

Abstraction, modularity, object-oriented programming

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Programming Language Concepts

Lecture 4, 3 February 2022

Stepwise refinement

- Begin with a high level description of the task

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begin  
  print first thousand prime numbers  
end
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Stepwise refinement

- Begin with a high level description of the task
- Refine the task into subtasks

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```

```
begin  
  declare table p  
  fill table p with first thousand primes  
  print table p  
end
```

Stepwise refinement

- Begin with a high level description of the task
- Refine the task into subtasks
- Further elaborate each subtask

```
begin
  print first thousand prime numbers
end
```

```
begin
✓ declare table p
  [fill table p with first thousand primes]
✓ print table p
end
```

```
begin
✓ integer array p[1:1000]
  [for k from 1 through 1000
    make p[k] equal to the kth prime number]
✓ for k from 1 through 1000
  print p[k]
```

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- Refine the task into subtasks
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- **Program refinement** — focus on code, not much change in data structures

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- Typical functions: `CreateAccount()`, `Deposit()/Withdraw()`, `PrintStatement()`

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- How do we represent each account?
 - Only need the current balance
 - Overall, an array of balances

Data refinement

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 - Typical functions: `CreateAccount()`, `Deposit()/Withdraw()`, `PrintStatement()`
- How do we represent each account?
 - Only need the current balance
 - Overall, an array of balances
- Refine `PrintStatement()` to include `PrintTransactions()`
 - Now we need to record transactions for each account
 - Data representation also changes
 - Cascading impact on other functions that operate on accounts

Modular software development

- Use refinement to divide the solution into **components**

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 - **Interfaces** — what is visible to other components, typically function calls
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- Simplest example of a component: a function
 - **Interfaces** — function header, arguments and return type
 - **Specification** — intended input-output behaviour
- Main challenge: suitable language to write specifications
 - Balance abstraction and detail, should not be another programming language!
 - Cannot algorithmically check that specification is met (halting problem!)

Z, VDM

Software Eng

Programming language support for abstraction

- Control abstraction
 - Functions and procedures
 - **Encapsulate** a block of code, reuse in different contexts

API

Programming language support for abstraction

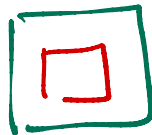
- Control abstraction

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- Data abstraction

- Abstract data types (ADTs)
- Set of values along with operations permitted on them
- Internal representation should not be accessible
- Interaction restricted to public interface
 - For example, when a stack is implemented as a list, we should not be able to observe or modify internal elements

Stack



push
pop
is-empty

Programming language support for abstraction

■ Control abstraction

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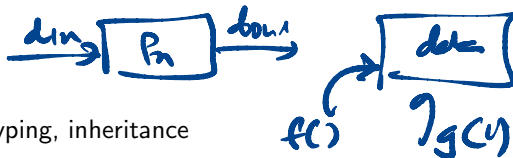
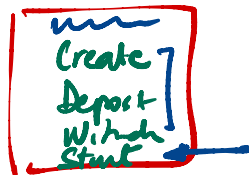
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■ Object-oriented programming

- Organize ADTs in a hierarchy
- Implicit reuse of implementations — subtyping, inheritance

Bank Account



20 data structures
1000 functions

} → Bug?

Inconsistency in
data struct.



20 data structures
50 functions per data structure

Examine all 1000
functions

Objects

- An **object** is like an abstract datatype
 - Hidden data with set of public operations
 - All interaction through operations — messages, **methods**, **member-functions**, ...



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- Uniform way of encapsulating different combinations of data and functionality
 - An object can hold single integer — e.g., a counter
 - An entire filesystem or database could be a single object
- Distinguishing features of object-oriented programming
 - Abstraction
 - Subtyping
 - Dynamic lookup
 - Inheritance

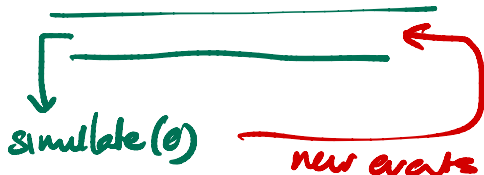
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Q := make-queue(first event)
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 - Queue must be well-typed, yet hold all types of events
 - Use a generic simulation operation across different types of events
 - Avoid elaborate checking of cases

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```

if e is of type 1 {--}

else if e is of type 2 {--}

...

e1 e2 e3

[s()] [s()] [s()]

Abstraction

- Objects are similar to abstract datatypes
 - Public interface
 - Private implementation
 - Changing the implementation should not affect interactions with the object

ADT \neq Data Structure
|
Abstract, functionality
Efficiency

Priority Queue
ADT

Heap
D.S.

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- Data-centric view of programming
 - Focus on what data we need to maintain and manipulate
- Recall that stepwise refinement could affect both code and data
 - Tying methods to data makes this easier to coordinate
 - Refining data representation naturally tied to updating methods that operate on the data

Subtyping

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- Arrange types in a hierarchy
 - A **subtype** is a specialization of a type
 - If **A** is a subtype of **B**, wherever an object of type **B** is needed, an object of type **A** can be used
 - Every object of type **A** is also an object of type **B**
 - Think subset — if $X \subseteq Y$, every $x \in X$ is also in Y

suggest less capabilities



Subtyping

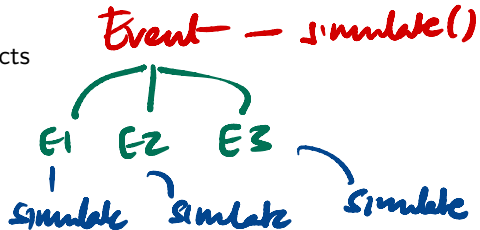
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- If **f()** is a method in **B** and **A** is a subtype of **B**, every object of **A** also supports **f()**
 - Implementation of **f()** can be different in **A**

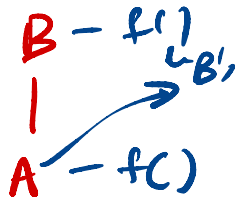


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- Dynamic lookup
 - A variable `v` of type `B` can refer to an object of subtype `A`
 - Static type of `v` is `B`, but method implementation depends on **run-time** type `A`

Inheritance

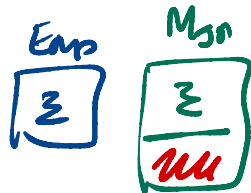
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 - Additional fields and functions: date of promotion, seniority (in current role)



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- Usually one hierarchy of types to capture both subtyping and inheritance
 - **A** can inherit from **B** iff **A** is a subtype of **B**

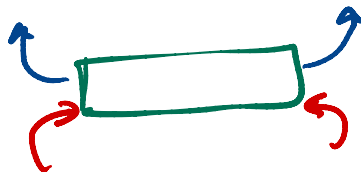
B
|
A

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- Usually one hierarchy of types to capture both subtyping and inheritance
 - **A** can inherit from **B** iff **A** is a subtype of **B**
- Philosophically, however the two are different
 - Subtyping is a relationship of interfaces
 - Inheritance is a relationship of implementations

Subtyping vs inheritance

- A **deque** is a double-ended queue
 - Supports `insert-front()`, `delete-front()`, `insert-rear()` and `delete-rear()`



Subtyping vs inheritance

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- We can implement a stack or a queue using a deque
 - Stack: use only `insert-front()`, `delete-front()`,
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- **Stack** and **Queue** inherit from **Deque** — reuse implementation
- But **Stack** and **Queue** are not subtypes of **Deque**
 - If `v` of type **Deque** points an object of type **Stack**, cannot invoke `insert-rear()`, `delete-rear()`
 - Similarly, no `insert-front()`, `delete-rear()` in **Queue**

B — `fc()`
|
A — `fc()`

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 - Supports `insert-front()`, `delete-front()`, `insert-rear()` and `delete-rear()`
- We can implement a stack or a queue using a deque
 - Stack: use only `insert-front()`, `delete-front()`,
 - Queue: use only `insert-rear()`, `delete-front()`,
- **Stack** and **Queue** inherit from **Deque** — reuse implementation
- But **Stack** and **Queue** are not subtypes of **Deque**
 - If `v` of type **Deque** points an object of type **Stack**, cannot invoke `insert-rear()`, `delete-rear()`
 - Similarly, no `insert-front()`, `delete-rear()` in **Queue**
- Interfaces of **Stack** and **Queue** are not compatible with **Deque**
 - In fact, **Deque** is a subtype of both **Stack** and **Queue**

Summary

- Solving a complex task requires breaking it down into manageable components
 - **Top down**: refine the task into subtasks; **Bottom up**: combine simple building blocks
- Modular description of components — interface and specification
 - Build prototype implementation to validate design
 - Reimplement the components independently, preserving interface and specification
- PL support for abstraction
 - Control flow: functions and procedures
 - Data: Abstract data types, object-oriented programming
- Distinguishing features of object-oriented programming
 - **Abstraction**: Public interface, private implementation, like ADTs
 - **Subtyping**: Hierarchy of types, compatibility of interfaces
 - **Dynamic lookup**: Choice of method implementation is determined at run-time
 - **Inheritance**: Reuse of implementations