## Advanced Programming, II Semester, 2014–2015 Quiz 2, 2 March 2015

Answer all questions in the space provided. Use the designated space for rough work, if any. Don't forget to fill your name!

1. Complete the following function definition so that it behaves as described below—that is, fill in the parameters for **f**() in the correct order, with default values, as appropriate.

def f(....):
 print("a",a,"b",b,"c",c,"d",d)

Expected behaviour:

```
>>> f(b=4,a=3)
a 3 b 4 c 10 d 15
>>> f(3,5,7)
a 3 b 5 c 7 d 15
>>> f(3,c=7)
Traceback (most recent call last):
   File "<stdin>", line 1, in <module>
TypeError: f() takes at least 2 arguments (2 given)
```

(4 marks)

Solution:

def f(a,b,c=10,d=15):

- $\bullet\,$  The first line gives us the default values for c and d.
- The second line tells us the positions for a and b.
- The third line tells us that neither **a** nor **b** has default values.

Rough work

2. The function minout(inl) takes a list inl of distinct natural numbers (i.e., integers from the set {0,1,2,...}) and returns the smallest natural number not present in inl. In other words, if 0 is not in inl, then minout(inl) returns 0, else if 0 is in inl but 1 is not in inl, them minout(inl) returns 1, etc.

For example, minout([1,3,2,4,17]) is 0 and minout([1,3,0,2,4]) is 5.

Here is a recursive algorithm to compute minout. (Note that inl is *not* assumed to be sorted, nor do we make any assumptions about the range of values in inl.)

- Suppose the length of inl is *n*. Construct two lists lower and upper, where lower contains elements smaller than  $\lfloor \frac{n}{2} \rfloor$  and upper contains the rest.
- If the size of lower is strictly smaller than  $\frac{n}{2}$ , the missing number is smaller than  $\frac{n}{2}$ , so recursively invoke minout on lower.
- Otherwise, invoke minout on upper. (All numbers in upper are bigger than  $\frac{n}{2}$ , so some offset is required.)

Analyze the running time of this algorithm. (Write a recurrence and solve it.) (6 marks)

## Solution:

The recurrence is:

- T(1) = 1
- $T(n) = T\left(\frac{n}{2}\right) + n$

Unwinding this we get  $n + \frac{n}{2} + \frac{n}{4} + \dots = O(n)$ .

Rough work