### Lecture 2: 10 January, 2023

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## Market-Basket Analysis

- Set of items  $I = \{i_1, i_2, ..., i_N\}$
- Set of transactions  $T = \{t_1, t_2, \dots, t_M\}$ 
  - A transaction is a set  $t \subseteq I$  of items
- Identify association rules  $X \rightarrow Y$ 
  - $X, Y \subseteq I, X \cap Y = \emptyset$
  - If  $X \subseteq t_j$  then it is likely that  $Y \subseteq t_j$
- Two thresholds
  - How frequently does  $X \subseteq t_j$  imply  $Y \subseteq t_j$ ?
  - How significant is this pattern overall?

## Setting thresholds

- For  $Z \subseteq I$ , Z.count =  $|\{t_j \mid Z \subseteq t_j\}|$
- How frequently does  $X \subseteq t_j$  imply  $Y \subseteq t_j$ ?
  - Fix a confidence level  $\chi$
  - Want  $\frac{(X \cup Y).count}{X.count} \ge \chi$
- How significant is this pattern overall?
  - Fix a support level  $\sigma$

• Want 
$$\frac{(X \cup Y).count}{M} \ge \sigma$$

Given sets of items *I* and transactions *T*, with confidence χ and support σ, find all valid association rules X → Y

### Frequent itemsets

- $X \to Y$  is interesting only if  $(X \cup Y)$ .count  $\geq \sigma \cdot M$
- First identify all frequent itemsets

•  $Z \subseteq I$  such that Z.count  $\geq \sigma \cdot M$ 

Naïve strategy: maintain a counter for each Z

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■ For each t_j \in T
For each Z \subseteq t_j
Increment the counter for Z
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- After scanning all transactions, keep Z with Z.count  $\geq \sigma \cdot M$
- Need to maintain 2<sup>|/|</sup> counters
  - Infeasible amount of memory
  - Can we do better?

## Sample calculation

• Let's assume a bound on each  $t_i \in T$ 

• No transaction has more than 10 items

• Say  $N = |I| = 10^6$ ,  $M = |T| = 10^9$ ,  $\sigma = 0.01$ 

• Number of possible subsets to count is  $\sum_{i=1}^{10} {10^6 \choose i}$ 

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A singleton subset {x} that is frequent is an item x that appears in at least 10<sup>7</sup> transactions

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• Number of possible subsets to count is  $\sum_{i=1}^{10} {10^6 \choose i}$ 

- A singleton subset {x} that is frequent is an item x that appears in at least 10<sup>7</sup> transactions
- Totally, T contains at most  $10^{10}$  items
- At most  $10^{10}/10^7 = 1000$  items are frequent!
- How can we exploit this?

• Clearly, if Z is frequent, so is every subset  $Y \subseteq Z$ 

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- We exploit the contrapositive

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Apriori observation

If Z is not a frequent itemset, no superset Y \supseteq Z can be

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If Z is not a frequent itemset, no superset Y \supseteq Z can be
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- For instance, in our earlier example, every frequent itemset must be built from the 1000 frequent items
- In particular, for any frequent pair {x, y}, both {x} and {y} must be frequent
- Build frequent itemsets bottom up, size 1,2,...

•  $F_i$ : frequent itemsets of size i — Level i

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- $F_i$ : frequent itemsets of size i Level i
- $F_1$ : Scan T, maintain a counter for each  $x \in I$



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- $F_2$ : Scan T, maintain a counter for each  $X \in C_2$
- $C_3 = \{\{x, y, z\} \mid \{x, y\}, \{x, z\}, \{y, z\} \in F_2\}$
- $F_3$ : Scan T, maintain a counter for each  $X \in C_3$

| $T=G \longrightarrow F_1$ |  |
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- $F_i$ : frequent itemsets of size i Level i
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- $C_k$  = subsets of size k, every (k-1)-subset is in  $F_{k-1}$
- $F_k$ : Scan T, maintain a counter for each  $X \in C_k$

- $C_k$  = subsets of size k, every (k-1)-subset is in  $F_{k-1}$
- How do we generate  $C_k$ ?

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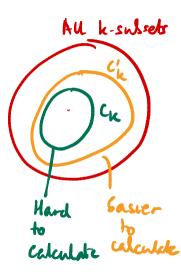
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Expensive!

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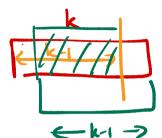
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- Items are ordered:  $i_1 < i_2 < \cdots < i_N$
- List each itemset in ascending order canonical representation

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- How do we generate  $C_k$ ?
- Naïve: enumerate subsets of size k and check each one
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- Observation: Any  $C'_k \supseteq C_k$  will do as a candidate set
- Items are ordered:  $i_1 < i_2 < \cdots < i_N$
- List each itemset in ascending order canonical representation
- Merge two (k-1)-subsets if they differ in last element

$$X = \{i_1, i_2, \dots, i_{k-2}, i_{k-1}\}$$
  

$$X' = \{i_1, i_2, \dots, i_{k-2}, i'_{k-1}\}$$
  

$$Merge(X, X') = \{i_1, i_2, \dots, i_{k-2}, i_{k-1}, i'_{k-1}\}$$



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## • Merge $(X, X') = \{i_1, i_2, \dots, i_{k-2}, i_{k-1}, i'_{k-1}\}$ • $X = \{i_1, i_2, \dots, i_{k-2}, i_{k-1}\}$ • $X' = \{i_1, i_2, \dots, i_{k-2}, i'_{k-1}\}$

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•  $X' = \{i_1, i_2, \dots, i_{k-2}, i'_{k-1}\}$ 

•  $C'_{k} = \{ Merge(X, X') \mid X, X' \in F_{k-1} \}$ 

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  - $X' = \{i_1, i_2, \dots, i_{k-2}, i'_{k-1}\}$
- $C'_k = \{ \operatorname{Merge}(X, X') \mid X, X' \in F_{k-1} \}$
- Claim  $C_k \subseteq C'_k$ 
  - Suppose  $Y = \{i_1, i_2, \dots, i_{k-1}, i_k\} \in C_k$
  - $X = \{i_1, i_2, \dots, i_{k-2}, i_{k-1}\} \in F_{k-1}$  and  $X' = \{i_1, i_2, \dots, i_{k-2}, i_k\} \in F_{k-1}$
  - $Y = Merge(X, X') \in C'_k$



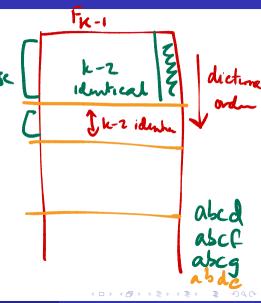
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- $C'_k = \{ \operatorname{Merge}(X, X') \mid X, X' \in F_{k-1} \}$

• Claim  $C_k \subseteq C'_k$ 

- Suppose  $Y = \{i_1, i_2, ..., i_{k-1}, i_k\} \in C_k$
- $X = \{i_1, i_2, \dots, i_{k-2}, i_{k-1}\} \in F_{k-1}$  and  $X' = \{i_1, i_2, \dots, i_{k-2}, i_k\} \in F_{k-1}$
- $Y = Merge(X, X') \in C'_k$
- Can generate  $C'_k$  efficiently
  - Arrange  $F_{k-1}$  in dictionary order
  - Split into blocks that differ on last element
  - Merge all pairs within each block



- $C_1 = \{\{x\} \mid x \in I\}$
- $F_1 = \{Z \mid Z \in C_1, Z. \text{count} \geq \sigma \cdot M\}$
- For  $k \in \{2, 3, ...\}$ 
  - $C'_k = \{ \operatorname{Merge}(X, X') \mid X, X' \in F_{k-1} \}$
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- When do we stop?

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- When do we stop?
- k exceeds the size of the largest transaction

= Fk is empty More generally, Ck is mply

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(x, y)

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- When do we stop?
- k exceeds the size of the largest transaction
- $F_k$  is empty

Next step: From frequent itemsets to association rules

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- Given sets of items / and transactions T. with confidence  $\chi$  and support  $\sigma$ , find all valid association rules  $X \to Y$ 
  - $X, Y \subseteq I, X \cap Y = \emptyset$ •  $\frac{(X \cup Y).count}{X.count} \ge \chi$  $\blacksquare \ \frac{(X \cup Y).count}{M} \ge \sigma$

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- Given sets of items *I* and transactions *T*, with confidence χ and support σ, find all valid association rules X → Y
  - $X, Y \subseteq I, X \cap Y = \emptyset$ •  $\frac{(X \cup Y).count}{X.count} \ge \chi$ •  $\frac{(X \cup Y).count}{M} \ge \sigma$
- For a rule X → Y to be valid, X ∪ Y should be a frequent itemset
- Apriori algorithm finds all  $Z \subseteq I$  such that Z.count  $\geq \sigma \cdot M$

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#### Naïve strategy

- For every frequent itemset Z
  - Enumerate all pairs  $X, Y \subseteq Z, X \cap Y = \emptyset$

• Check  $\frac{(X \cup Y).count}{X.count} \ge \chi$ 

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Can we do better?

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Can we do better?

- Sufficient to check all partitions of Z
  - Suppose X, Y ⊆ Z, X → Y is a valid association rule, but X ∪ Y is a proper subset of Z

### Naïve strategy

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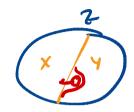
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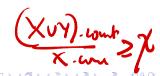
- Sufficient to check all partitions of Z
  - Suppose X, Y ⊆ Z, X → Y is a valid association rule, but X ∪ Y is a proper subset of Z
  - $\bullet X \cup Y = Z' \subsetneq Z$
  - Z' is also a frequent itemset (a priori)
  - X, Y partitions Z'

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- Sufficient to check all partitions of Z
- Suppose  $Z = X \uplus Y$ ,  $X \to Y$  is a valid rule and  $y \in Y$
- What about  $(X \cup \{y\}) \to Y \setminus \{y\}$ ?



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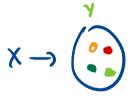
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- Suppose  $Z = X \uplus Y$ ,  $X \to Y$  is a valid rule and  $y \in Y$
- What about  $(X \cup \{y\}) \to Y \setminus \{y\}$ ?

• Know 
$$\frac{(X \cup Y).count}{X.count} \ge \chi$$
  
• Check  $\frac{(X \cup Y).count}{(X \cup \{y\}).count} \ge \chi$ 

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- Sufficient to check all partitions of Z
- Suppose  $Z = X \uplus Y$ ,  $X \to Y$  is a valid rule and  $y \in Y$
- What about  $(X \cup \{y\}) \to Y \setminus \{y\}$ ?
  - Know  $\frac{(X \cup Y).count}{X.count} \ge \chi$ • Check  $\frac{(X \cup Y).count}{(X \cup \{y\}).count} \ge \chi$
  - X.count  $\geq (X \cup \{y\})$ .count, always
  - Second fraction has smaller denominator, so  $(X \cup \{y\}) \rightarrow Y \setminus \{y\}$  is also a valid rule



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Observation: Can use apriori principle again!

## Apriori for association rules

- If  $X \to Y$  is a valid rule, and  $y \in Y$ ,  $(X \cup \{y\}) \to Y \setminus \{y\}$  must also be a valid rule
- If  $X \to Y$  is not a valid rule, and  $x \in X$ ,  $(X \setminus \{x\}) \to Y \cup \{x\}$  cannot be a valid rule

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## Apriori for association rules

- If  $X \to Y$  is a valid rule, and  $v \in Y$ .  $(X \cup \{y\}) \rightarrow Y \setminus \{y\}$  must also be a valid rule
- If  $X \to Y$  is not a valid rule, and  $x \in X$ .  $(X \setminus \{x\}) \to Y \cup \{x\}$  cannot be a valid rule
- Start by checking rules with single element on the right

 $\blacksquare Z \setminus \{z\} \to \{z\}$ 

- For  $X \to \{x, y\}$  to be a valid rule, both  $(X \cup \{x\}) \rightarrow \{y\}$  and  $(X \cup \{y\}) \rightarrow \{x\}$  must be valid
- Explore partitions of each frequent itemset "level by level"

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Classify documents by topic

• Consider the table on the right

| Words in document           | Topic     |
|-----------------------------|-----------|
| student, teach, school      | Education |
| student, school             | Education |
| teach, school, city, game   | Education |
| cricket, football           | Sports    |
| football, player, spectator | Sports    |
| cricket, coach, game, team  | Sports    |
| football, team, city, game  | Sports    |

- Classify documents by topic
- Consider the table on the right
- Items are regular words and topics
- Documents are transactions set of words and one topic

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- Classify documents by topic
- Consider the table on the right
- Items are regular words and topics
- Documents are transactions set of words and one topic
- Look for association rules of a special form
  - {student, school}  $\rightarrow$  {Education}
  - {game, team}  $\rightarrow$  {Sports}

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- Right hand side always a single topic
- Class Association Rules

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- Market-basket analysis searches for correlated items across transactions
- Formalized as association rules
- Apriori principle helps us to efficiently
  - identify frequent itemsets, and
  - split these itemsets into valid rules
- Class association rules simple supervised learning model