

PLC : End-semester examination

April 28, 2020. 2.00 pm – 5.00 pm

Marks: 100

Weightage: 30

General instructions:

1. Submit your solutions as a PDF scan of your handwritten answers or as a plain text file on Moodle. The file should be named `<un>-endsem.txt` (or `<un>-endsem.pdf`, if you are submitting a PDF), where `un` is your username. For example, I would submit a file named `spsuresh-endsem.txt`.
2. If you are submitting a text file, use the Haskell notation `\x -> M` for the $\lambda x.M$ when writing expressions in the lambda calculus. (For $\lambda x y z.M$, write `\x y z -> M`.) Use the notation `:=` for syntactic equality, and `-->` and `=` for (many-step) beta-reduction and beta-equality, respectively. Use `M[x <- N]` for substitution. Also use `<m>` for the Church encoding of m , and `x^{y}` for x^y .
3. Properly parenthesize your lambda expressions and use spacing to keep it readable.
4. Recall that the Church encoding of n , denoted `[n]`, is the expression $\lambda f x.f^n x$, where $f^0 x$ is defined to be just x , and $f^{i+1} x := f(f^i x)$.
5. You can assume the following standard encodings introduced in the lecture slides, with the appropriate behaviour:

$$\begin{aligned} [n] &= \lambda f x.f^n x \\ \mathbf{succ} &= \lambda p f x.f(p f x) \\ \mathbf{plus} &= \lambda p q f x.p f(q f x) \\ \mathbf{mult} &= \lambda p q f.p(q f) \\ \mathbf{true} &= (\lambda x y.x) \\ \mathbf{false} &= (\lambda x y.y) \\ \mathbf{pair} &= (\lambda x y w.w x y) \\ \mathbf{fst} &= (\lambda p.p \mathbf{true}) \\ \mathbf{snd} &= (\lambda p.p \mathbf{false}) \\ \mathbf{ite} &= (\lambda b x y.b x y) \\ \mathbf{iszero} &= (\lambda x.(x(\lambda z.\mathbf{false})) \mathbf{true}) \end{aligned}$$

1. Reduce the following terms to β -normal form. (Note the need to rename bound variables appropriately.)

- (a) $(\lambda x y. x y)(\lambda v. v u)(\lambda x. x)$
- (b) $(\lambda x y. x y)((\lambda v. v u)(\lambda x. x))$
- (c) $(\lambda x. x(x(yz))x)(\lambda u. uv)$
- (d) $(\lambda x. xxy)(\lambda y. yz)$
- (e) $(\lambda x y. xyy)(\lambda u. uyx)$

(15 marks)

2. Find a lambda-calculus encoding for the function $S(n)$ defined by $S(n) = 0 + 1 + \dots + n$ for $n \geq 0$. Prove that your encoding is correct. Do not use any fixed-point combinator.

(20 marks)

3. Consider the following encoding of the empty list and **cons** operator.

$$[] = \lambda x y. y$$

$$H : T = \lambda x y. xHT$$

Find λ -expressions with the following behaviour:

- (a) **null**: if $M = []$ then $\mathbf{null} M \xrightarrow{*}_{\beta} \mathbf{true}$, and if M is of the form $H : T$ then $\mathbf{null} M \xrightarrow{*}_{\beta} \mathbf{false}$.
- (b) **head**, which extracts the head of a list and is defined only on inputs of the form $H : T$.
- (c) **tail**, which extracts the tail of a list and is defined only on inputs of the form $H : T$.

(20 marks)

4. Give an example of a *closed* λ -expression (an expression without free variables) that belongs to each of the following types:

- (a) $a \rightarrow (a \rightarrow b) \rightarrow b$
- (b) $(a \rightarrow b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c)$
- (c) $(a \rightarrow b \rightarrow c) \rightarrow (b \rightarrow a \rightarrow c)$
- (d) $((b \rightarrow c) \rightarrow a \rightarrow b) \rightarrow (b \rightarrow c) \rightarrow a \rightarrow c$
- (e) $a \rightarrow ((b \rightarrow a) \rightarrow c) \rightarrow c$

(25 marks)

5. Consider the following two threads running concurrently, with x and y being shared variables, with initial values 7 and 3 respectively. Assume that each assignment statement is atomic.

Thread A

$x := y + 1$

$x := 2*x*y$

$x := x - y$

Thread B

$y := x + 1$

$y := 2*y*y$

$y := y - x$

Show how to generate the following combinations of (x,y) values at the end of the execution of both threads.

(a) $(-10, 60)$

(b) $(30, 10)$

(c) $(21, 947)$

(d) $(4480, -2176)$

Argue that at the completion of both threads, the value of $x+y$ is always even. (20 marks)