = 2n + 1$ |
| = 9 | |
| $a_{\rm f} = 2a_{\rm f} - a_{\rm z}$ | |
| = 2(9) - 7 | |
| = 11 | |

What is any? Chances are it is 2(100)+1=201.

an = 2n+1<u>Clarine</u>: will induction work? The statement is certainly true for n=1. Suppose it is true for n=k, ie. suppose we know Q = 2k-1 for some k 2.2. (we know it for k=1 e 2 by inspection.) a = 2ak ~ ak-1 $= 2(2k+1) - a_{k-1}$ What do me do $= 4k + 2 - (a_{k-1})$ with this?

We are stuck. So the principle of induction (in the form we have given it) does not quite help here. What if we assumed that for \$7, the statement is true n=k AND for n=k-17 Let ne do that. We have already seen that $a_{k+1} = 4k+2 - a_{k-1}$ Since the statement is assumed true for n=k-1 also, the abone gines = 42+2-(2(k-1)+1)Q 2-41 = 4k+2 - (2k - 1)= 212+3 = 2(k+1) + 1Which means $S_{\mathbf{k}} + S_{\mathbf{k}-1} \implies S_{\mathbf{k}+1}$ (ج) Snice the statement is true for n=1 & 2, by &) it is true for n=3. None the statement is true for n=2 and n=3. So it must be true for n=4. Continuing this way we see that Sn is true for all nGIN. We dearly need to strongthen our principle of induction in light of the above example. The Strong Brinciple of Mathematical Induction Let S1, S2, ..., be a seguence of statements ench that: (Base care) Si is true and for all kGN (Inductive step) For all & G.N., Spece is true whenever Sis Sz, Sz. ..., Sz are true. Then, Su is true for all ne IN. (usual induction)

 $\phi^{a-1}(1+\phi) - \psi^{a-1}(1+\psi)$ Ξ φ-ψ ψ^{k-1} (ψ²⁾ $\frac{\phi^{(l)-1}(\phi^2)}{\phi-\psi}$ λΨ π-1=D (since a sation Ξ φ = +1 - ψ = +1 Ξ **φ-**Ψ This means Sket is also true. q.e.d.