

Programming and Data Structures with Python

Madhavan Mukund

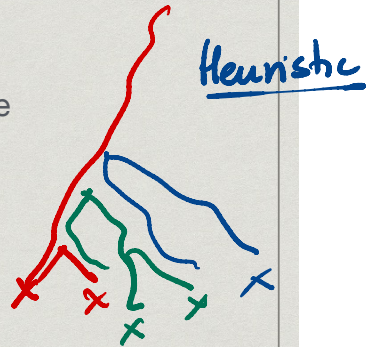
<https://www.cmi.ac.in/~madhavan>

6 December, 2021

Backtracking

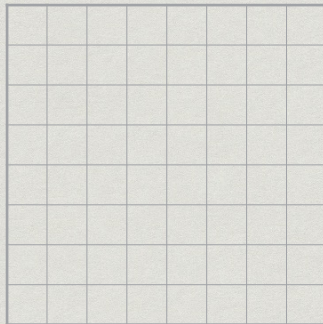
Playing Sudoku
Exiting a maze

- * Systematically search for a solution
- * Build the solution one step at a time
- * If we hit a dead-end
 - * Undo the last step
 - * Try the next option



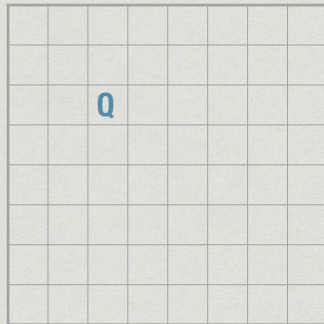
Eight queens

- * Place 8 queens on a chess board so that none of them attack each other
- * In chess, a queen can move any number of squares along a row column or diagonal



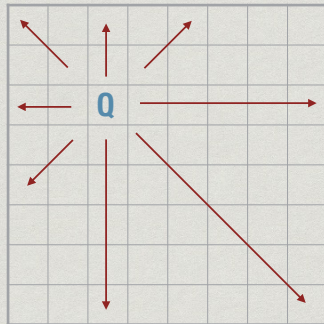
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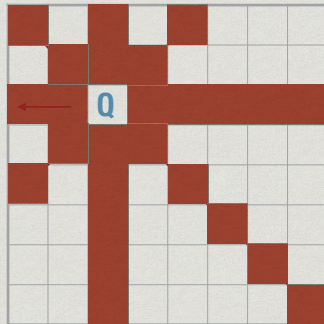
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N queens

- * Place N queens on an N x N chess board so that none attack each other
- * $N = 2, 3$ impossible

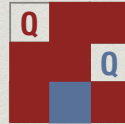
N queens

- * Place N queens on an $N \times N$ chess board so that none attack each other
- * $N = 2, 3$ impossible



N queens

- * Place N queens on an N x N chess board so that none attack each other
- * N = 2, 3 impossible
- * N = 4 is possible



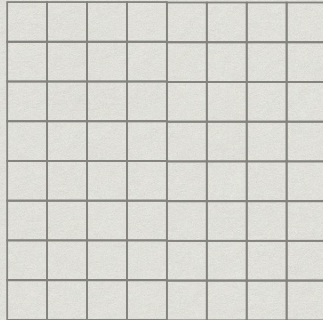
N queens

- * Place N queens on an $N \times N$ chess board so that none attack each other
- * $N = 2, 3$ impossible
- * $N = 4$ is possible
- * And all bigger N as well



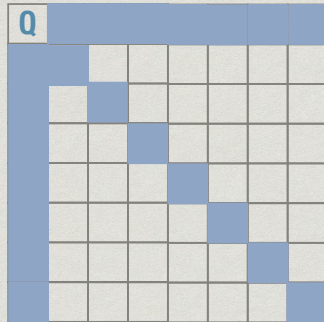
8 queens

- * Clearly, exactly one queen in each row, column
- * Place queens row by row
- * In each row, place a queen in the first available column



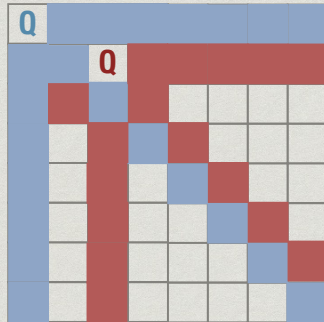
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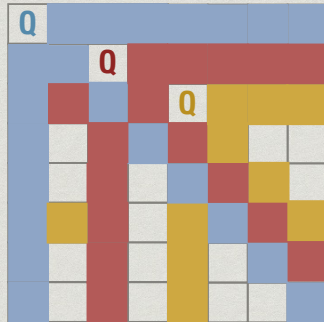
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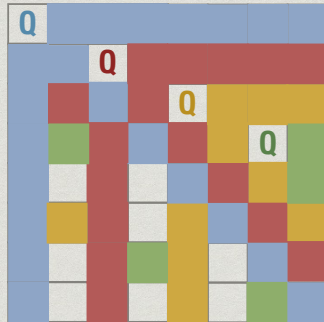
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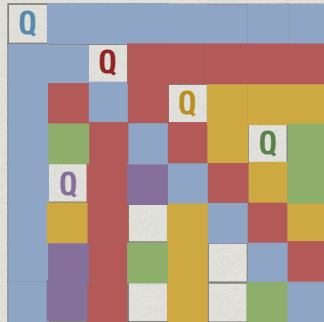
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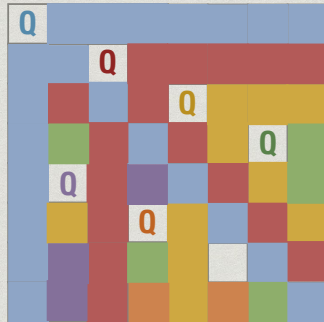
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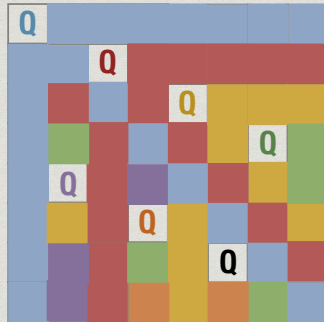
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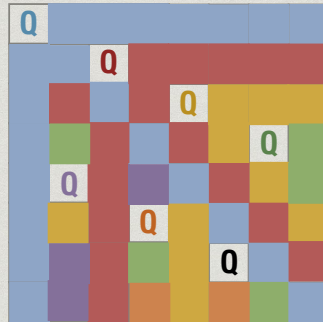
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8 queens

- * Clearly, exactly one queen in each row, column
- * Place queens row by row
- * In each row, place a queen in the first available column
- * Can't place a queen in the 8th row!



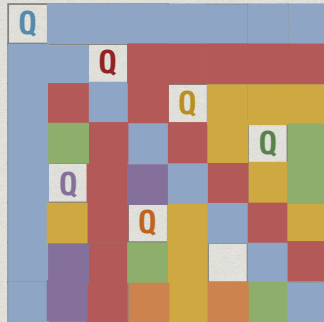
8 queens

- * Can't place the a queen in the 8th row!



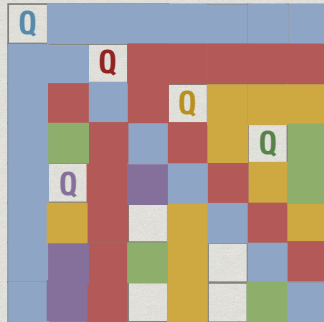
8 queens

- * Can't place the a queen in the 8th row!
- * Undo 7th queen, no other choice



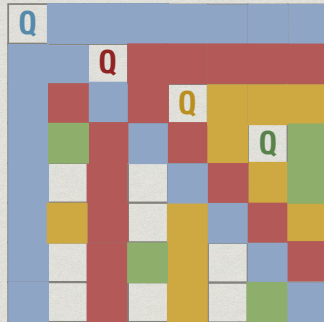
8 queens

- * Can't place the a queen in the 8th row!
- * Undo 7th queen, no other choice
- * Undo 6th queen, no other choice



8 queens

- * Can't place the a queen in the 8th row!
- * Undo 7th queen, no other choice
- * Undo 6th queen, no other choice
- * Undo 5th queen, try next



8 queens

- * Can't place the a queen in the 8th row!
- * Undo 7th queen, no other choice
- * Undo 6th queen, no other choice
- * Undo 5th queen, try next



Backtracking

- * Keep trying to extend the next solution
- * If we cannot, undo previous move and try again
- * Exhaustively search through all possibilities
- * ... but systematically!

Coding the solution

- * How do we represent the board?
- * $n \times n$ grid, number rows and columns from 0 to $n-1$
 - * $\text{board}[i][j] == 1$ indicates queen at (i, j)
 - * $\text{board}[i][j] == 0$ indicates no queen
- * We know there is only one queen per row
- * Single list board of length n with entries 0 to $n-1$
 - * $\text{board}[i] == j$: queen in row i , column j , i.e. (i, j)

$\text{board}[5][7]$
 $= 1$

$\forall j \neq 7$

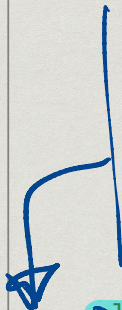
$\text{board}[5][j]$
 $= 0$

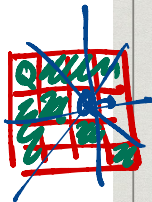
Overall structure

```
def placequeen(i,board): # Trying row i
    for each c such that (i,c) is available:
        place queen at (i,c) and update board
        if i == n-1:
            return(True) # Last queen has been placed
        else:
            extendsoln = placequeen(i+1,board)
            if extendsoln:
                return(True) # This solution extends fully
            else:
                undo this move and update board
            else:
            return(False) # Row i failed
```

current state of board
row 0 to i-1
are occupied

updated board





Updating the board

- * Our 1-D and 2-D representations keep track of the queens
- * Need an efficient way to compute which squares are free to place the next queen
- * $n \times n$ attack grid
 - * $\text{attack}[i][j] == 1$ if (i, j) is attacked by a queen
 - * $\text{attack}[i][j] == 0$ if (i, j) is currently available
- * How do we undo the effect of placing a queen?
 - * Which $\text{attack}[i][j]$ should be reset to 0?

board[i][j] keeps track of queens

Updating the board

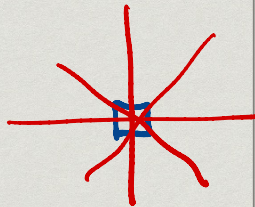
- * Queens are added row by row
- * Number the queens 0 to $n-1$
- * Record earliest queen that attacks each square
 - * `attack[i][j] == k` if (i, j) was first attacked by queen k
 - * `attack[i][j] == -1` if (i, j) is free
- * Remove queen k — reset `attack[i][j] == k` to -1
 - * All other squares still attacked by earlier queens

Updating the board

- * `attack` requires n^2 space
 - * Each update only requires $O(n)$ time
 - * Only need to scan row, column, two diagonals
- * Can we improve our representation to use only $O(n)$ space?

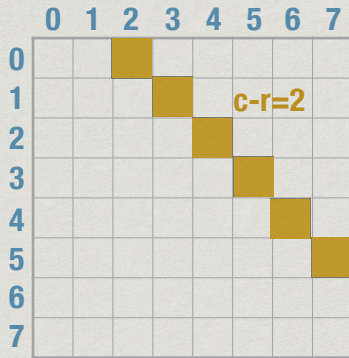
A better representation

- * How many queens attack row i ?
- * How many queens attack row j ?
- * An individual square (i,j) is attacked by upto 4 queens
 - * Queen on row i and on column j
 - * One queen on each diagonal through (i,j)



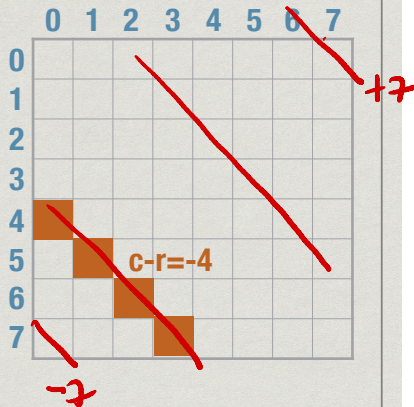
Numbering diagonals

- * Decreasing diagonal:
(column - row) is invariant



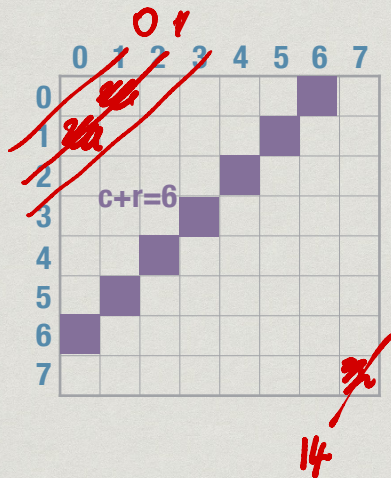
Numbering diagonals

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Numbering diagonals

- * Decreasing diagonal:
column - row is invariant
- * Increasing diagonal:
column + row is invariant



Numbering diagonals

- * Decreasing diagonal:
column - row is invariant
- * Increasing diagonal:
column + row is invariant

	0	1	2	3	4	5	6	7
0								
1								
2								
3								
4								
5								
6								
7								

$c+r=12$

Numbering diagonals

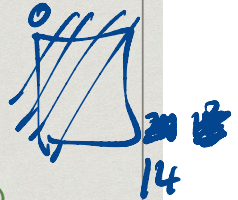
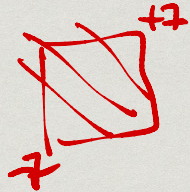
- * Decreasing diagonal:
column - row is invariant
- * Increasing diagonal:
column + row is invariant
- * (i,j) is attacked if
 - * row i is attacked
 - * column j is attacked
 - * diagonal $j-i$ is attacked
 - * diagonal $j+i$ is attacked

	0	1	2	3	4	5	6	7
0								
1								
2								
3								
4								
5								
6								
7								

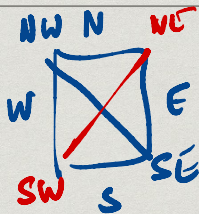
$c+r=12$

$O(n)$ representation

- * $\text{row}[i] == 1$ if row i is attacked, $0..N-1$
- * $\text{col}[i] == 1$ if column i is attacked, $0..N-1$
- * $\text{NWtoSE}[i] == 1$ if NW to SE diagonal i is attacked, $-(N-1)$ to $(N-1)$
- * $\text{SWtoNE}[i] == 1$ if SW to NE diagonal i is attacked, 0 to $2(N-1)$



Updating the board



- * (i, j) is free if
 $\text{row}[i] == \text{col}[j] == \text{NWtoSE}[j-i] == \text{SWtoNE}[j+i] == 0$
- * Add queen at (i, j)
 $\text{board}[i] = j$
 $(\text{row}[i], \text{col}[j], \text{NWtoSE}[j-i], \text{SWtoNE}[j+i]) =$
 $(1, 1, 1, 1)$
- * Remove queen at (i, j)
 $\text{board}[i] = -1$
 $(\text{row}[i], \text{col}[j], \text{NWtoSE}[j-i], \text{SWtoNE}[j+i]) =$
 $(0, 0, 0, 0)$

n^2 attack
matrix
with $O(n)$
update time
to
 $O(n)$ attack
arrays
 $O(1)$ update
time

Implementation details

- * Maintain `board` as nested dictionary
 - * `board['queen'][i] = j` : Queen located at (i, j)
 - * `board['row'][i] = 1` : Row i attacked
 - * `board['col'][i] = 1` : Column i attacked
 - * `board['nwtose'][i] = 1` : NWtoSW diagonal i attacked
 - * `board['swtone'][i] = 1` : SWtoNE diagonal i attacked

Overall structure

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    for each c such that (i,c) is available:
        place queen at (i,c) and update board
        if i == n-1:
            return(True) # Last queen has been placed
        else:
            extendsoln = placequeen(i+1,board)
            if extendsoln:
                return(True) # This solution extends fully
            else:
                undo this move and update board
    else:
        return(False) # Row i failed
```

All solutions?

```
def placequeen(i,board): # Try row i
    for each c such that (i,c) is available:
        place queen at (i,c) and update board
        if i == n-1:
            record solution # Last queen placed
        else:
            extendsoln = placequeen(i+1,board)
            undo this move and update board
```

Global variables

- * Can we avoid passing `board` explicitly to each function?
- * Can we have a single `global` copy of `board` that all functions can update?

Scope of name

- * Scope of name is the portion of code where it is available to read and update
- * By default, in Python, scope is local to functions
 - * But actually, only if we update the name inside the function

Two examples

```
def f():  
    y = x  
    print(y)
```

```
x = 7  
f()
```



Fine!

Two examples

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def f():  
    y = x  
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x = 7  
f()
```

Fine!

```
def f():  
    y = x  
    print(y)  
    x = 22
```

```
x = 7  
f()
```

Error!

different x

*makes x
local to f()*

Two examples

```
def f():  
    y = x  
    print(y)
```

```
x = 7  
f()
```

Fine!

```
def f():  
    y = x  
    print(y)  
    x = 22
```

```
x = 7  
f()
```

Error!

- * If `x` is not found in `f()`, Python looks at enclosing function for **global** `x`
- * If `x` is updated in `f()`, it becomes a **local** name!

Global variables

- * Actually, this applies only to immutable values
- * Global names that point to mutable values can be updated within a function

```
def f():  
    y = x[0]  
    print(y)  
    x[0] = 22
```

```
x = [7]
```

```
f()
```

Fine!

*x = [22]
world change
x*

Global immutable values

- * What if we want a global integer
- * Count the number of times a function is called
- * Declare a name to be `global`

```
def f():  
    global x  
    y = x  
    print(y)  
    x = 22
```

```
x = 7  
f()  
print(x)
```

Global immutable values

- * What if we want a global integer
- * Count the number of times a function is called
- * Declare a name to be `global`

```
def f():  
    global x  
    y = x  
    print(y)  
    x = 22  
x = 7  
f()  
print(x)
```

Handwritten annotations:

- A red circle around `global x`.
- A green circle around `x` in `y = x`.
- A green circle around `x` in `x = 22`.
- A green circle around `x` in `x = 7`.
- A green arrow from the `x` in `y = x` to the `x` in `x = 22`.
- A green arrow from the `x` in `x = 7` to the `x` in `x = 22`.
- The word *same* in green, with arrows pointing to the `x` in `x = 22` and the `x` in `x = 7`.
- The number **22** in red at the bottom right.

Nest function definitions

- * Can define local “helper” functions
- * `g()` and `h()` are only visible to `f()`
- * Cannot be called directly from outside

```
def f():  
    def g(a):  
        return(a+1)  
  
    def h(b):  
        return(2*b)  
  
    global x  
    y = g(x) + h(x)  
    print(y)  
    x = 22
```

```
x = 7
```

```
f()
```

~~`g(7)`~~ X

Nest function definitions

- * If we look up `x`, `y` inside `g()` or `h()` it will first look in `f()`, then outside
- * Can also declare names global inside `g()`, `h()`
- * Intermediate scope declaration: `nonlocal`
- * See Python documentation

```
def f():  
    def g(a):  
        return(a+1)
```

```
    def h(b):  
        return(2*b)
```

```
    global x  
    y = g(x) + h(x)  
    print(y)  
    x = 22
```

```
x = 7  
f()
```

Generating permutations

- * Often useful when we need to try out all possibilities
 - * Each potential columnwise placement of N queens is a permutation of $\{0, 1, \dots, N-1\}$
- * Given a permutation, generate the next one
- * For instance, what is the next sequence formed from $\{a, b, \dots, m\}$, in dictionary order after

d c h b a e g l k o n m j i

$\{a, b, c, d\}$

Smallest ?

a b c d

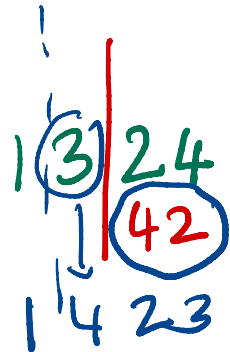
1 2 3 4

largest ?

d c b a

4 3 2 1

a c b d



Generating permutations

- * Smallest permutation — all elements in ascending order

a b c d e f g h i j k l m

- * Largest permutation — all elements in descending order

m l k j i h g f e d c b a

- * Next permutation — find shortest suffix that can be incremented
 - * Or longest suffix that cannot be incremented

Next permutation

- * Longest suffix that cannot be incremented
- * Already in descending order

d c h b a e g l **k** o n m j i

The diagram illustrates the next permutation algorithm. The sequence of letters is "d c h b a e g l k o n m j i". A red circle highlights the letter 'k', which is the pivot element. A red line connects 'k' to the letter 'o', indicating a swap. A blue line underlines the entire sequence, and a blue arrow points from the end of the underline to the letter 'k', indicating the pivot element.

Next permutation

- * Longest suffix that cannot be incremented
 - * Already in descending order

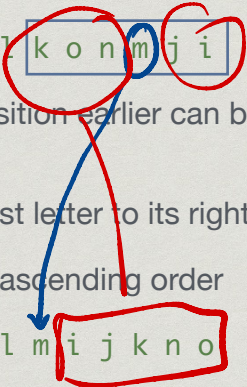
d c h b a e g l k o n m j i

- * The suffix starting one position earlier can be incremented

Next permutation

- * Longest suffix that cannot be incremented
 - * Already in descending order

d c h b a e g l k o n m j i




- * The suffix starting one position earlier can be incremented
 - * Replace **k** by next largest letter to its right, **m**
 - * Rearrange **k o n j i** in ascending order

d c h b a e g l m i j k n o

Implementation


- * From the right, identify first decreasing position

d c h b a e g l k o n m j i



- * Swap that value with its next larger letter to its right

d c h b a e g l m o n k j i



- * Finding next larger letter is similar to insert
- * Reverse the increasing suffix

d c h b a e g l m i j k n o