NPTEL MOOC PROGRAMMING, DATA STRUCTURES AND ALGORITHMS IN PYTHON

Week 6, Lecture 4

Madhavan Mukund, Chennai Mathematical Institute http://www.cmi.ac.in/~madhavan

Data structures

- * Algorithms + Data Structures = Programs Niklaus Wirth
- * Arrays/lists sequences of values
- Dictionaries key-value pairs
- * Python also has sets as a built in datatype

Sets in Python

List with braces, duplicates automatically removed

colours = {'red', 'black', 'red', 'green'}

>>> print(colours)
{'black', 'red', 'green'}

* Create an empty set

colours = set()

* Note, not colours = {} - empty dictionary!

Sets in Python

* Set membership

```
>>> 'black' in colours
True
```

* Convert a list into a set

```
>>> numbers = set([0,1,3,2,1,4])
>>> print(numbers)
{0, 1, 2, 3, 4}
```

```
>>> letters = set('banana')
>>> print(letters)
{'a', 'n', 'b'}
```

```
Set operations
```

```
odd = set([1,3,5,7,9,11])
prime = set([2,3,5,7,11])
```

- * Union
 odd | prime → {1, 2, 3, 5, 7, 9, 11}
- * Intersection odd & prime \rightarrow {3, 11, 5, 7}
- * Set difference odd - prime \rightarrow {1, 9}
- * Exclusive or odd \wedge prime \rightarrow {1, 2, 9}

Stacks

- * Stack is a last-in, first-out list
 - * push(s,x) add x to stack s
 - * pop(s) return most recently added element
- * Maintain stack as list, push and pop from the right
 - * push(s,x) is s.append(x)
 - * s.pop() Python built-in, returns last element

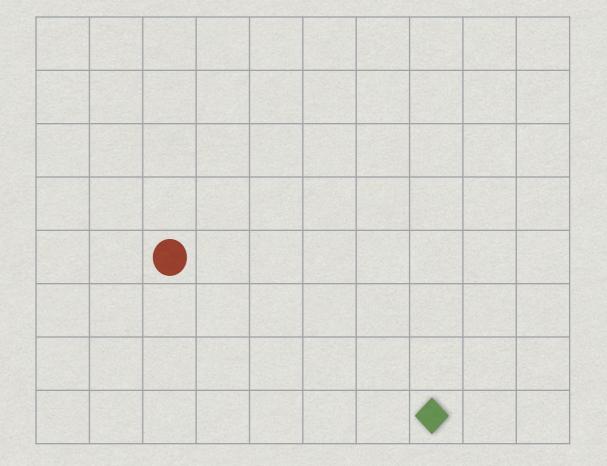
Stacks

- Stacks are natural to keep track of recursive function calls
- In 8 queens, use a stack to keep track of queens added
 - * Push the latest queen onto the stack
 - * To backtrack, pop the last queen added

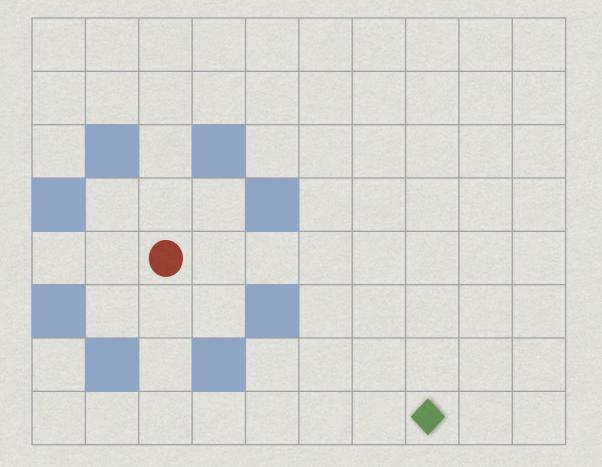
Queues

- * First-in, first-out sequences
 - * addq(q,x) adds x to rear of queue q
 - * removeq(q) removes element at head of q
- * Using Python lists, left is rear, right is front
 - * addq(q,x) is q.insert(0,x)
 - * l.insert(j,x), insert x before position j
 - * removeq(q) is q.pop()

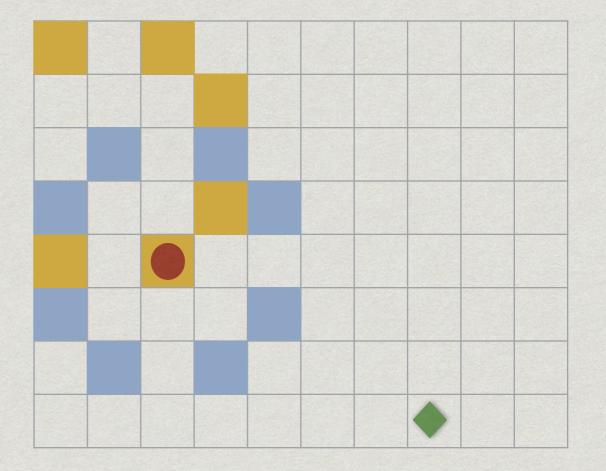
- Rectangular m x n grid
- Chess knight starts at (sx,sy)
 - * Usual knight moves
- * Can it reach a target square (tx,ty)?



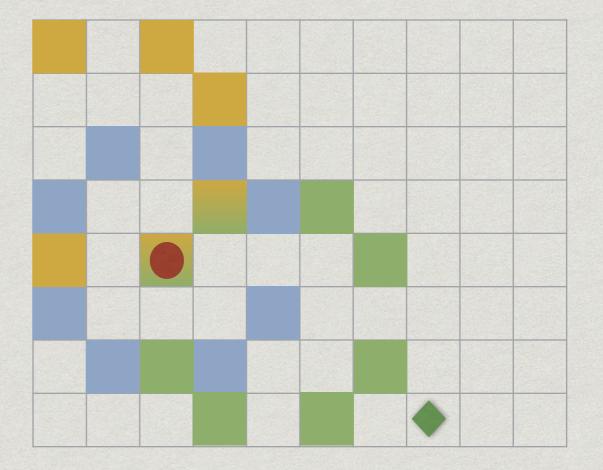
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- * X1 all squares reachable in one move from (sx,sy)
- * X2 all squares reachable from X1 in one move

- * Don't explore an already marked square
- * When do we stop?

*

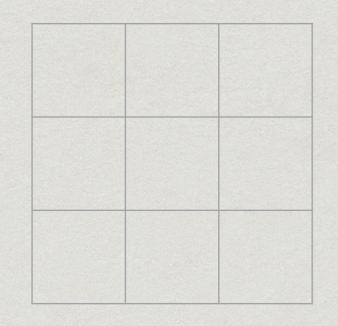
- * If we reach target square
- * What if target is not reachable?

- * Maintain a queue Q of cells to be explored
- Initially Q contains only start node (sx,sy)
 - Remove (ax,ay) from head of queue
 - Mark all squares reachable in one step from (ax,ay)
 - * Add all newly marked squares to the queue
- * When the queue is empty, we have finished

```
def explore((sx,sy),(tx,ty)):
 marked = [[0 for i in range(n)]
                      for j in range(m)]
 marked[sx][sy] = 1
  queue = [(sx, sy)]
 while queue != []:
    (ax,ay) = queue.pop()
    for (nx,ny) in neighbours((ax,ay)):
      if !marked[nx][ny]:
        marked[nx][ny] = 1
        queue.insert(0,(nx,ny))
  return(marked[tx][ty])
```

Example

src = (0,1) tgt =(1,1)



* This is an example of breadth first search

Summary

- Data structures are ways of organising information that allow efficient processing in certain contexts
- * Python has a built-in implementation of sets
- Stacks are useful to keep track of recursive computations
- * Queues are useful for breadth-first exploration