NPTEL MOOC, JAN-FEB 2015 Week 3, Module 5

DESIGN AND ANALYSIS OF ALGORITHMS

Applications of BFS and DFS

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Graphs, formally

- G = (V,E)
- Set of vertices V
- Set of edges E
 - * E is a subset of pairs (v,v'): E ⊆ V × V
 - * Undirected graph: (v,v') and (v',v) are the same edge
 - * Directed graph:
 - * (v,v') is an edge from v to v'
 - * Does not guarantee that (v',v) is also an edge

Exploring graph structure

- * Breadth first search
 - * Level by level exploration
- * Depth first search
 - * Explore each vertex as soon as it is visited
 - * DFS numbering
- * What can we find out about a graph using BFS/ DFS?

Connectivity





Connected graph

Disconnected graph

Connectivity





Connected graph

Disconnected graph Connected components

Identifying connected components

- * Vertices {1,2,...,N}
- * Start BFS or DFS from 1
 - * All nodes marked Visited form a connected component
 - * Pick first unvisited node, say j, and run BFS or DFS from j
 - * Repeat till all nodes are visited
- Update BFS/DFS to label each visited node with component number



- * Add a counter comp to number components
- Increment counter each time a fresh BFS/DFS starts
- * Label each visited node j with component[j] = comp



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Acyclic graph

Graph with cycles



- * Edges explored by BFS form a tree
 - * Acyclic graph = connected, with n-1 edges



- * Edges explored by BFS form a tree
 - * Acyclic graph = connected, with n-1 edges
- * Any non-tree edge generates a cycle









pre



pre

2

post




















































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- A directed graph has a cycle if and only if DFS reveals a back edge
- * Can classify edges using pre and post numbers
 - * Tree/Forward edge (u,v) : Interval [pre(u),post(u)] contains [(pre(v),post(v)]
 - Backward edge (u,v):
 Interval [pre(v),post(v)] contains [(pre(u),post(u)]
 - * Cross edge (u,v): Intervals [(pre(u),post(u)] and [(pre(v),post(v)] disjoint

Directed acyclic graphs

- Directed graphs without cycles are useful for modelling dependencies
 - * Courses with prerequisites
 - Edge (Algebra, Calculus) indicates that Algebra is a prerequisite for Calculus
- * Will look at Directed Acyclic Graphs (DAGs) soon

Connectivity in directed graphs

- Need to take directions into account
- Nodes i and j are strongly connected if there is a path from i to j and a path from j to i
- Directed graph can be decomposed into strongly connected components (SCCs)
 - All pairs of nodes in an SCC are strongly connected

Computing SCCs



 DFS numbering (pre and post) can be used to compute SCCs

[Dasgupta, Papadimitriou,Vazirani]


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Other properties

- A number of other structural properties can be inferred from DFS numbering
- Articulation points (vertices)
 - * Removing such a vertex disconnects the graph
- * Bridges (edges)
 - * Removing such an edge disconnects the graph