NPTEL MOOC PROGRAMMING, DATA STRUCTURES AND ALGORITHMS IN PYTHON

Week 1, Lecture 3

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Algorithm for gcd(m,n)

- To find the largest common factor, start at the end and work backwards
- * Let i run from min(m,n) to 1
- * First common factor that we find will be gcd!

- * Suppose d divides both m and n, and m > n
- * Then m = ad, n = bd
- * Som-n = ad bd = (a-b)d
- * d divides m-n as well!
- * So gcd(m,n) = gcd(n,m-n)

- * Consider gcd(m,n) with m > n
- * If n divides m, return n
- * Otherwise, compute gcd(n,m-n) and return that value

```
Euclid's algorithm
```

```
def gcd(m,n):
 # Assume m \ge n
 if m < n:
   (m,n) = (n,m)
 if (m\%n) == 0:
   return(n)
 else:
   diff = m-n
   # diff > n? Possible!
   return(gcd(max(n,diff),min(n,diff))
```

Euclid's algorithm, again

```
def gcd(m,n):
```

```
if m < n: # Assume m >= n
(m,n) = (n,m)
```

```
while (m%n) != 0:
diff = m-n
# diff > n? Possible!
(m,n) = (max(n,diff),min(n,diff))
```

return(n)

Even better

- Suppose n does not divide m
- * Then m = qn + r, where q is the quotient, r is the remainder when we divide m by n
- * Assume d divides both m and n
- * Then m = ad, n = bd
- * So ad = q(bd) + r
- It follows that r = cd, so d divides r as well

- * Consider gcd(m,n) with m > n
- * If n divides m, return n
- * Otherwise, let r = m%n
- * Return gcd(n,r)

def gcd(m,n):

if m < n: # Assume m >= n
(m,n) = (n,m)

if (m%n) == 0:
 return(n)

else:

return(gcd(n,m%n)) # m%n < n, always!</pre>

Euclid's algorithm, revisited

def gcd(m,n):

if m < n: # Assume m >= n
(m,n) = (n,m)

while (m%n) != 0:
(m,n) = (n,m%n) # m%n < n, always!</pre>

return(n)

Efficiency

- Can show that the second version of Euclid's algorithm takes time proportional to the number of digits in m
- If m is 1 billion (10⁹), the naive algorithm takes billions of steps, but this algorithm takes tens of steps