

\* Purpose of the paper is to extend a modified version of a result of Vande Vate (1989) which characterizes stable matchings as the extreme points of some polytope.

Vande Vate gave an LP formulation for the standard stable matching problem b/w  $n$  men &  $n$  women each of who has a strictly ordered preference list which doesn't leave anybody of the opposite gender out.

In this paper, the author also considers stable matchings that allow staying single (unmatched), so  $|M| = |W|$  is not necessary.

$M$ : men ,  $W$ : women

For each  $m$ ,  $W_m \subseteq W$  is the (strict) preference list of  $m$ .  
For each  $w$ ,  $M_w$  is similarly defined.

For  $w_1, w_2 \in W_m$ ,  $w_1 >_m w_2$  means  $m$  prefers  $w_1$  to  $w_2$   
 $>_w$  is similarly defined on  $M_w$

$\leq_K$ ,  $\leq_K$ ,  $\geq_K$  have the usual meaning for  $K \in M \cup W$

$W_m$  is acceptable to  $m$

$M_w$  " " " " $w$

$A = \{(m, w) \mid m \in M_w \text{ & } w \in W_m\} \leftarrow \text{acceptable pairs}$

A matching from  $M$  to  $W$  is represented by an assignment matrix  $x = \{x_{ij}\}_{i \in M, j \in W}$  whose elements are all integral & we want to stipulate that

$$x_{ij} = \begin{cases} 1 & \text{if } i \text{ & } j \text{ are matched to each other} \\ 0 & \text{otherwise} \end{cases}$$

- this is what we want the matrix  $x$  is characterize.  
So, we define the following constraints :

$$\sum_{j \in W} x_{ij} \leq 1 \quad \text{for all } i \in M \quad \left\{ \begin{array}{l} \text{Each row in } x \text{ can have at} \\ \text{most one '1' value which is to} \\ \text{ensure each man is matched to at} \\ \text{most one woman} \end{array} \right\} \quad \text{①}$$

Similarly to ensure each woman can be matched to at most one man in a matching, we have

$$\sum_{i \in M} x_{ij} \leq 1 \text{ for all } j \in W \quad \text{--- (2)}$$

Note: We are considering at most one mate because one is allowed to remain unmatched here.

$$x_{ij} \geq 0 \quad \forall i \in M \text{ & } j \in W \quad \left\{ \begin{array}{l} \text{Don't allow negative values} \\ \text{--- (3)} \end{array} \right.$$

So, if  $x$  is integer (meaning each value is integral) then clearly,  $x_{ij} \in \{0, 1\}$ .

$w(x, m)$  : mate of  $m$  in  $x$

$m(x, w)$  : mate of  $w$  in  $x$

\* A matching is stable if no one is matched to an unacceptable person & the following holds:

Say  $x$  is the assignment matrix representing the matching

There is no  $(m, w) \in A$  s.t. (the pair is acceptable)

(i) both  $m$  &  $w$  have mates in  $x$  but  $w >_m w(x, m)$  and  $m >_w m(x, w)$  // Both prefer each other to their mates

(ii)  $m$  has a mate,  $w$  doesn't but  $w >_m w(x, m)$   
//  $m$  prefers  $w$  to his current mate

(iii)  $w$  has a mate,  $m$  doesn't but  $m >_w m(x, m)$   
//  $w$  prefers  $m$  to her current mate

(iv) both  $m$  &  $w$  don't have mates.

These conditions for stability can be expressed by the following constraints:

$$x_{ij}=0 \quad \forall (i, j) \in (M \times W) \setminus A \quad \left\{ \begin{array}{l} \text{Unacceptable pairs} \\ \text{can't be matched} \end{array} \right. \quad \text{--- (4)}$$

Conditions (i), (ii), (iii), (iv) can be incorporated by the following constraint:

$$\sum_{\substack{j \in W_m \\ j >_m w}} x_{mj} + \sum_{\substack{i \in M_w \\ i >_w m}} x_{iw} + x_{mw} \geq 1 \quad \text{--- (5)}$$

This inequality ensures that for a given stable matching at least one of the 3 terms must be 1, i.e., either  $x_{mw} = 1$  // m is matched to w OR

$$\sum_{j >_m w} x_{mj} = 1 \quad // \text{m prefers his mate to } w \quad \left. \begin{array}{l} \text{allowing} \\ \text{partial or} \\ \text{fractional} \\ \text{matching} \end{array} \right\}$$

$$\sum_{i >_w m} x_{iw} = 1 \quad // \text{w prefers her mate to } m \quad \left. \begin{array}{l} \text{allowing} \\ \text{partial or} \\ \text{fractional} \\ \text{matching} \end{array} \right\}$$

\* Lemma 1 : Let  $\alpha$  be a matching. Then  $\alpha$  is stable if & only if  $\alpha$  satisfies (1) & (5).

Pf: It is evident that (5) is violated iff for some  $(m, w) \in A$

$$\sum_{j >_m w} x_{mj} = \sum_{i >_w m} x_{iw} = x_{mw} = 0$$

which precisely characterizes one or more of the 4 conditions for violating stability.

\* By lemma 1, it follows that  $\alpha = \{x_{ij}\}_{i \in M, j \in W}$  represents a stable matching iff  $\alpha$  is an integer sol<sup>n</sup> of (1) - (5)

Now we will see some results about the solutions (not necessary integral) of (1) - (5)

Let  $\alpha$  satisfy (1) - (5). Definitions:

$$S_m(\alpha) = \{m \in M \mid \sum_{j \in W} x_{mj} > 0\} \quad // \text{men who are matched (maybe partially) in } \alpha$$

$$S_w(\alpha) = \{w \in W \mid \sum_{i \in M} x_{iw} > 0\} \quad // \text{women who are matched in } \alpha \text{ (maybe partially or fractionally)}$$

$$w^*(\alpha, m), w_* (\alpha, m) \in \{j \in W \mid x_{mj} > 0\}$$

Out of all partners of m in  $\alpha$ ,  $w^*(\alpha, m)$  is the best partner, i.e., most preferred w.r.t.  $\succ_m$  order for  $m \in S_m(\alpha)$ .  $m^*(\alpha, w)$  &  $m_*(\alpha, w)$  are similarly defined for  $w \in S_w(\alpha)$ .

\* Note : ④ assures that  $\{j \in W \mid x_{mj} > 0\} \subseteq W_m$  so  $x_m$  is defined on it.

\* Observe that  $x$  is integer iff  $m \in S_m(x)$ ,

$$\sum_{j \in W} x_{mj} = 1 \quad \& \quad w^*(x, m) = w_*(x, m) \quad \text{i.e.,}$$

For any man who is matched, he is completely matched & has exactly one mate  $w(x, m) = w^*(x, m) = w_*(x, m)$ . We can switch the roles of  $m$  &  $w$  & say the same thing.

Lemma 2 : Let  $x$  satisfy ① - ⑤ &  $(m, w) \notin A$ . Then

$$m \notin S_m(x) \text{ or } m \in S_m(x) \quad \& \quad w \geq_m w^*(x, m)$$

$$\Rightarrow \sum_{i \in M} x_{iw} = 1 \quad \& \quad m \leq_w m_*(x, w) \quad \text{--- ⑥}$$

Consider  $m \notin S_m(x)$ , so  $m$  is not matched at all, then clearly  $w$  must be fully matched otherwise it will violate ⑤, intuitively, we can add more fractional values to the mates of  $w$  in the matching to make it fully matched. So,  $\sum_{i \in M} x_{iw} = 1$

Also, since  $m$  is not matched at all,  $m \neq m_*(x, w)$  &  $m >_w m_*(x, w)$  is also not possible as it violates stability ::  $w$  prefers  $m$  to her worst partner but isn't matched to  $m$ . In case of an integral stable matching, it is easy to see, in case of a fractional matching,  $m >_w m_*(x, w)$  will violate ⑤.

$$\text{So } m \leq_w m_*(x, w)$$

Now suppose  $m \in S_m(x)$  &  $w \geq_m w^*(x, m)$ , i.e.,  $m$  is matched &  $w$  is at least as preferred as the best partner of  $m$  in  $x$ . So  $w$  must be fully matched as otherwise it would violate stability or more formally, ⑤ for similar reasons as previously seen.

Also,  $m \leq_w m_*(x, w)$  because  $w$  is equally or

more preferred to the best mate of  $m$  by  $w$ , so to avoid instability,  $m$  must be at most as liked as the worst partner of  $w$  & not more than that. Hence ⑥ holds.

$$m \in S_m(x) \text{ & } w = w^*(x, m) \iff \sum_{i \in M} x_{iw} = 1 \text{ & } m = m_*(x, w) \quad \text{--- (7)}$$

Pf.  $\Rightarrow$ ) Suppose  $m$  is matched &  $w$  is the best mate of  $m$ . Then by the previous result ⑥, it follows that  $\sum_{i \in M} x_{iw} = 1$ , i.e.,  $w$  is fully matched &

$m \leq_w m_*(x, w)$ , i.e.,  $w$  prefers  $m$  not more than her worst mate.

Suppose  $m <_w m_*(x, w)$ , but  $\therefore w$  is the best mate of  $m$ ,  $m$  is matched to  $w$ , i.e.,  $x_{mw} > 0$  which by definition means that  $m \geq_w m_*(x, w) \Rightarrow \Leftarrow$

Hence,  $m = m_*(x, w)$  //

\* Now we derail a bit & using this result, we prove that  $m$  is matched in  $x$  iff  $m$  is fully matched in  $x$

Subclaim:

$$m \in S_m(x) \iff \sum_{j \in W} x_{mj} = 1 \quad \text{--- fm}$$

Pf of subclaim:  $(\Leftarrow)$  direction is trivial.

$(\Rightarrow)$ : Define a set  $F_w(x) = \{w \in W \mid \sum_{i \in M} x_{iw} = 1\}$  // fully matched women

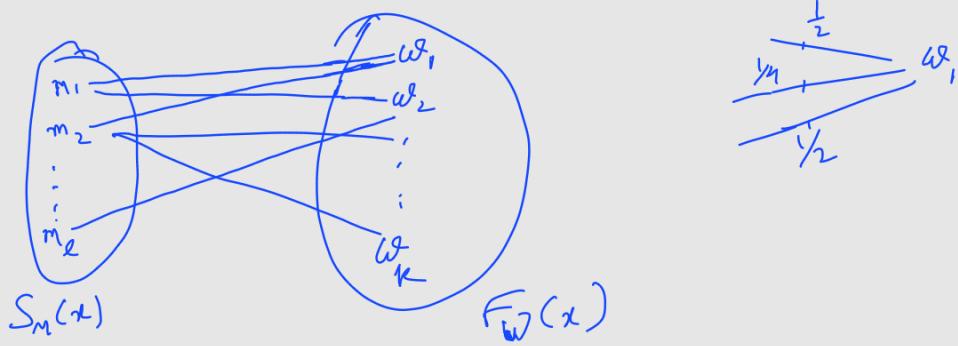
want to prove:  $w^*(x, \cdot) : S_m(x) \rightarrow F_w(x)$  is a bijection.

By the previously proven result, it is clear that if  $m \in S_m(x)$ ,  $w^*(x, m) \in F_w(x)$ . Also observe  $w^*(x, \cdot)$  is injective because suppose  $m_1, m_2 \in S_m(x)$  s.t.  $w^*(x, m_1) = w^*(x, m_2) = w$

By previous result, we know  $m_1 = m_*(x, w)$  &  $m_2 = m_*(x, w)$

$\Rightarrow m_1 = m_2 = m_*(x, w)$   $\because$  the worst partner of  $w$  is unique. Hence,  $|F_w(x)| \geq |S_m(x)|$

Consider the corresponding bipartite graph of the matching  $\chi$  (not necessarily integral) :



We know that for each  $w \in F_w(x)$ ,  $\sum_{i \in M} x_{iw} = 1$

Hence in this graph restricted to  $S_m(x) \times F_w(x)$ , the sum of all values of edges =  $|F_w(x)|$ . This gives 2 conclusions : i)  $|F_w(x)| = |S_m(x)|$  because, since the edges are incident on elements in  $S_m(x)$ , we can take the sum of all edges from the left side & it should give  $k = |F_w(x)|$  but if  $|S_m(x)| < k$ , then we know that edges outgoing from a single  $m \in S_m(x)$  can add up to at most 1, sum over all  $m \in S_m(x)$  can add up to at most  $|S_m(x)| < |F_w(x)| = k \Rightarrow \infty$ .

ii)  $\sum_{j \in W} x_{mj} = 1$ , i.e., each  $m \in S_m(x)$  must be fully matched since their edges must sum up to  $k = |S_m(x)| = |F_w(x)|$  & edges for each man can sum up to at most 1. This establishes  $\text{f}_m$

$\therefore w^*(x, \cdot)$  is injective &  $|S_m(x)| = |F_w(x)|$ ,

$w^*(x, \cdot)$  is bijective.

Now, we are ready to prove  $(\Leftarrow)$  of ⑦.

Suppose  $w$  is fully matched, i.e.,  $\sum_{i \in M} x_{iw} = 1$  &  $m = m_*(x, w)$

Then  $w \notin F_w(x)$  but  $\because w^*(x, \cdot)$  is onto,  $\exists m' \in S_m(x)$  s.t.

$w = w^*(x, m')$ , i.e.,  $w$  is the best partner of some  $m' \in S_m(x)$

Then, we can apply  $(\Rightarrow)$  of ⑦ to say  $m' = m_*(x, w) = m$

Hence,  $m' = m \in S_m(x)$  &  $m = m_*(x, w)$  holds.  $\square$

So, ⑦ holds.

Consider

$$\sum_{j \geq_m w} x_{mj} + \underbrace{\sum_{i \geq_w m} x_{iw} + x_{mw}}_{=1} = 1 \quad \text{--- (8)}$$

If (8) holds, then converse of (6) holds as well.

It is easy to see that if  $w$  is fully matched &  $m \leq_w m^*(x, w)$ , i.e.,  $m$  is at most as liked as the worst mate of  $w$  in  $x$ , then  $\sum_{i \geq_w m} x_{iw} = 1$

$\Rightarrow \sum_{j \geq_m w} x_{mj} = 0$  which clearly implies that either

$m$  is unmatched or if  $m$  is matched, his best mate is at most as good as  $w$ , i.e.,  $w^*(x, m) \leq_m w$ .  $\square$

Lemma 3 is just lemma 2 with roles of  $m$  &  $w$  reversed, it is true by symmetric args.

Lemma 3: let  $x$  satisfy (1)-(5) &  $(m, w) \in A$ . Then

$w \notin S_w(x)$  or  $w \in S_w(x) \wedge m \geq_w m^*(x, w)$

$$\Rightarrow \sum_{j \in W} x_{mj} = 1 \wedge w \leq_m w^*(x, m) \quad \text{--- (9)}$$

&

$$w \in S_w(x) \wedge m = m^*(x, w) \Leftrightarrow \sum_{j \in W} x_{mj} = 1 \wedge w = w^*(x, m) \quad \text{--- (10)}$$

$w \in S_w(x)$  iff  $\sum_{i \in M} x_{iw} = 1$  — (F<sub>w</sub>)

If (8) holds then converse of (9) holds.

Now we are ready to prove the main result.

Theorem 1: let  $C$  be the solution of (1)-(5). Then the integer points in  $C$  are precisely its extreme points.

Pf: ( $\Rightarrow$ ) Easy, suppose  $x$  is an integer sol' of (1)-(5). Then  $x$  is an integer solution of (1)-(3) which are constraints for the general bipartite matching  $\Delta$  for that we have already seen in class that integer points satisfying (1)-(3) are precisely the extreme points of the corresponding polytope. (Also can be derived by Birkhoff's theorem)

Hence,  $x$  is an extreme point of  $\mathcal{D} - \mathcal{S}$  &  $\therefore$  an extreme point of  $\mathcal{D} - \mathcal{G}$ . //

( $\Leftarrow$ ): The tricky part. Suppose  $x$  is an extreme point of  $C$ . We define 3 matrices of order  $(M \times W)$ .

For  $(m, w) \in (M, W)$ ,

$$(\bar{z}^*)_{mw} := \begin{cases} 1, & \text{if } m \text{ is matched in } x \text{ & } w \text{ is his best mate} \\ & \quad [m \in S_m(x) \text{ & } w = w^*(x, m)] \\ 0, & \text{otherwise} \end{cases} \quad \text{--- (11)}$$

$$(\bar{z}_*)_{mw} := \begin{cases} 1, & \text{if } m \text{ is matched in } x \text{ & } w \text{ is his worst mate} \\ & \quad [m \in S_m(x) \text{ & } w = w_*(x, m)] \\ 0, & \text{otherwise} \end{cases} \quad \text{--- (12)}$$

$$(\bar{z})_{mw} = (\bar{z}^*)_{mw} - (\bar{z}_*)_{mw} \quad \text{--- (13)}$$

If we show that  $\bar{z} = 0$ , then that implies

$\bar{z}^* = \bar{z}_*$  which means,  $\forall m \in S_m(x)$ , i.e., for all men who are matched, their best & worst partners are the same i.e.,  $w^*(x, m) = w_*(x, m)$  hence  $m$  has exactly one partner & by  $\textcircled{f}_m$ , we know  $m$  is fully matched so by  $\textcircled{*}$ ,  $x$  is integer. So, it is enough to prove  $\bar{z} = 0$ .

Also note, by using Lemma 2 & 3, we can equivalently define  $\bar{z}^*$  &  $\bar{z}_*$  as:

$$(\bar{z}^*)_{mw} = \begin{cases} 1 & \text{if } w \in S_w(x) \text{ & } m = m^*(x, w) \\ 0, & \text{o/w} \end{cases} \quad \text{--- (14)}$$

$$(\bar{z}_*)_{mw} = \begin{cases} 1 & \text{if } w \in S_w(x) \text{ & } m = m_*(x, w) \\ 0, & \text{o/w} \end{cases} \quad \text{--- (15)}$$

Observe how  $\bar{z}^*$  &  $\bar{z}_*$  are formed:

In  $x$ , For all the men who are matched (maybe fractionally)

$z^*$  matches them to only their best partners, fully.  
 Similarly,  $\bar{z}_*$  matches them to only their worst partners fully. Similarly from the side of matched women but "best" & "worst" gets interchanged.

We prove the following :

$$\sum_{j \in W} z_{ij} = 1 \Rightarrow \sum_{j \in W} \bar{z}_{ij} = 0, i \in M \quad \text{--- (16)}$$

$$\sum_{i \in M} z_{ij} = 1 \Rightarrow \sum_{i \in M} \bar{z}_{ij} = 0, j \in W \quad \text{--- (17)}$$

$$z_{ij} = 0 \Rightarrow \bar{z}_{ij} = 0, i \in M, j \in W \quad \text{--- (18)}$$

$$\bar{z}_{ij} = 0, (i, j) \in (M \times W) \setminus A \quad \text{--- (19)}$$

For  $(m, w) \in A$

$$\underbrace{\sum_{j >_m w} z_{mj} + \sum_{i >_w m} z_{iw} + z_{mw}}_{(21)} = 1 \Rightarrow \underbrace{\sum_{j >_m w} \bar{z}_{mj} + \sum_{i >_w m} \bar{z}_{iw} + \bar{z}_{mw}}_{(22)} = 0$$

If (16) - (20) hold, then for sufficiently small  $\epsilon$ ,  
 $x - \epsilon z$  &  $x + \epsilon z$  satisfy (1) - (5) & hence  $\epsilon \in C$

Then  $x = \frac{1}{2}(x + \epsilon z) + \frac{1}{2}(x - \epsilon z)$  but  $\therefore x$  is extreme,  
 $z = 0$  & we are done!

It remains to prove (16) - (20)

(16) is easy to see because if some  $m$  is fully matched in  $x$ , then it is fully matched to one partner in  $z^*$  & in  $\bar{z}_*$  & the difference of their sums would be zero as  $|1 - 1|$  will cancel each other out.

(17) can be argued similarly but from the POV of women.

In (18), if  $z_{ij} = 0$  then either  $i$  is completely unmatched in which case  $i$  is completely unmatched in  $z^*$  &  $\bar{z}_*$  also & if  $i$  is matched, then  $j$  can't be his best or worst partner so  $(\bar{z}^*)_{ij} = 0 = (\bar{z}_*)_{ij}$

(19) immediately follows from (4) & (18)

Only (20) remains to be shown. Assume (21), claim: (22)

3 cases, Case 1:  $m \notin S_m(x)$  or  $m \in S_m(x) \wedge w \geq_m w^*(x, m)$

Then (11) & (12) implies

$$\sum_{j \geq_m w} (\tilde{g}^*)_{mj} = \sum_{j \geq_m w} (\tilde{g}^*)_{mj} = 0 \quad \text{--- (23)}$$

has 1 exactly at  $j = w^*(x, m)$   
 but  $w^*(x, m) \not\geq_m w$ .

Also, by (6),  $w \in S_w(x) \wedge m \leq_w m^*(x, w) \leq_w m^*(x, w)$

So, by (9) & (5),

$$\sum_{i \geq_w m} (\tilde{g}^*)_{iw} + (\tilde{g}^*)_{mw} = \sum_{i \geq_w m} (\tilde{g}^*)_{iw} + (\tilde{g}^*)_{mw} = 1 \quad \text{--- (24)}$$

1 if  $m <_w m^*(x, w)$   
 0 if  $m = m^*(x, w)$

0 if  $m <_w m^*(x, w)$   
 1 if  $m = m^*(x, w)$

similarly for  $m \leq_w m^*(x, w)$  cases.

(23) & (24)  $\Rightarrow$  (22)

Case 2:  $w \notin S_w(x)$  or  $w \in S_w(x) \wedge m \geq_w m^*(x, w)$

Symmetric arguments as the previous case.

Case 3:  $m \in S_m(x)$ ,  $w \in S_w(x)$ ,  $w \leq_m w^*(x, m) \wedge m <_w m^*(x, w)$

$\therefore m \in S_m(x) \wedge w \leq_m w^*(x, m)$ , it follows that

$$\sum_{j \geq_m w} (\tilde{g}^*)_{mj} = 1 \quad \text{as } j = w^*(x, m) \text{ will yield } (\tilde{g}^*)_{mj} = 1, \text{ others 0.}$$

$$\therefore (\tilde{g}^*)_{mw} = 0$$

We want to prove  $\sum_{i \geq_w m} (\tilde{g}^*)_{iw} = 0$

It is enough to prove that  $m^*(x, w) \leq_w m$

$\therefore (21) = (8)$  holds, Converse of (6) holds. If

$m^*(x, w) \geq_w m$ , then  $w \geq_m w^*(x, m)$  ( $\because m \in S_m(x)$  holds)

$\Rightarrow \infty$ , hence  $m^*(x, w) <_w m$  & we are done.

So, we have  $\sum_{j \geq_m w} (\bar{g}^*)_{mj} + \sum_{i >_w m} (\bar{g}^*)_{iw} + (\bar{g}^*)_{mw} = 1$  — (26)

By symmetric arguments, we can show

$$\underbrace{\sum_{j \geq_m w} (\bar{g}^*)_{mj}}_0 + \underbrace{\sum_{i >_w m} (\bar{g}^*)_{iw}}_1 + \underbrace{(\bar{g}^*)_{mw}}_0 = 1 \quad — (27)$$

Clearly, (26) & (27)  $\Rightarrow$  (22) in this final case.  $\square$ .