## Hints/answers and marking scheme to Assignment 1

## 14/09/2012

- 1. Both  $O(\log n)$ . Correct answer for sorted, reverse sorted: 1.5 marks each.
- 2. DFS + listing of tree edges (as they are obvious from the DFS): 3 marks Other types of edges: 1 mark each. (DFS tree has no cross edges).
- 3. Prim's algorithm with a as root node adds edges in the order (a, b), (b, c), (c, e), (d, a), (d, g), (e, f). Cost: 25. Correct answer: 5 marks.
- 4. Let T,T' be two MSTs in G. Consider  $T \oplus T'$  and choose the smallest weight edge e from it. Let  $e \in T$  but  $e \notin T$ . Then add e to T'. The resulting cycle has another edge  $f \in T'$ ,  $f \notin T$ , w(f) > w(e) by the above assumption. Deleting f from T' and adding e reduces its weight, contradicting that T' is an MST. Correct proof: 6 marks
- 5. MST remains MST as relative ordering among edges remains the same A shortest path may not remain a shortest path. For example, if there are two paths from u to v of weights 3+11=14 and 5+10=15, they become 9+121=130 and 25+100=125 after squaring. Thus the relative order among paths does not remain the same.
  - Correct answers: 1 mark each, Correct explanation/counter example: 1 mark each
- 6. Do a topological sort in linear time. Go over the vertices in the topological order, maintain an array of counts count[i] for number of s to i paths. For each vertex i, count[i] = sum of counts of its in-neighbours. Due to topological ordering, their counts are available while processing i.

Correct algorithm: 6 marks