

Lecture 5, 22 January 2026
Java: abstract classes, interfaces
Storage allocation

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Programming Language Concepts
January–April 2026

Grouping together classes

- Sometimes we collect together classes under a common heading
- Classes `Circle`, `Square` and `Rectangle` are all shapes
- Create a class `Shape` so that `Circle`, `Square` and `Rectangle` extend `Shape`
- We want to force every `Shape` to define a function
`public double perimeter()`
- Could define a function in `Shape` that returns an absurd value
`public double perimeter() { return(-1.0); }`
- Rely on the subclass to redefine this function
- What if this doesn't happen?
 - Should not depend on programmer discipline

*Simula
event queue*

Abstract classes

- A better solution

- Provide an **abstract definition** in **Shape**

```
public abstract double perimeter();
```



Abstract classes

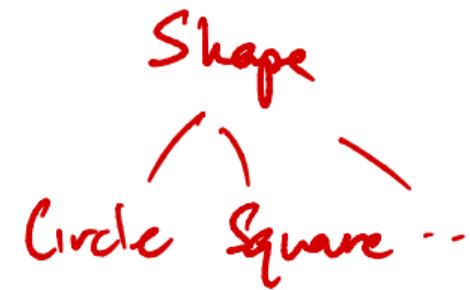
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Abstract classes

- A better solution
 - Provide an **abstract definition** in **Shape**
- Forces subclasses to provide a concrete implementation
- Cannot create objects from a class that has abstract functions
- **Shape** must itself be declared to be **abstract**

```
public abstract class Shape{  
    ...  
    public abstract double perimeter();  
    ...  
}
```

Abstract classes . . .

- Can still declare variables whose type is an abstract class

Abstract classes . . .

- Can still declare variables whose type is an abstract class

```
Shape shapearr[] = new Shape[3];
int sizearr[] = new int[3];

shapearr[0] = new Circle(...);
shapearr[1] = new Square(...);
shapearr[2] = new Rectangle(...);

for (i = 0; i < 3; i++){
    sizearr[i] = shapearr[i].perimeter();    ← Dynamic dispatch
    // each shapearr[i] calls the appropriate method
    ...
}
```

Generic functions

- Use abstract classes to specify generic properties

```
public abstract class Comparable{  
    public abstract int cmp(Comparable s);  
    // return -1 if this < s,  
    //          0 if this == 0,  
    //          +1 if this > s  
}
```

Generic functions

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public abstract class Comparable{  
    public abstract int cmp(Comparable s);  
    // return -1 if this < s,  
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}
```

- Now we can sort any array of objects that extend Comparable

```
public class SortFunctions{  
    public static void quicksort(Comparable[] a){  
        ...  
        // Usual code for quicksort, except that  
        // to compare a[i] and a[j] we use a[i].cmp(a[j])  
    }  
}
```

any subtype of Comparable

Generic functions ...

```
public class SortFunctions{  
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        ...  
    }  
}
```

Generic functions ... — “Structural polymorphism”

```
public class SortFunctions{  
    public static void quicksort(Comparable[] a){  
        ...  
    }  
}
```

- To use this definition of `quicksort`, we write

```
public class Myclass extends Comparable{  
    private double size; // quantity used for comparison  
  
    public int cmp(Comparable s){  
        if (s instanceof Myclass){  
            // compare this.size and ((Myclass) s).size  
            // Note the cast to access s.size  
        }  
    }  
}
```

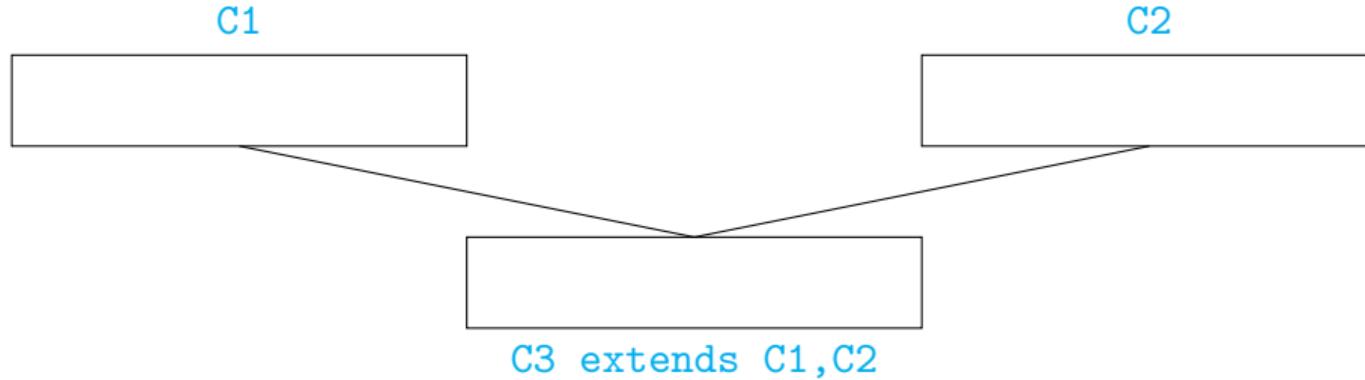


Multiple inheritance

- Can we sort `Circle` objects using the generic functions in `SortFunctions`?
 - `Circle` already extends `Shape`
 - Need `Circle` to also extend `Comparable`

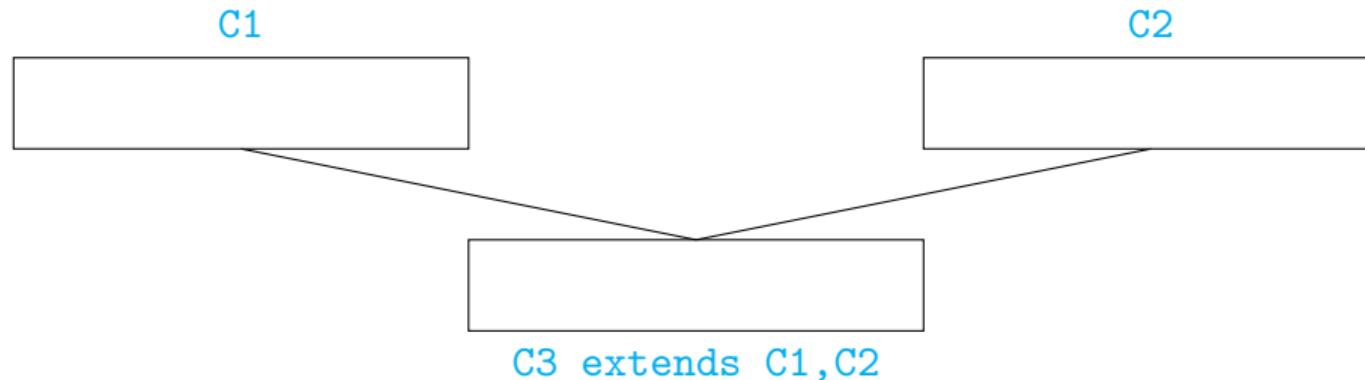
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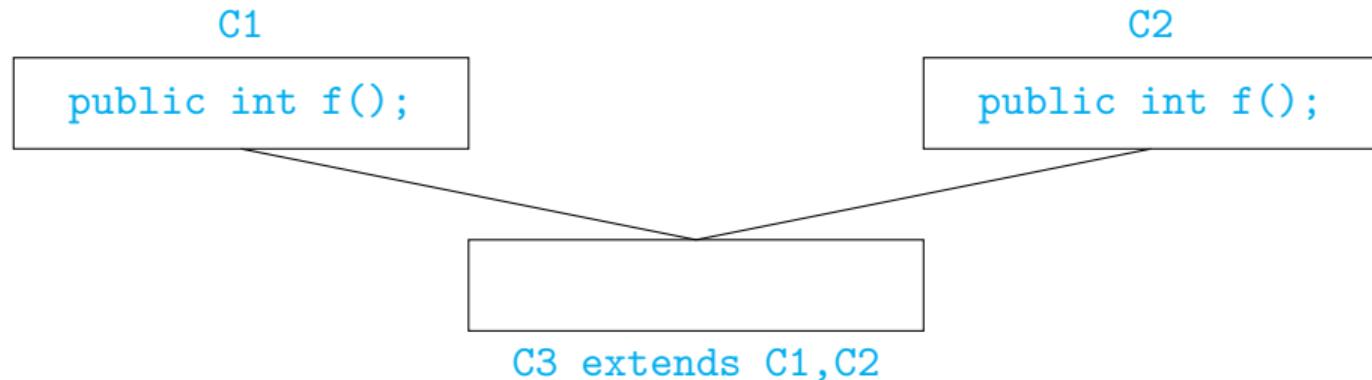
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Multiple inheritance

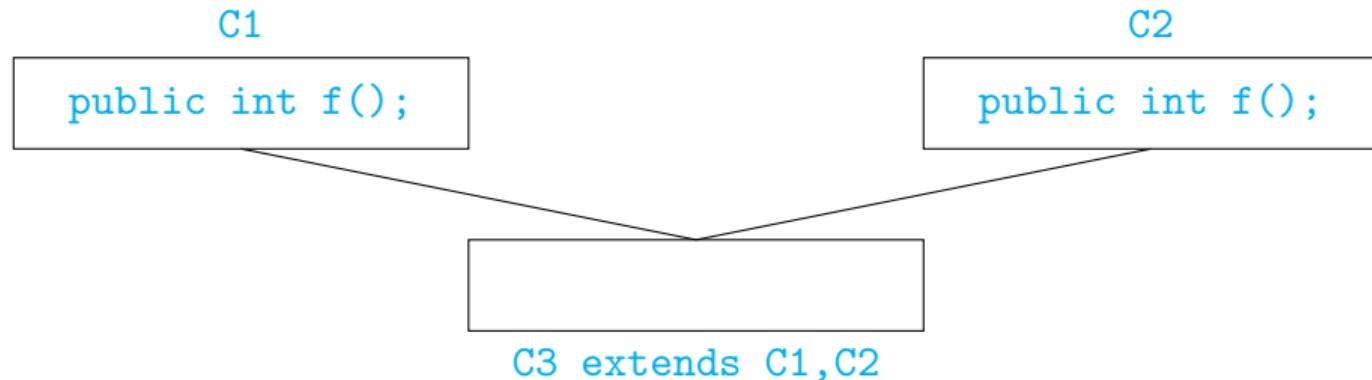
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- If `f()` is not overridden, which `f()` do we use in `C3`?

Multiple inheritance

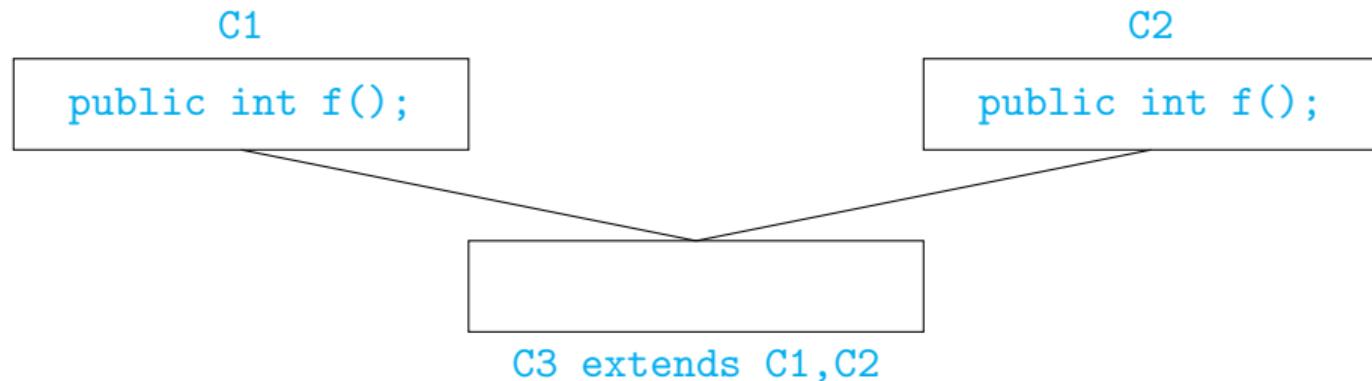
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Multiple inheritance

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- If `f()` is not overridden, which `f()` do we use in `C3`?
- Java does not allow multiple inheritance
- C++ allows this if `C1` and `C2` have no conflict

Interfaces and Multiple inheritance

- An **interface** is an abstract class with no concrete components

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- A class that extends an interface is said to **implement** it:

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public class Circle extends Shape implements Comparable{  
    public double perimeter(){...}  
    public int cmp(Comparable s){...}  
    ...  
}
```

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But . . .

Java now allows
concrete fns
in interfaces

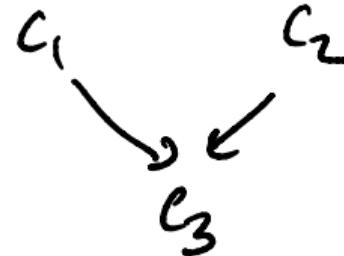
class A extends B implements
C,D,E {

- Can extend only one class, but can implement multiple interfaces
- Interfaces describe relevant aspects of a class
 - Abstract functions describe a specific “slice” of capabilities
 - Another class only needs to know about these capabilities

Java class hierarchy

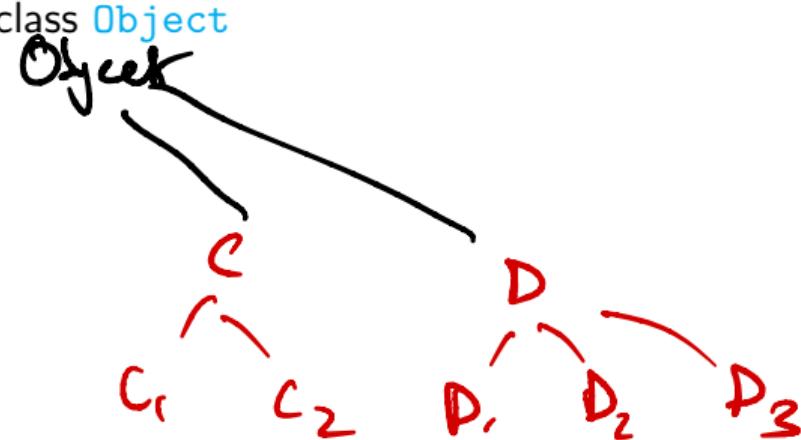
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Forest of classes



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- Useful methods defined in `Object`

```
public boolean equals(Object o) // defaults to pointer equality
public String toString()           // converts the values of the
                                    // instance variables to String
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- For Java objects `x` and `y`, `x == y` invokes `x.equals(y)`
- To print `o`, use `System.out.println(o+");`
 - Implicitly invokes `o.toString()`

Java class hierarchy

- Can exploit the tree structure to write generic functions
 - Example: search for an element in an array

```
public int find (Object[] objarr, Object o){  
    int i;  
    for (i = 0; i < objarr.length(); i++){  
        if (objarr[i] == o) {return i};  
    }  
    return (-1);  
}
```

Object.equals()

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    return (-1);  
}
```

objarr[i].equals(o)

- Recall that `==` is pointer equality, by default
- If a class overrides `equals()`, dynamic dispatch will use the redefined function instead of `Object.equals()` for `objarr[i] == o`

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            (this.year == d.year));  
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boolean equals(Date d)  
does not override  
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```

- Unfortunately,

`boolean equals(Date d)`
does not override
`boolean equals(Object o)`!

- Should write, instead

```
public boolean equals(Object d){  
    if (d instanceof Date){  
        Date myd = (Date) d;  
        return ((this.day == myd.day) &&  
                (this.month == myd.month)  
                (this.year == myd.year));  
    }  
    return(false);  
}
```

- Note the run-time type check and the cast

Overriding functions

- Overriding looks for “closest” match

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Object

Employee



Manager

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- Suppose we have `public boolean equals(Employee e)` but no `equals()` in `Manager`
- Consider

```
Manager m1 = new Manager(...);  
Manager m2 = new Manager(...);  
...  
if (m1.equals(m2)){ ... }
```

↳ subtype of Employee, then Object

Overriding functions

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- Suppose we have `public boolean equals(Employee e)` but no `equals()` in `Manager`

- Consider

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- `public boolean equals(Manager m)` is compatible with both `boolean equals(Employee e)` and `boolean equals(Object o)`
- Use `boolean equals(Employee e)`

Subclasses, subtyping and inheritance

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- Class hierarchy provides both **subtyping** and **inheritance**
- **Subtyping**
 - Capabilities of the subtype are a superset of the main type
 - If **B** is a subtype of **A**, wherever we require an object of type **A**, we can use an object of type **B**
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 - `Employee e = new Manager(...);` is legal
- **Inheritance**
 - Subtype can reuse code of the main type
 - **B** inherits from **A** if some functions for **B** are written in terms of functions of **A**
 - `Manager.bonus()` uses `Employee.bonus()`

Subtyping vs inheritance

- Consider the following example
 - queue, with methods `insert-rear`, `delete-front`
 - stack, with methods `insert-front`, `delete-front`
 - deque, with methods `insert-front`, `delete-front`, `insert-rear`, `delete-rear`

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 - `deque` has more functionality than `queue` or `stack`
 - `deque` is a subtype of both these types

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- What are the subtype and inheritance relationships between these classes?
- **Subtyping**
 - `deque` has more functionality than `queue` or `stack`
 - `deque` is a subtype of both these types
- **Inheritance**
 - Can suppress two functions in a `deque` and use it as a `queue` or `stack`
 - Both `queue` and `stack` inherit from `deque`

Subclasses, subtyping and inheritance

- Class hierarchy represents both **subtyping** and **inheritance**
- **Subtyping**
 - Compatibility of interfaces.
 - **B** is a subtype of **A** if every function that can be invoked on an object of type **A** can also be invoked on an object of type **B**.
- **Inheritance**
 - Reuse of implementations.
 - **B** inherits from **A** if some functions for **B** are written in terms of functions of **A**.
- Using one idea (hierarchy of classes) to implement both concepts blurs the distinction between the two

Variables, functions and storage

- Variables represent data residing in a memory location
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- Functions represent blocks of (reusable) code
 - Complexities introduced by **recursion**
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 - Need a way to keep track of all copies of a local **x**
 - Figure out which copy of **x** is referred to at any point of the execution

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- **Scope** and **lifetime** of variables

Scope

- Consider the following program block

```
{  
    int x = 2;  
    int y = 4;  
    {  
        int y = 3;  
        x = x+2; y = x+y;  
        print(x,y);  
    }  
    x = x+2; y = x+y;  
    print(x,y);  
}
```

```
for (int i=0; i<n; i++) {  
    --  
}
```

Scope

- Consider the following program block

```
{  
    int x = 2;  
    int y = 4;  
  
    {  
        int y = 3;          Outer y is hidden.  
        x = x+2; y = x+y;  Updated y value is not propagated outside  
        print(x,y);        4, 7  
    }  
  
    x = x+2; y = x+y;  
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```

Outer y is hidden.

Updated y value is not propagated outside
4, 7

Outer y value and updated x value
6, 10

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      ...
    }
  }
}
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```
✓{ int x = ...;
  ✓{ int y = ...;
    { int x = ...;
      X
      ...
    }
  }
}
```

- Scope of outer `x` is the two outer blocks

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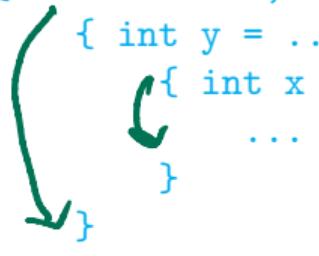


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{ int x = ...;  
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      ...  
    }  
  }  
}
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- Scope of outer `x` is the two outer blocks
- Scope of the inner `x` is the innermost block
- Lifetime of inner `x` is the time during which innermost block is active

Scope and Lifetime

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- Consider the example below

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- Scope of the inner **x** is the innermost block
- Lifetime of inner **x** is the time during which innermost block is active
- Lifetime of outer **x** is the time during which outermost block is active (includes the lifetime of inner **x**)

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- **static** variables are associated with a class as a whole
- Do not require instantiation of objects

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```
public class A {  
    static int howManyAs = 0;  
    int id;  
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        howManyAs += 1;  
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- The **static** variable **howManyAs** counts the number of instances of **A** created
- Lifetime of **howManyAs** spans the execution of the entire program
- Scope of **howManyAs** is limited to the class **A**

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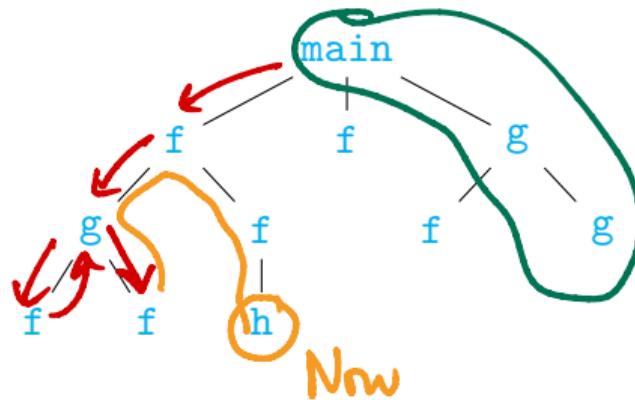
- For local variables and function parameters, we need to store one copy for each **function invocation (or activation)**
- **Activation record** — collection of all data related to a function invocation
- Includes space for local variables, parameters, intermediate results, and some pointers

Call graph

- A **call graph** helps us visualize the function calls during a program execution

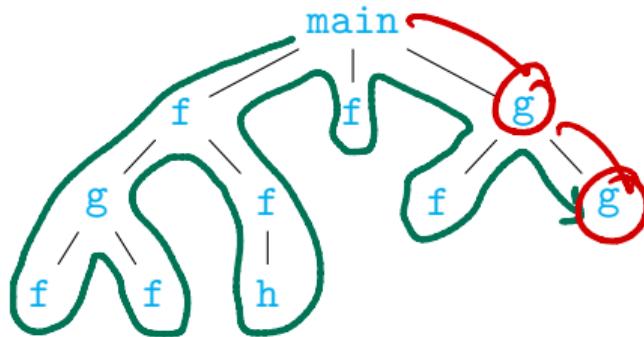
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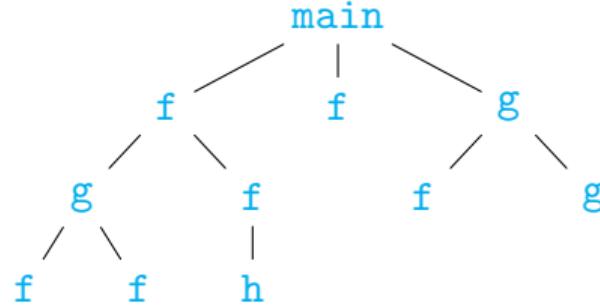
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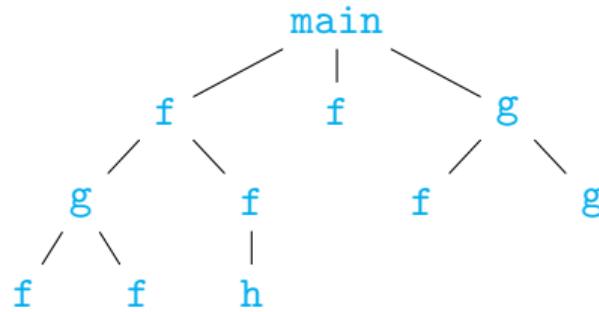
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- The set of **active function calls** at any point of time lies on the path from the root to the right most leaf
- If **f** calls **g**, then **g** is completed before **f**



Call graph

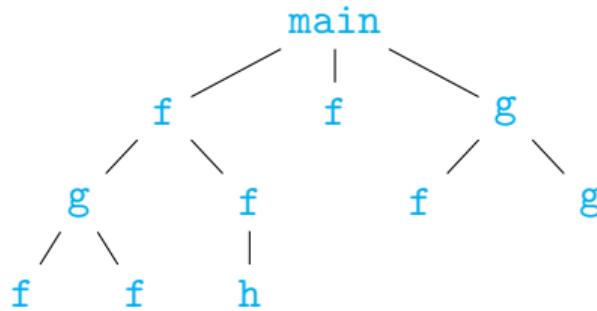
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- The set of **active function calls** at any point of time lies on the path from the root to the right most leaf
- If **f** calls **g**, then **g** is completed before **f**
- Store the activation records on a **stack**

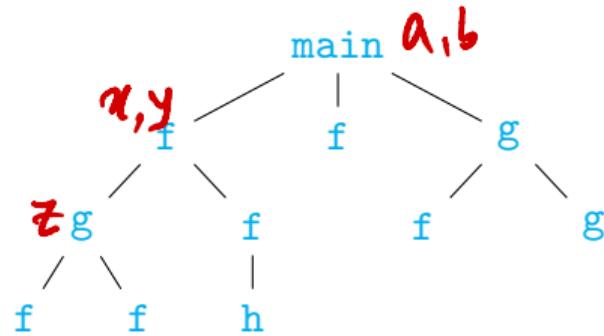
Call graph

- A **call graph** helps us visualize the function calls during a program execution



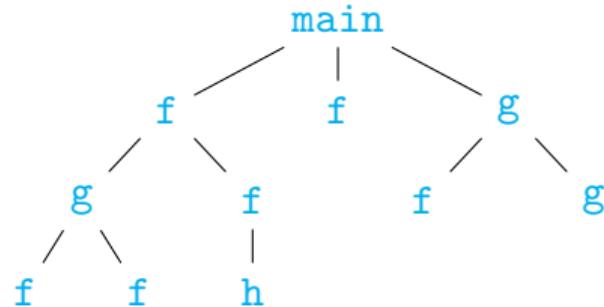
- The set of **active function calls** at any point of time lies on the path from the root to the right most leaf
- If **f** calls **g**, then **g** is completed before **f**
- Store the activation records on a **stack**
- Activation record is also called a **stack frame**

Activation records on stack



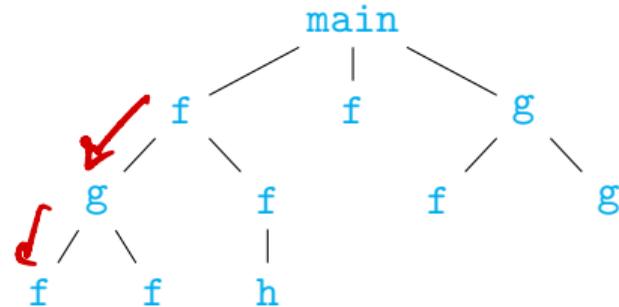
- Assume that `main` has local variables `a` and `b`, `f` has `x` and `y`, and `g` has `z`

Activation records on stack



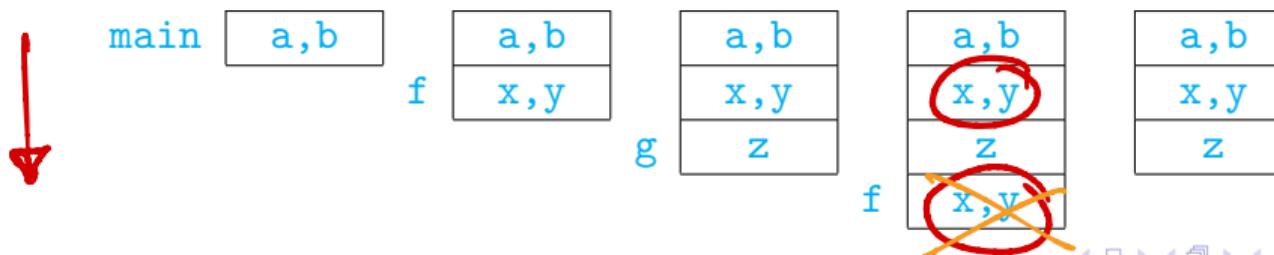
- Assume that `main` has local variables `a` and `b`, `f` has `x` and `y`, and `g` has `z`
- Place activation records on a stack — grows and shrinks as a program executes

Activation records on stack



- Assume that `main` has local variables `a` and `b`, `f` has `x` and `y`, and `g` has `z`
- Place activation records on a stack — grows and shrinks as a program executes

- The stack evolves as follows:



General layout of a program in memory

