

Lecture 3

Discrete-time Markov Chains...

Dr. Dave Parker



Department of Computer Science
University of Oxford

Next few lectures...

- Today:
 - Discrete-time Markov chains (continued)
- Mon 2pm:
 - Probabilistic temporal logics
- Wed 3pm:
 - PCTL model checking for DTMCs
- Thur 12pm:
 - PRISM

Overview

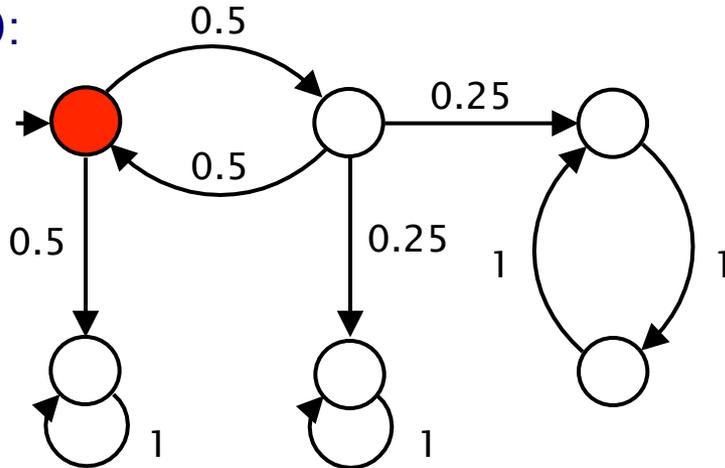
- Transient state probabilities
- Long-run / steady-state probabilities
- Qualitative properties
 - repeated reachability
 - persistence

Transient state probabilities

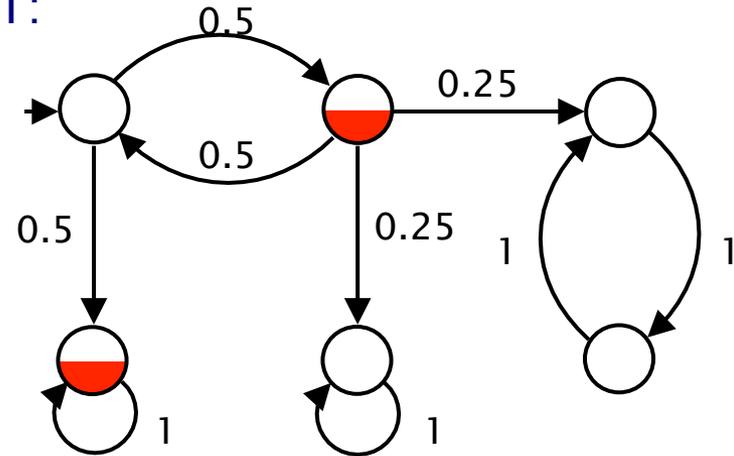
- What is the probability, having started in state s , of being in state s' at time k ?
 - i.e. after exactly k steps/transitions have occurred
 - this is the **transient state probability**: $\pi_{s,k}(s')$
- **Transient state distribution**: $\underline{\pi}_{s,k}$
 - vector $\underline{\pi}_{s,k}$ i.e. $\pi_{s,k}(s')$ for all states s'
- **Note**: this is a **discrete** probability distribution
 - so we have $\underline{\pi}_{s,k} : S \rightarrow [0,1]$
 - rather than e.g. $\Pr_s : \Sigma_{\text{Path}(s)} \rightarrow [0,1]$ where $\Sigma_{\text{Path}(s)} \subseteq 2^{\text{Path}(s)}$

Transient distributions

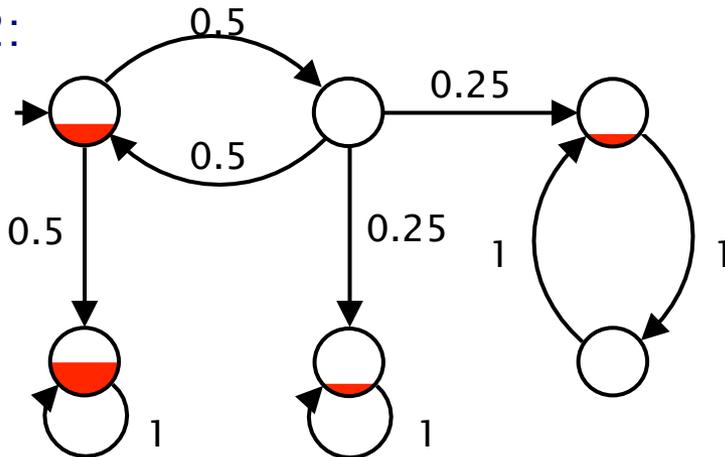
k=0:



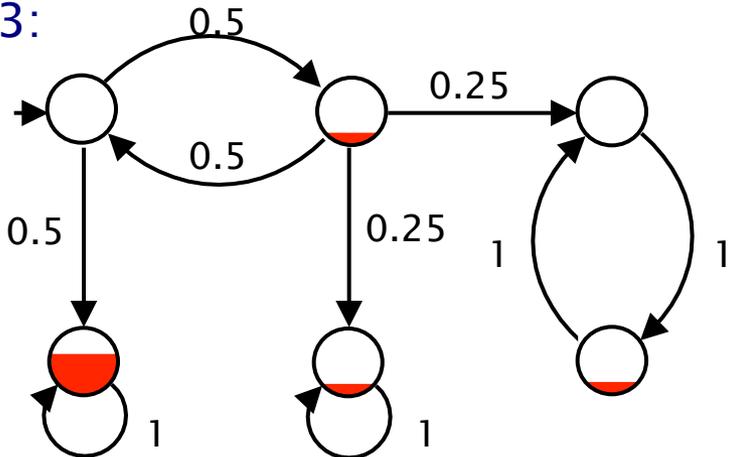
k=1:



k=2:



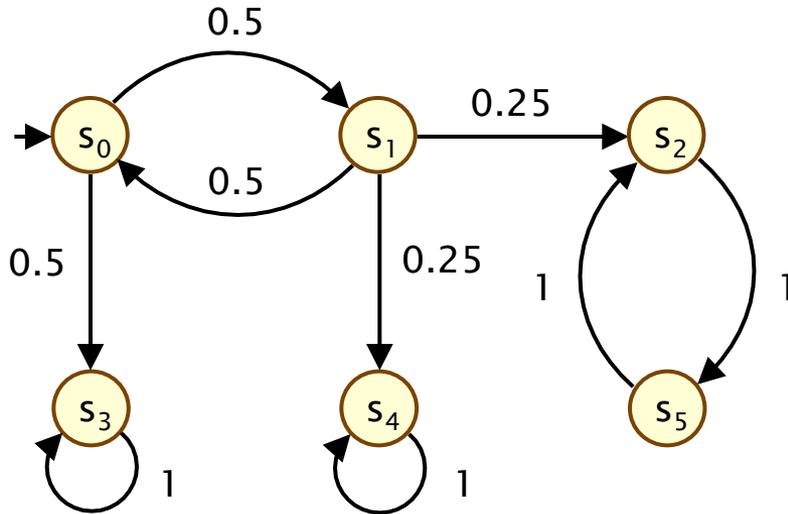
k=3:



Computing transient probabilities

- Transient state probabilities:
 - $\pi_{s,k}(s') = \sum_{s'' \in S} \mathbf{P}(s'', s') \cdot \pi_{s,k-1}(s'')$
 - (i.e. look at incoming transitions)
- Computation of transient state distribution:
 - $\underline{\pi}_{s,0}$ is the initial probability distribution
 - e.g. in our case $\underline{\pi}_{s,0}(s') = 1$ if $s' = s$ and $\underline{\pi}_{s,0}(s') = 0$ otherwise
 - $\underline{\pi}_{s,k} = \underline{\pi}_{s,k-1} \cdot \mathbf{P}$
- i.e. successive vector–matrix multiplications

Computing transient probabilities



$$\mathbf{P} = \begin{bmatrix}
 0 & 0.5 & 0 & 0.5 & 0 & 0 \\
 0.5 & 0 & 0.25 & 0 & 0.25 & 0 \\
 0 & 0 & 0 & 0 & 1 & 0 \\
 0 & 0 & 0 & 1 & 0 & 0 \\
 0 & 0 & 0 & 0 & 1 & 0 \\
 0 & 0 & 1 & 0 & 0 & 0
 \end{bmatrix}$$

$$\underline{\pi}_{s_0,0} = \begin{bmatrix} 1, 0, 0, 0, 0, 0 \end{bmatrix}$$

$$\underline{\pi}_{s_0,1} = \begin{bmatrix} 0, \frac{1}{2}, 0, \frac{1}{2}, 0, 0 \end{bmatrix}$$

$$\underline{\pi}_{s_0,2} = \begin{bmatrix} \frac{1}{4}, 0, \frac{1}{8}, \frac{1}{2}, \frac{1}{8}, 0 \end{bmatrix}$$

$$\underline{\pi}_{s_0,3} = \begin{bmatrix} 0, \frac{1}{8}, 0, \frac{5}{8}, \frac{1}{8}, \frac{1}{8} \end{bmatrix}$$

...

Computing transient probabilities

- $\underline{\pi}_{s,k} = \underline{\pi}_{s,k-1} \cdot \mathbf{P} = \underline{\pi}_{s,0} \cdot \mathbf{P}^k$
- k^{th} matrix power: \mathbf{P}^k
 - \mathbf{P} gives one-step transition probabilities
 - \mathbf{P}^k gives probabilities of k -step transition probabilities
 - i.e. $\mathbf{P}^k(s,s') = \pi_{s,k}(s')$
- A possible optimisation: iterative squaring
 - e.g. $\mathbf{P}^8 = ((\mathbf{P}^2)^2)^2$
 - only requires $\log k$ multiplications
 - but potentially inefficient, e.g. if \mathbf{P} is large and sparse
 - in practice, successive vector-matrix multiplications preferred

Notion of time in DTMCs

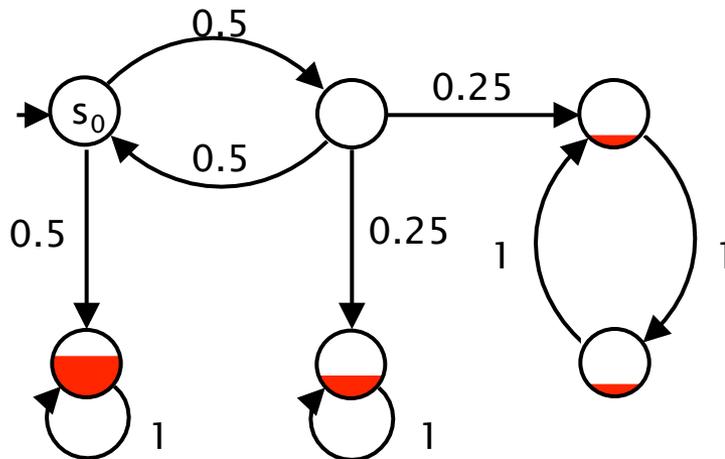
- Two possible views on the timing aspects of a system modelled as a DTMC:
- Discrete time-steps model time accurately
 - e.g. clock ticks in a model of an embedded device
 - or like dice example: interested in number of steps (tosses)
- Time-abstract
 - no information assumed about the time transitions take
 - e.g. simple Zeroconf model
- In the latter case, transient probabilities are not very useful
- In both cases, often beneficial to study long-run behaviour

Long-run behaviour

- Consider the limit: $\underline{\pi}_s = \lim_{k \rightarrow \infty} \underline{\pi}_{s,k}$
 - where $\underline{\pi}_{s,k}$ is the transient state distribution at time k having starting in state s
 - this limit, where it exists, is called the **limiting distribution**
- Intuitive idea
 - the percentage of time, in the long run, spent in each state
 - e.g. reliability: “in the long-run, what percentage of time is the system in an operational state”

Limiting distribution

- Example:



$$\underline{\pi}_{s_0,0} = \begin{bmatrix} 1, 0, 0, 0, 0, 0 \end{bmatrix}$$

$$\underline{\pi}_{s_0,1} = \begin{bmatrix} 0, \frac{1}{2}, 0, \frac{1}{2}, 0, 0 \end{bmatrix}$$

$$\underline{\pi}_{s_0,2} = \begin{bmatrix} \frac{1}{4}, 0, \frac{1}{8}, \frac{1}{2}, \frac{1}{8}, 0 \end{bmatrix}$$

$$\underline{\pi}_{s_0,3} = \begin{bmatrix} 0, \frac{1}{8}, 0, \frac{5}{8}, \frac{1}{8}, \frac{1}{8} \end{bmatrix}$$

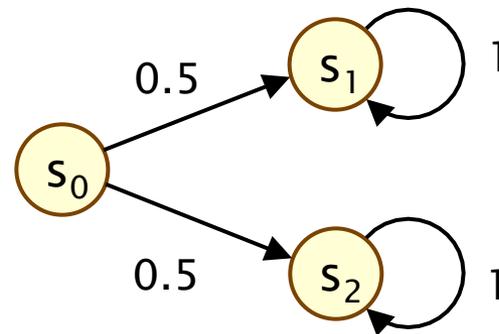
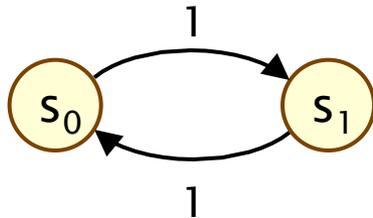
...

$$\underline{\pi}_{s_0} = \begin{bmatrix} 0, 0, \frac{1}{12}, \frac{2}{3}, \frac{1}{6}, \frac{1}{12} \end{bmatrix}$$

Long-run behaviour

- Questions:

- when does this limit exist?
- does it depend on the initial state/distribution?



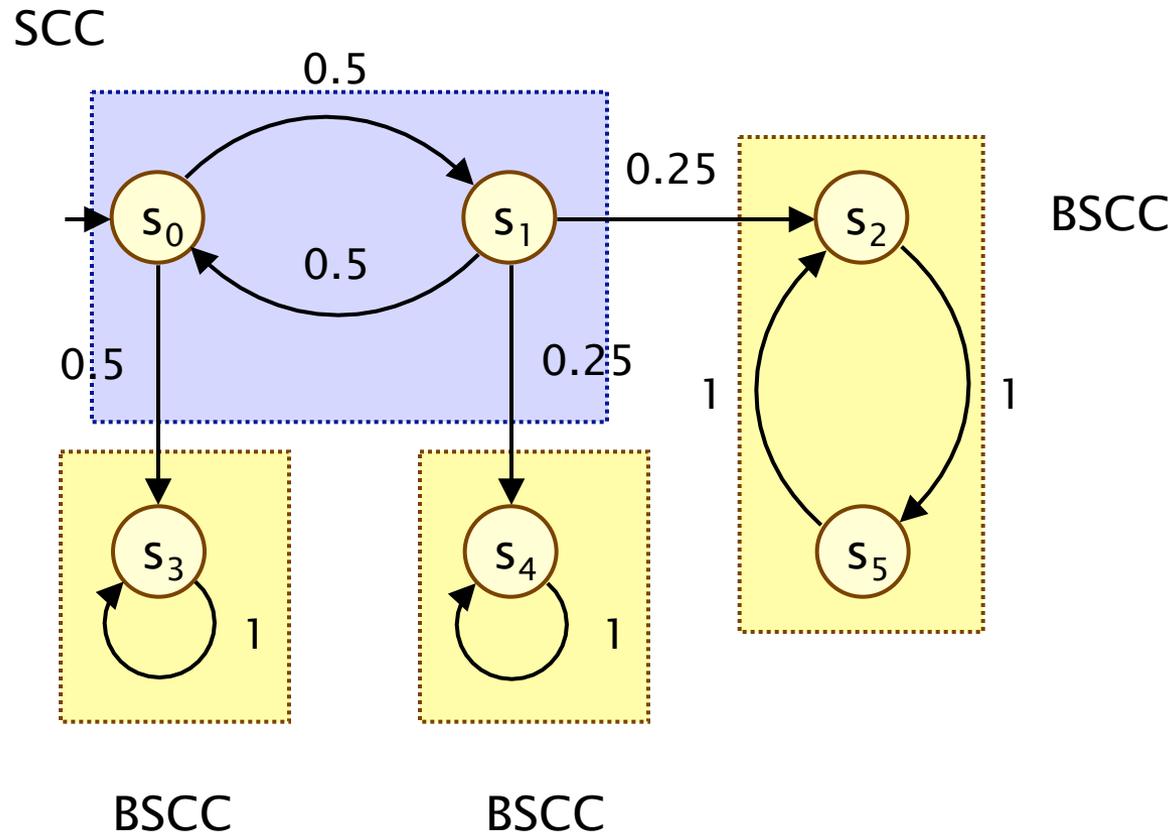
- Need to consider underlying graph

- (V,E) where V are vertices and $E \subseteq V \times V$ are edges
- $V = S$ and $E = \{ (s,s') \text{ s.t. } P(s,s') > 0 \}$

Graph terminology

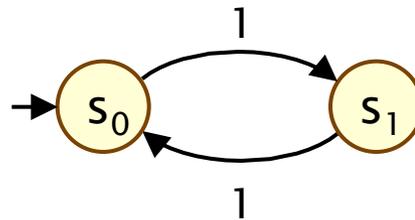
- A state s' is **reachable** from s if there is a finite path starting in s and ending in s'
- A subset T of S is **strongly connected** if, for each pair of states s and s' in T , s' is reachable from s passing only through states in T
- A **strongly connected component** (SCC) is a maximally strongly connected set of states (i.e. no superset of it is also strongly connected)
- A **bottom strongly connected component** (BSCC) is an SCC T from which no state outside T is reachable from T
- Alternative terminology: “ s communicates with s' ”, “communicating class”, “closed communicating class”

Example – (B)SCCs



Graph terminology

- Markov chain is **irreducible** if all its states belong to a single BSCC; otherwise reducible



- A state s is **periodic**, with period d , if
 - the greatest common divisor of the set $\{ n \mid f_s^{(n)} > 0 \}$ equals d
 - where $f_s^{(n)}$ is the probability of, when starting in state s , returning to state s in exactly n steps
- A Markov chain is **aperiodic** if its period is 1

Steady-state probabilities

- For a finite, irreducible, aperiodic DTMC...
 - limiting distribution always exists
 - and is independent of initial state/distribution
- These are known as steady-state probabilities
 - (or equilibrium probabilities)
 - effect of initial distribution has disappeared, denoted $\underline{\pi}$
- These probabilities can be computed as the unique solution of the linear equation system:

$$\underline{\pi} \cdot P = \underline{\pi} \quad \text{and} \quad \sum_{s \in S} \underline{\pi}(s) = 1$$

Steady-state – Balance equations

- Known as **balance equations**

$$\underline{\pi} \cdot P = \underline{\pi} \quad \text{and} \quad \sum_{s \in S} \underline{\pi}(s) = 1$$

- That is:

- $\underline{\pi}(s') = \sum_{s \in S} \underline{\pi}(s) \cdot P(s, s')$

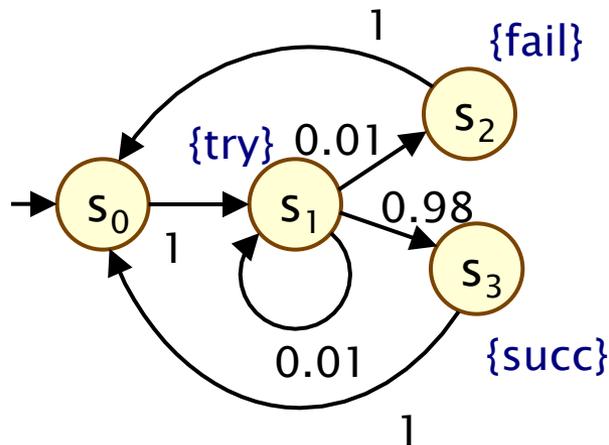
balance the probability of leaving and entering a state s'

- $\sum_{s \in S} \underline{\pi}(s) = 1$

normalisation

Steady-state – Example

- Let $\underline{x} = \underline{\pi}$
- Solve: $\underline{x} \cdot \mathbf{P} = \underline{x}$, $\sum_s \underline{x}(s) = 1$



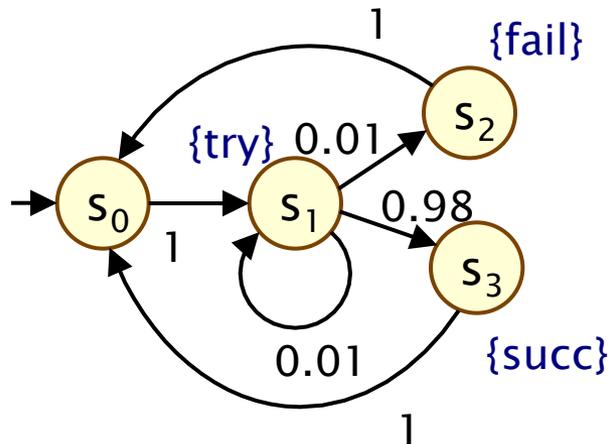
$$\mathbf{P} = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0.01 & 0.01 & 0.98 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}$$

$$\underline{x} \approx [0.332215, 0.335570, \\ 0.003356, 0.328859]$$

Steady-state – Example

- Let $\underline{x} = \underline{\pi}$
- Solve: $\underline{x} \cdot \mathbf{P} = \underline{x}, \sum_s \underline{x}(s) = 1$

$$\underline{x} \approx [0.332215, 0.335570, 0.003356, 0.328859]$$



$$\mathbf{P} = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0.01 & 0.01 & 0.98 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}$$

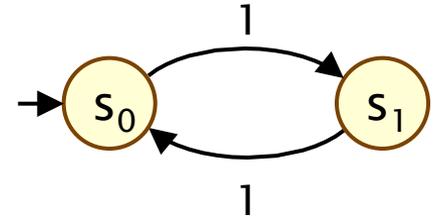
Long-run percentage of time spent in the state “try”
 $\approx 33.6\%$

Long-run percentage of time spent in “fail”/”succ”
 $\approx 0.003356 + 0.328859$
 $\approx 33.2\%$

Periodic DTMCs

- For (finite, irreducible) periodic DTMCs, this limit:

$$\underline{\pi}_s(s') = \lim_{k \rightarrow \infty} \underline{\pi}_{s,k}(s')$$



- does not exist, but this limit does:

$$\lim_{n \rightarrow \infty} \frac{1}{n} \cdot \sum_{k=1}^n \underline{\pi}_{s,k}(s')$$

(and where both limits exist,
e.g. for aperiodic DTMCs,
these 2 limits coincide)

- Steady-state probabilities for these DTMCs can be computed by solving the same set of linear equations:

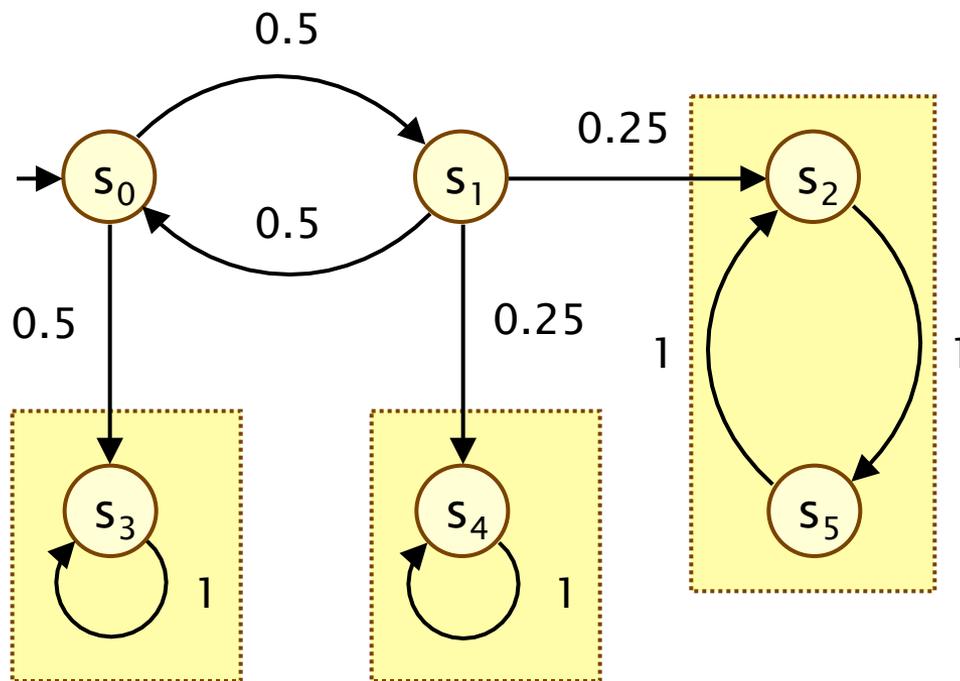
$$\underline{\pi} \cdot P = \underline{\pi} \quad \text{and} \quad \sum_{s \in S} \underline{\pi}(s) = 1$$

Steady-state – General case

- General case: reducible DTMC
 - compute vector $\underline{\pi}_s$
 - (note: distribution depends on initial state s)
- Compute BSCCs for DTMC; then two cases to consider:
- (1) s is in a BSCC T
 - compute steady-state probabilities \underline{x} in sub-DTMC for T
 - $\underline{\pi}_s(s') = \underline{x}(s')$ if s' in T
 - $\underline{\pi}_s(s') = 0$ if s' not in T
- (2) s is not in any BSCC
 - compute steady-state probabilities \underline{x}_T for sub-DTMC of each BSCC T and combine with reachability probabilities to BSCCs
 - $\underline{\pi}_s(s') = \text{ProbReach}(s, T) \cdot \underline{x}_T(s')$ if s' is in BSCC T
 - $\underline{\pi}_s(s') = 0$ if s' is not in a BSCC

Steady-state – Example 2

- $\underline{\pi}_s$ depends on initial state s



$$\underline{\pi}_{s_3} = [0 \ 0 \ 0 \ 1 \ 0 \ 0]$$

$$\underline{\pi}_{s_4} = [0 \ 0 \ 0 \ 0 \ 1 \ 0]$$

$$\underline{\pi}_{s_2} = \underline{\pi}_{s_5} = \left[0, 0, \frac{1}{2}, 0, 0, \frac{1}{2} \right]$$

$$\underline{\pi}_{s_0} = \left[0, 0, \frac{1}{12}, \frac{2}{3}, \frac{1}{6}, \frac{1}{12} \right]$$

$$\underline{\pi}_{s_1} = \dots$$

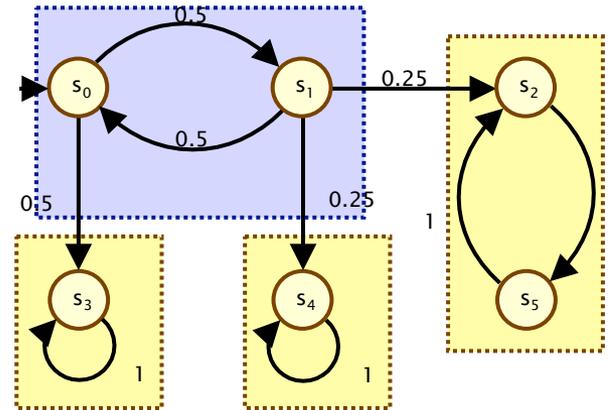
Qualitative properties

- **Quantitative** properties:
 - “what is the probability of event A?”
- **Qualitative** properties:
 - “the probability of event A is 1” (“almost surely A”)
 - or: “the probability of event A is > 0 ” (“possibly A”)
- For finite DTMCs, qualitative properties do not depend on the transition probabilities – only need underlying graph
 - e.g. to determine “is target set T reached with probability 1?” (see DTMC model checking lecture)
 - computing BSCCs of a DTMCs yields information about long-run qualitative properties...

Fundamental property

- Fundamental property of (finite) DTMCs...

- With probability 1, a BSCC will be reached and all of its states visited infinitely often

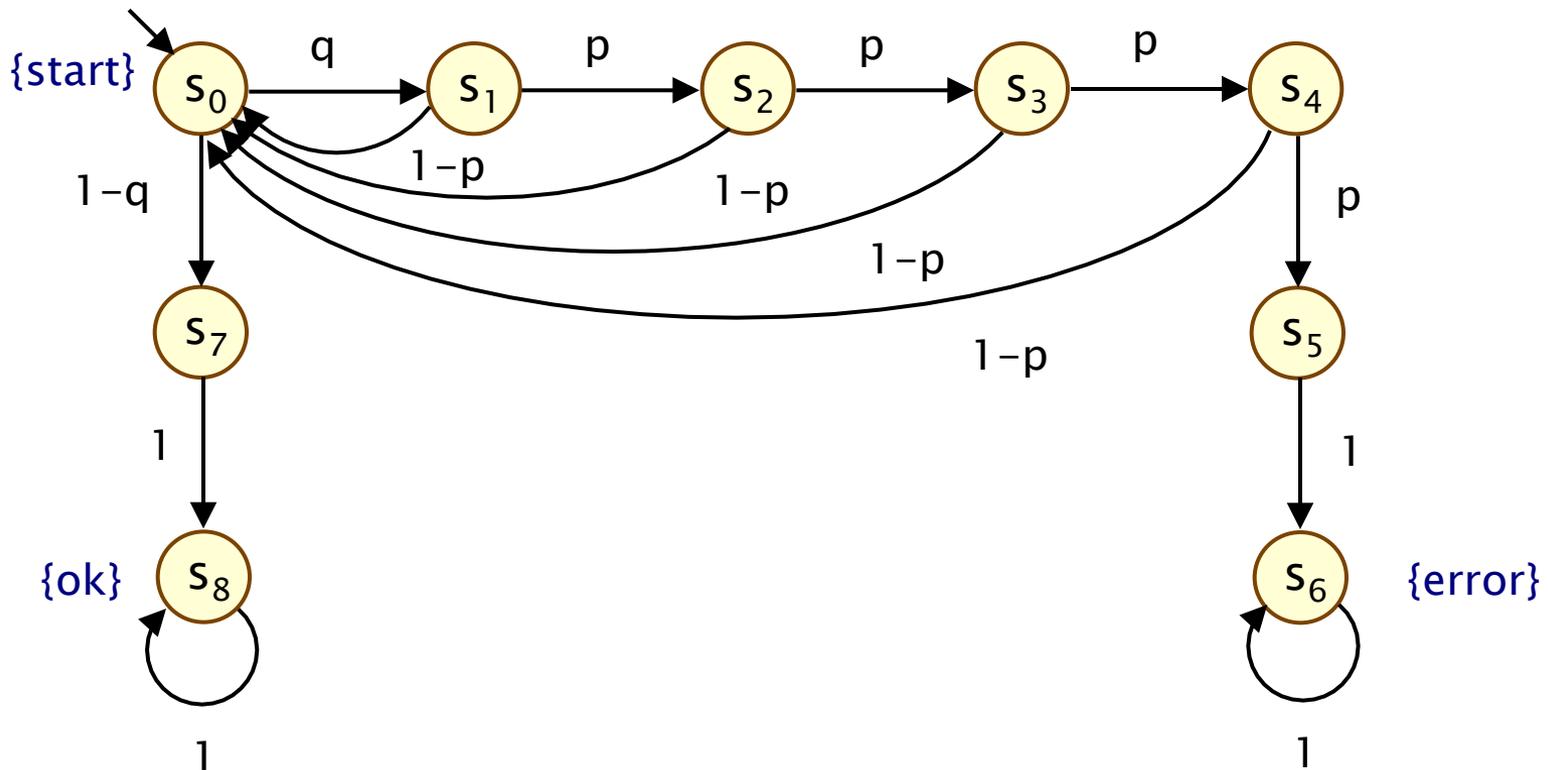


- Formally:

$$\begin{aligned} & - \Pr_{s_0} (s_0 s_1 s_2 \dots \mid \exists i \geq 0, \exists \text{ BSCC } T \text{ such that} \\ & \quad \forall j \geq i \ s_j \in T \text{ and} \\ & \quad \forall s \in T \ s_k = s \text{ for infinitely many } k) = 1 \end{aligned}$$

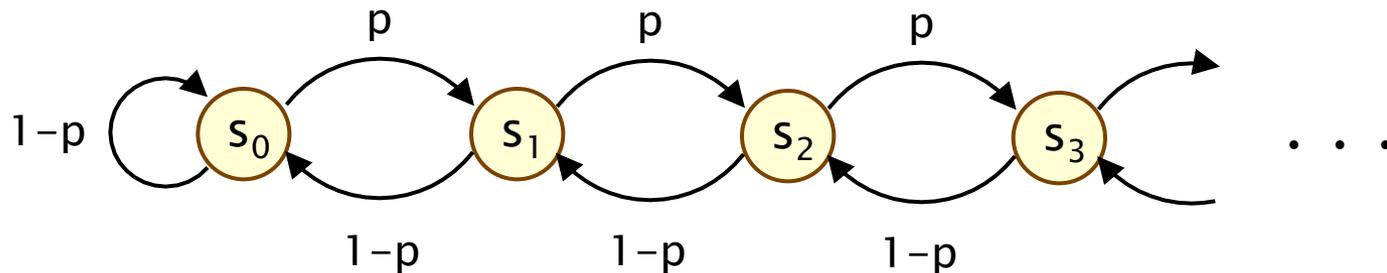
Zeroconf example

- 2 BSCCs: $\{s_6\}$, $\{s_8\}$
- Probability of trying to acquire a new address infinitely often is 0



Aside: Infinite Markov chains

- Infinite-state random walk



- Value of probability p **does** affect qualitative properties
 - $\text{ProbReach}(s, \{s_0\}) = 1$ if $p \leq 0.5$
 - $\text{ProbReach}(s, \{s_0\}) < 1$ if $p > 0.5$

Repeated reachability

- Repeated reachability:
 - “always eventually...”, “infinitely often...”
- $\Pr_{s_0} (s_0 s_1 s_2 \dots \mid \forall i \geq 0 \exists j \geq i s_j \in B)$
 - where $B \subseteq S$ is a set of states
- e.g. “what is the probability that the protocol successfully sends a message infinitely often?”
- Is this measurable? Yes...
 - set of satisfying paths is: $\bigcap_{n \geq 0} \bigcup_{m \geq n} C_m$
 - where C_m is the union of all cylinder sets $\text{Cyl}(s_0 s_1 \dots s_m)$ for finite paths $s_0 s_1 \dots s_m$ such that $s_m \in B$

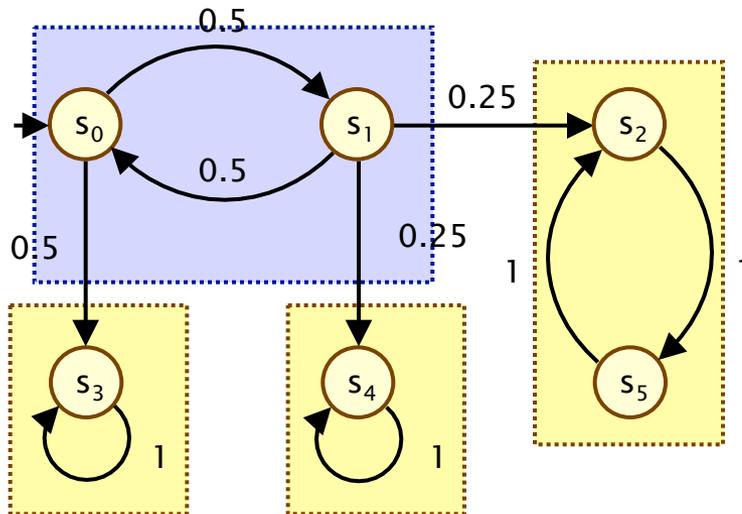
Qualitative repeated reachability

- $\Pr_{s_0} (s_0 s_1 s_2 \dots \mid \forall i \geq 0 \exists j \geq i s_j \in B) = 1$
 $\Pr_{s_0} (\text{“always eventually B”}) = 1$

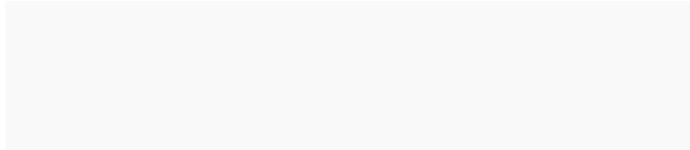
if and only if

- $T \cap B \neq \emptyset$ for each BSCC T that is reachable from s_0

Example:
 $B = \{ s_3, s_4, s_5 \}$



Persistence

- Persistence properties:
 - “eventually forever...”
- $\Pr_{s_0} (s_0 s_1 s_2 \dots \mid \exists i \geq 0 \forall j \geq i s_j \in B)$
 - where $B \subseteq S$ is a set of states
- e.g. “what is the probability of the leader election algorithm reaching, and staying in, a stable state?”
- e.g. “what is the probability that an irrecoverable error occurs?”
- Is this measurable? Yes... 

Qualitative persistence

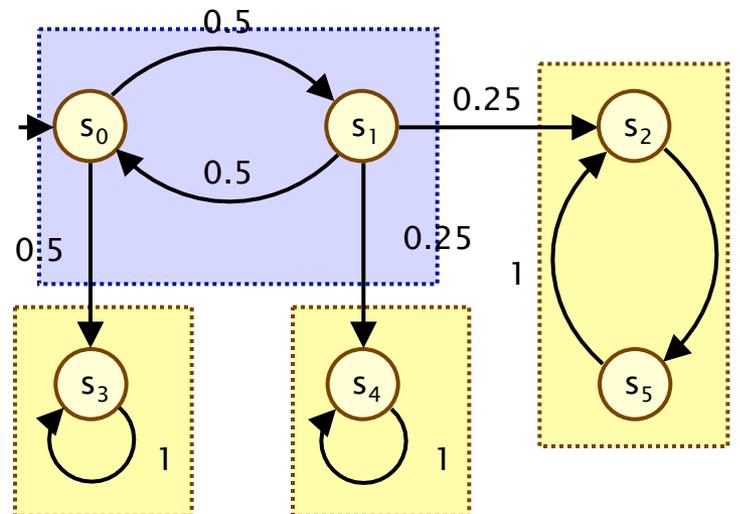
- $\Pr_{s_0} (s_0 s_1 s_2 \dots \mid \exists i \geq 0 \forall j \geq i s_j \in B) = 1$
 $\Pr_{s_0} (\text{“eventually forever B”}) = 1$

if and only if

- $T \subseteq B$ for each BSCC T that is reachable from s_0

Example:

$$B = \{ s_2, s_3, s_4, s_5 \}$$



Summing up...

- **Transient state probabilities**
 - successive vector–matrix multiplications
- **Long–run/steady–state probabilities**
 - requires graph analysis
 - irreducible case: solve linear equation system
 - reducible case: steady–state for sub–DTMCs + reachability
- **Qualitative properties**
 - repeated reachability
 - persistence