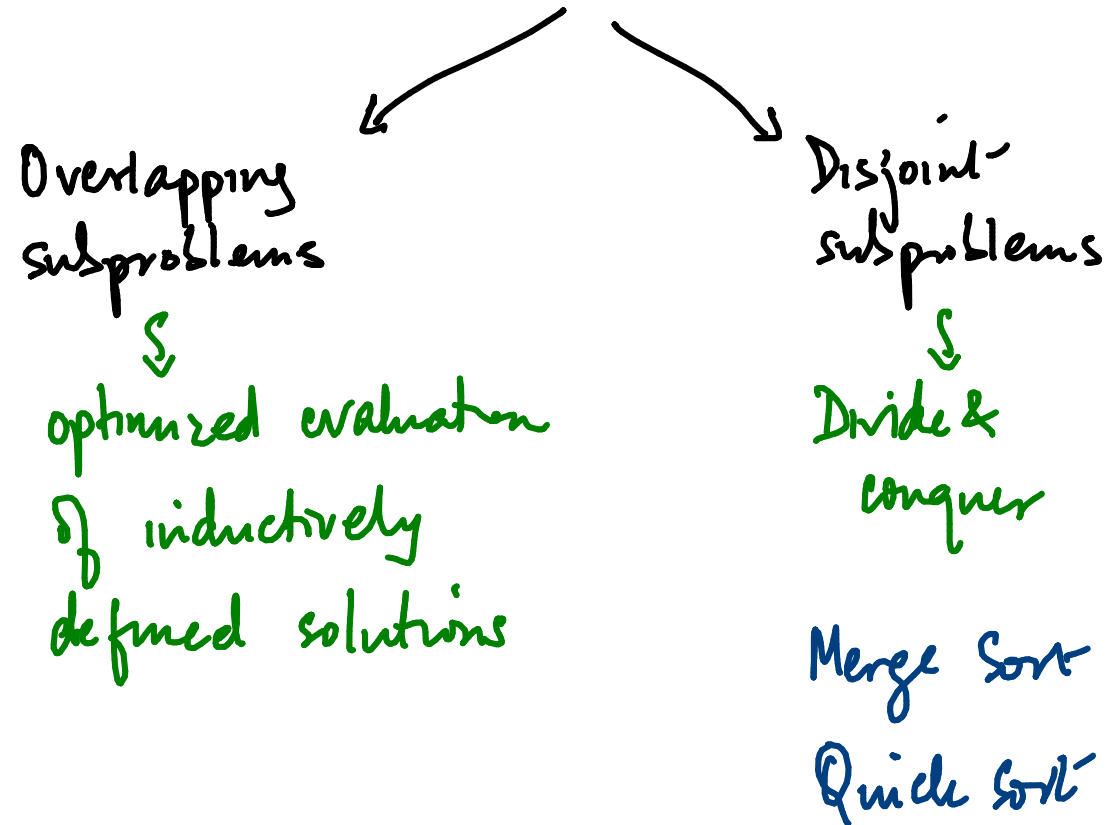


Greedy does not work  $\Rightarrow$  need to consider all subproblems



Recommendation systems / collaborative filtering

Customers like you also bought XYZ

Assume customers rate their purchases

Customers with similar rankings are "alike"

One proposal:

Items  $i_1, i_2, \dots, i_n$

Customer 1 ranks them  $r_1 > r_2 \dots > r_n$

2 " "  $s_1 > s_2 \dots > s_n$

Count how many disagreements on  $(i_j, i_k)$

Simplification:

Take customer 1 as reference:  $1, 2, \dots, n$

customer 2 is a permutation:  $\pi_1, \pi_2, \dots, \pi_n$

Count inversions w.r.t. sorted order

Problem

Given a permutation of  $1..n$ , count inversions  
(i.e.  $i > j$  but  $i$  appears before  $j$  in permutation)

HINT: DIVIDE & CONQUER

Split list into two equal parts

- Count inversions in each half (inductively)
- Count inversions across halves

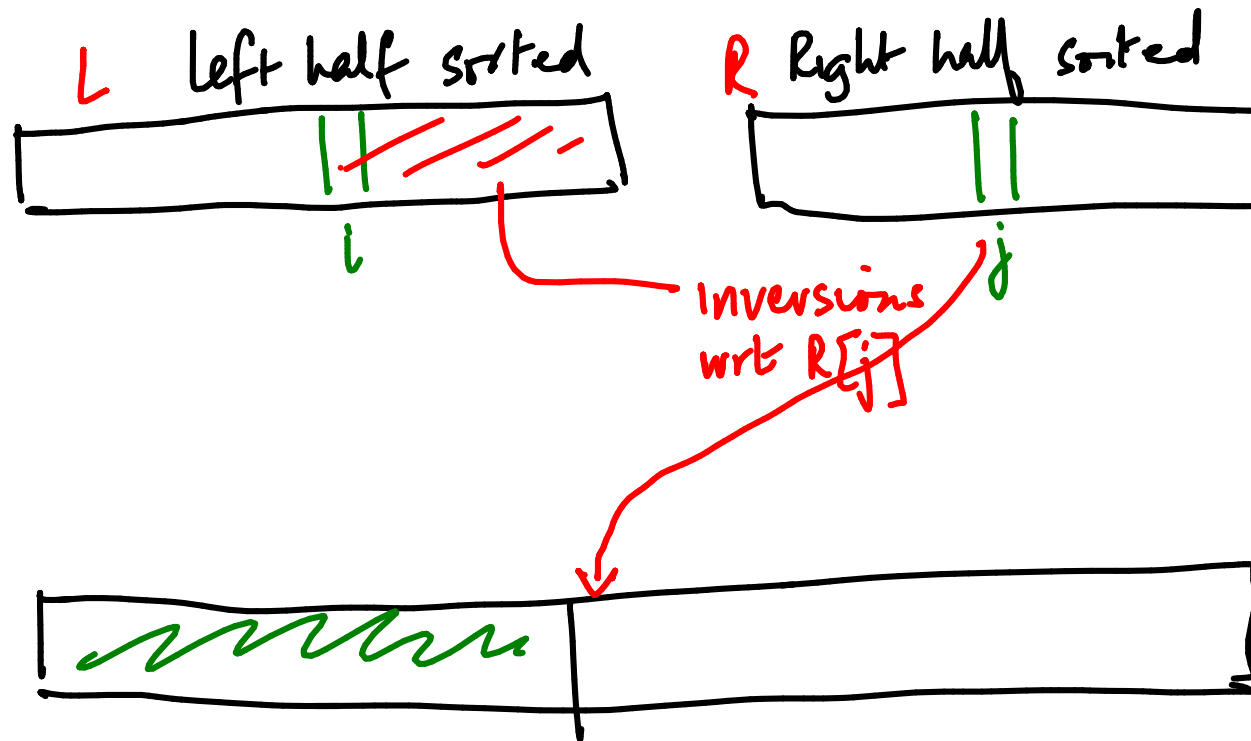
$i$  in left,  $j$  in right s.t.  $i > j$

Sort using merge sort

Combine step is merge

Each time we add to merged list from right list,

we witness as many inversions as current size of left list



Inductively, know inversion counts in L, R  
 While merging, update count each time we pick  
 an element from R — Overall  $O(n \log n)$

## Geometry

Given  $n$  2D points  $(x_1, y_1), (x_2, y_2) \dots (x_n, y_n)$ ,

find the closest pair (usual Euclidean distance  
 $\sqrt{(x_j - x_i)^2 + (y_j - y_i)^2}$ )

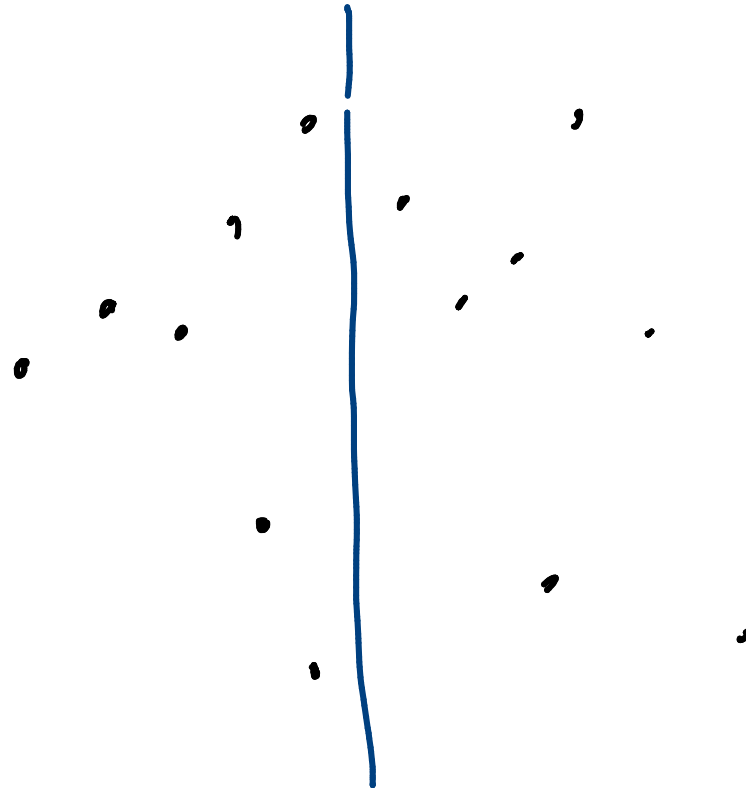
Trivial  $O(n^2)$  soln

- Try every pair of points

Divide & conquer?

Assume all  $x_i, y_j$  are distinct

Sort by  $x$ -coord & split



Inductively, know  
closest pair on  
left & right,  
respectively

Compute distances  
across separator

Sufficient to  
consider a "band"

Band:

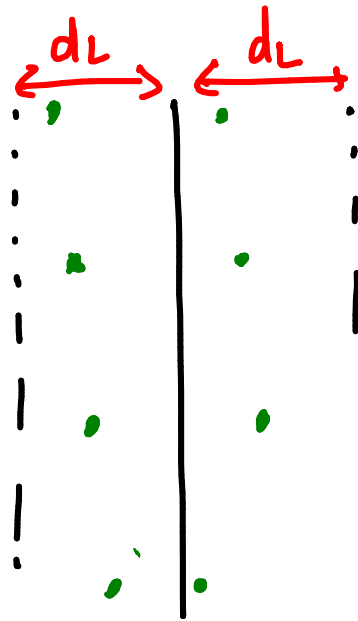
Let  $d_L, d_R$  be min distances on left, right

Wlog suppose  $d_L < d_R$

Need not consider points more than  $d_L$  away  
from separator

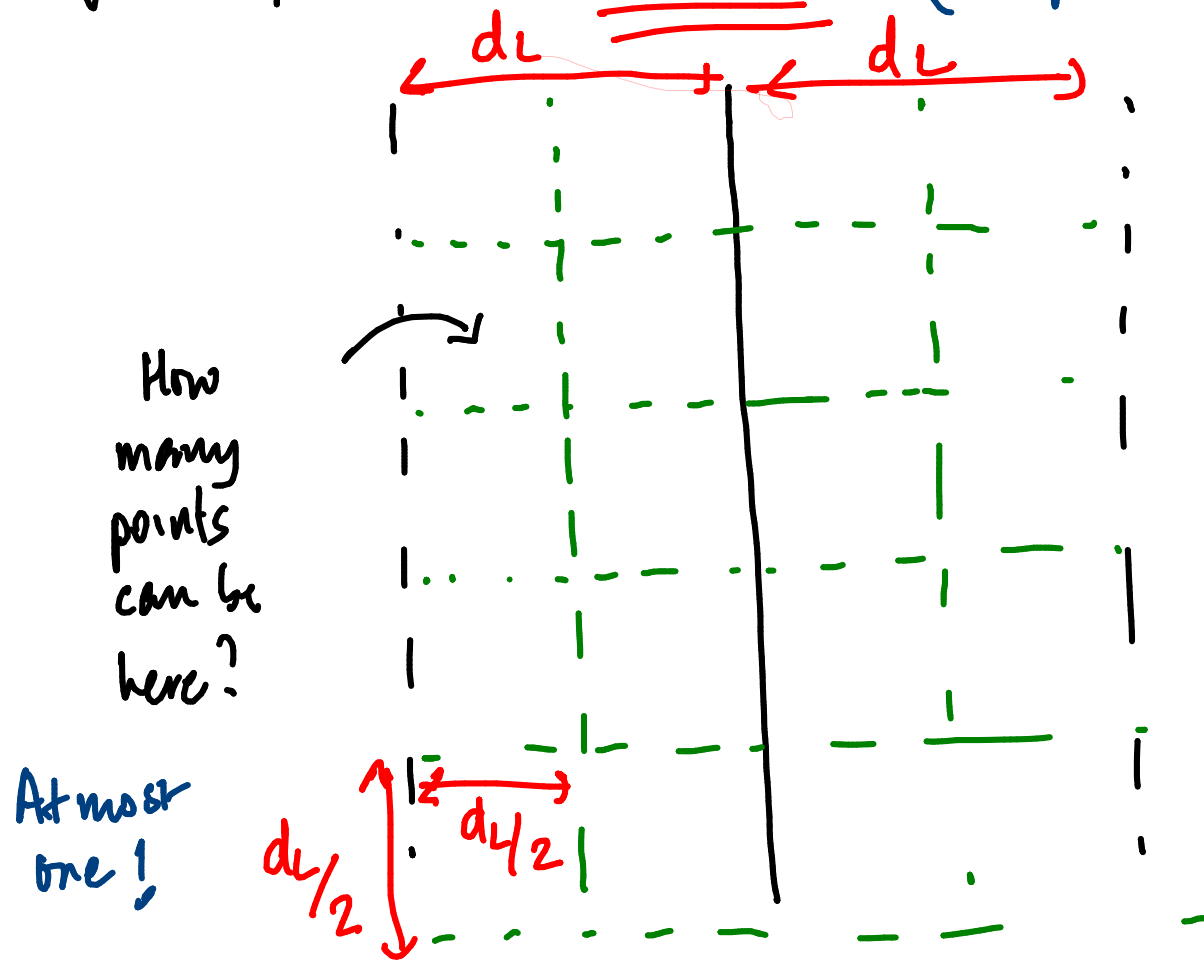
But...

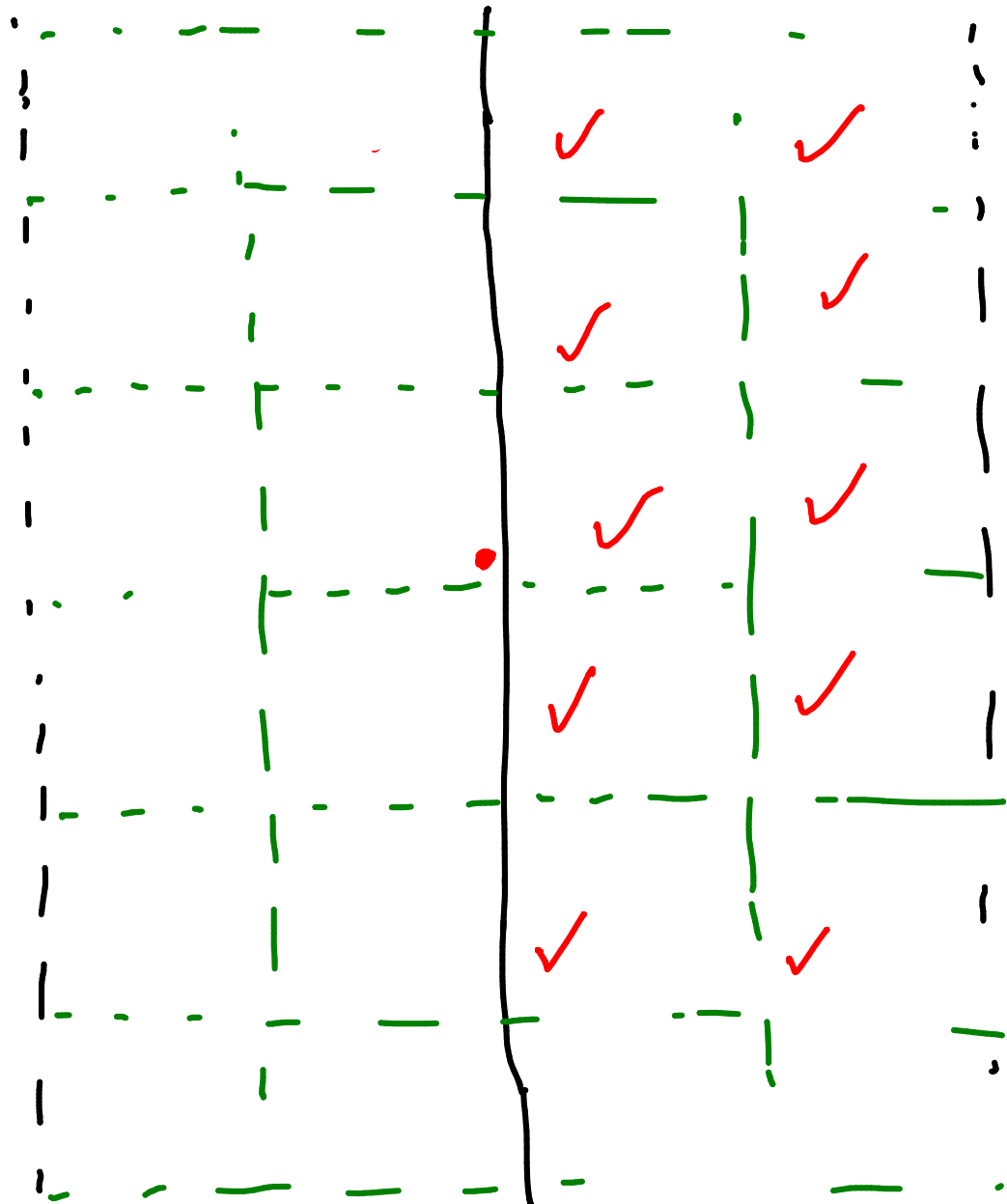
How do we  
ensure





Within band, limit comparisons made for a given point to a constant (independent of  $n$ )





Can bound  
the number  
of squares  
to check

Kleinberg &  
Tardos

Loose bound  
of 16

Merge step - process points in y-order

