

NPTEL MOOC, JAN-FEB 2015
Week 1, Module 5

DESIGN AND ANALYSIS OF ALGORITHMS

MADHAVAN MUKUND, CHENNAI MATHEMATICAL INSTITUTE
<http://www.cmi.ac.in/~madhavan>

Analysis of algorithms

- * Measuring efficiency of an algorithm
 - * Time: How long the algorithm takes (running time)
 - * Space: Memory requirement

Time and space

- * Time depends on processing speed
 - * Impossible to change for given hardware
- * Space is a function of available memory
 - * Easier to reconfigure, augment
- * Typically, we will focus on time, not space

Measuring running time

- * Analysis independent of underlying hardware
 - * Don't use actual time
 - * Measure in terms of “basic operations”
- * Typical basic operations
 - * Compare two values
 - * Assign a value to a variable
- * Other operations may be basic, depending on context
 - * Exchange values of a pair of variables

Input size

- * Running time depends on input size
 - * Larger arrays will take longer to sort
- * Measure time efficiency as function of input size
 - * Input size n
 - * Running time $t(n)$
- * Different inputs of size n may each take a different amount of time
- * Typically $t(n)$ is **worst case estimate**

Example 1: Sorting

- * Sorting an array with n elements
 - * Naïve algorithms : time proportional to n^2
 - * Best algorithms : time proportional to $n \log n$
- * How important is this distinction?
- * Typical CPUs process up to 10^8 operations per second
 - * Useful for approximate calculations

Example 1: Sorting ...

- * Telephone directory for mobile phone users in India
 - * India has about 1 billion = 10^9 phones
- * Naïve n^2 algorithm requires 10^{18} operations
 - * 10^8 operations per second $\Rightarrow 10^{10}$ seconds
 - * 2778000 hours
 - * 115700 days
 - * 300 years!
- * Smart $n \log n$ algorithm takes only about 3×10^{10} operations
 - * About 300 seconds, or 5 minutes

Example 2: Video game

- * Several objects on screen
- * Basic step: find closest pair of objects
- * Given n objects, naïve algorithm is again n^2
 - * For each pair of objects, compute their distance
 - * Report minimum distance over all such pairs
- * There is a clever algorithm that takes time $n \log n$

Example 2: Video game ...

- * High resolution monitor has 2500 x 1500 pixels
 - * 3.75 million points
- * Suppose we have 500,000 = 5×10^5 objects
- * Naïve algorithm takes 25×10^{10} steps = 2500 seconds
 - * 2500 seconds = 42 minutes response time is unacceptable!
- * Smart $n \log n$ algorithm takes a fraction of a second

Orders of magnitude

- * When comparing $t(n)$ across problems, focus on orders of magnitude
 - * Ignore constants
- * $f(n) = n^3$ eventually grows faster than $g(n) = 5000 n^2$
 - * For small values of n , $f(n)$ is smaller than $g(n)$
 - * At $n = 5000$, $f(n)$ overtakes $g(n)$
 - * What happens in the limit, as n increases :
asymptotic complexity

Typical functions

- * We are interested in orders of magnitude
- * Is $t(n)$ proportional to $\log n$, ..., n^2 , n^3 , ..., 2^n ?
- * Logarithmic, polynomial, exponential ...

Typical functions $t(n)$...

| Input | $\log n$ | n | $n \log n$ | n^2 | n^3 | 2^n | $n!$ |
|-----------|----------|-----------|------------|-----------|-----------|-----------|------------|
| 10 | 3.3 | 10 | 33 | 100 | 1000 | 1000 | 10^6 |
| 100 | 6.6 | 100 | 66 | 10^4 | 10^6 | 10^{30} | 10^{157} |
| 1000 | 10 | 1000 | 10^4 | 10^6 | 10^9 | | |
| 10^4 | 13 | 10^4 | 10^5 | 10^8 | 10^{12} | | |
| 10^5 | 17 | 10^5 | 10^6 | 10^{10} | | | |
| 10^6 | 20 | 10^6 | 10^7 | | | | |
| 10^7 | 23 | 10^7 | 10^8 | | | | |
| 10^8 | 27 | 10^8 | 10^9 | | | | |
| 10^9 | 30 | 10^9 | 10^{10} | | | | |
| 10^{10} | 33 | 10^{10} | | | | | |