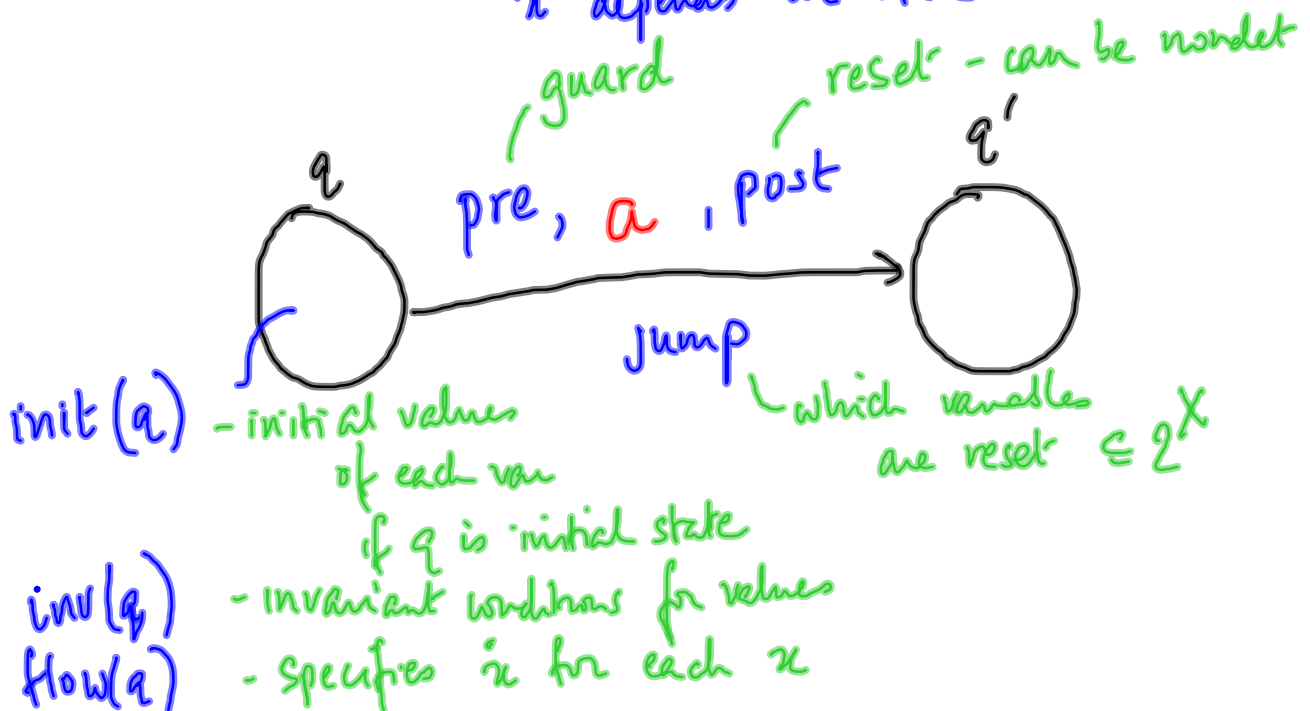


Hybrid Automata

Clocks $\rightarrow$ arbitrary real valued variables

x depends on state



What's decidable about hybrid automata?

Henzinger et al, TCSS

$$X = \{x_1, x_2, \dots, x_n\}$$

$\subseteq \mathbb{R}^n$

$$V: X \rightarrow \mathbb{R}$$

Rectangular hybrid automata

$R \subseteq \mathbb{R}^n$ is a rectangle if it is a product of intervals

Rectangular hybrid automaton

flow(a)
inv(q)
init(q)

pre(t)
post(t)

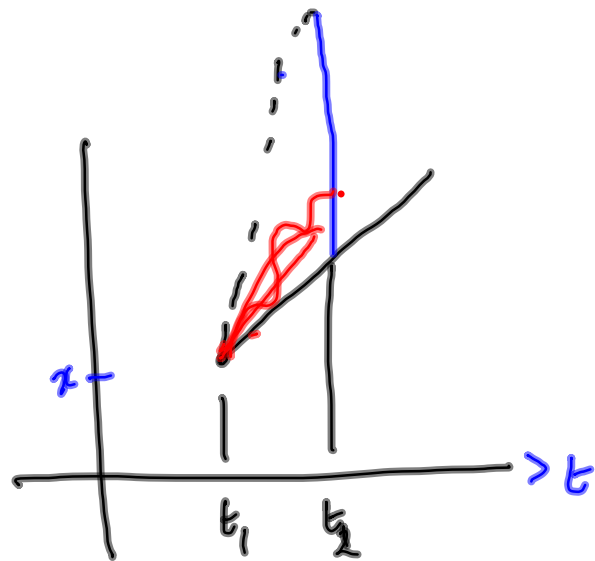
} a conjunction of
one rectangular
constraint per variable

post(t) : $x = 3y + 7z + 8y^3$ X

$x \in [l, u)$

$\text{flow}(\pi) \in [1, 5)$

effectively assume
evolution is
piecewise linear



Initialized automaton

$q \xrightarrow{a} q'$ if $\text{flow}(q)(x) \neq \text{flow}(q')(x)$
then $x \in \text{jump}(t)$

Whenever i changes, x is reset

Initialized + rectangular $\Leftrightarrow$ positive

Deterministic reset

Reachability is decidable for initialized rectangular automata

reachability of (q, Z)

where Z is a "zone" or rectangle

Initialized Rectangular



Initialized Singular

- each π is a constant
deterministic resets.



Initialized Stopwatch

$\pi \in \{0,1\}$

deterministic resets



Timed Automata

with constant resets

Initialized Stopwatch



Timed Automata

$(q, x_1, x_2, \dots, x_n, \text{some } \dot{x}_i = 0)$

If $\dot{x}_i = 0$, x_i is $\text{post}(t)$
for last transition when
 x_i went from 1 to 0

$(q, x_1, \dots, x_n, f: X \rightarrow \mathbb{R} \cup \{\perp\})$

If $\dot{x}_i = 1$, same value as stopwatch and

If $\dot{x}_i = 0$, $f(x) = \perp$
 $f(x) = \text{stopped value of } x$

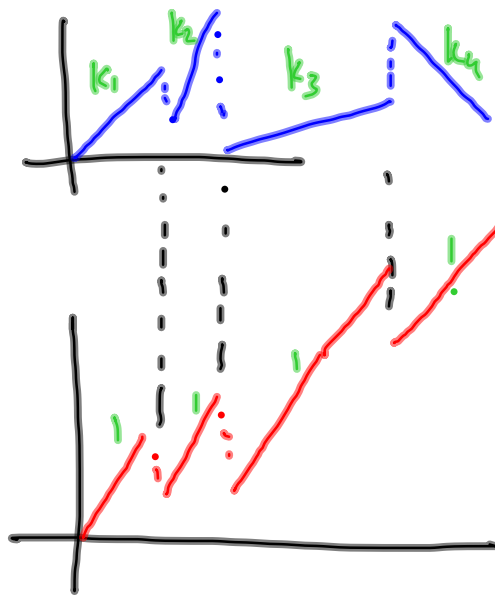
Initialized singular



Initialized stopwatch

Simple scaling
relates configurations

$\forall q \exists k \ x \in [k, k]$ in q .
det. jumps



Initialized Rectangular

$$x \in [l, u]$$



Initialized Singular

$$\begin{aligned} & \swarrow \searrow \\ & x_l = l \quad x_u = u \end{aligned}$$

"Compact" rectangles
"bounded & closed"

