#### Data Mining and Machine Learning

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### Information retrieval on the Internet

- Traditional IR
  - Books published after editing, review trustworthy content
- IR for Internet
  - Internet documents are self-published, unverified
  - Economic incentive to boost rankings through fraudulent means
  - Ranking algorithms should try not to be fooled
- Easy to add invisible content in HTML to misdirect search
  - Merging text and background colour, overlay text with images, unreadable font size
- Self published documents may omit useful search terms
  - IBM webpage did not mention the word "computer"

## Exploiting hypertext

- Hypertext links refer from one document to another
  - > <a href="https://www.cmi.ac.in"> CMI webpage </a>
  - Target location : https://www.cmi.ac.in
  - Anchor text : CMI webpage
- Use anchor text to index document at target location
  - Reliable indicator of what target document is about
- Hyperlinks also connect internet documents as a directed graph
  - Reason about the World Wide Web (WWW) as a gigantic graph
  - Use techniques from social network analysis

### Social network analysis — prestige

- Consider the film industry
  - When is an actor a star? When is a director famous?
  - Stars are sought out by famous directors
  - Famous directors get stars to work in their films
  - Recursive definition
- Network (graph) of actors and directors, matrix M



### Social network analysis — prestige

- Each actor *i* has star value *S*[*i*]
- Each director j has fame F[j]
- Actors derive star value from the famous directors they work with

$$S[i] = \sum_{j} M[i, j] \cdot F[j], \text{ or } S = M \cdot F$$

- Directors derive fame from the stars who work with them  $F[j] = \sum_{i} M[i, j] \cdot S[i], \text{ or } F = M^{\top} \cdot S$
- Substituting F from second equation,  $S = M \cdot M^{\top} \cdot S$
- Substituting S from first equation,  $F = M^{\top} \cdot M \cdot F$
- Solve for S, F to compute star ratings, fame

## Prestige for webpages

- Each document *i* has prestige *P*[*i*]
- Prestigious (reliable) documents confer prestige on documents they link to
  - P[i] is shared equally among all outgoing links
- A document derives prestige from documents that link to it
  - P[i] is sum of prestige transferred by incoming links
- Structure of the internet, adjacency matrix A

Webpagesj
$$i$$
 $\vdots$  $\dots$  $\dots$ 

Prestige for webpages ...

• Suppose 
$$A = \begin{bmatrix} 0 & 1 & 1 \\ 1 & 0 & 0 \\ 1 & 1 & 0 \end{bmatrix}$$

• Each document initially has prestige 1,  $P = \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix}$ 

• If a webpage points to n other pages, each of them gets 1/n of P[i]

• Prestige transfer matrix, 
$$A^* = \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \\ 1 & 0 & 0 \\ \frac{1}{2} & \frac{1}{2} & 0 \end{bmatrix}$$
  
• One step:  $P^{\top} \cdot A^* = \begin{bmatrix} 1 & 1 & 1 \end{bmatrix} \cdot \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \\ 1 & 0 & 0 \\ \frac{1}{2} & \frac{1}{2} & 0 \end{bmatrix} = \begin{bmatrix} 1.5 & 1 & 0.5 \end{bmatrix}$ 

### Page rank

- Stable solution:  $P^{\top} \cdot A^* = P^{\top}$
- *P*[*i*] is Page rank of webpage *i* 
  - Larry Page, co-founder of Google with Sergey Brin
- How do we compute  $P^{\top}$ ?
- A\* is a stochastic matrix each row sums to 1

$$orall i \sum_j A^*[i,j] = 1$$

- Intepret A\*[i, j] as probability of moving from document i to document j — random web surfer
- Use theory of Markov chains

#### Markov chains

- Finite set of states, with transition probabilities between states
- For us, states are documents
  - Henceforth, write A\* as A for convenience



• P[j] is probability of being in document j

• Start in document 1, so initially 
$$P = \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}$$

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#### Markov chains ...

• After one step: 
$$P^{\top}A = \begin{bmatrix} 1 & 0 & 0 \end{bmatrix} \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \\ 1 & 0 & 0 \\ \frac{1}{2} & \frac{1}{2} & 0 \end{bmatrix} = \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \end{bmatrix}$$
  
• After second step:  $\begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \\ 1 & \frac{1}{2} & 0 \end{bmatrix} \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \\ 1 & 0 & 0 \\ \frac{1}{2} & \frac{1}{2} & 0 \end{bmatrix} = \begin{bmatrix} \frac{3}{4} & \frac{1}{4} & 0 \end{bmatrix}$ 

- After k steps, P[j] is probability of being in state j
- Continuing our example,

 $\left[\begin{array}{ccc} \frac{3}{4} & \frac{1}{4} & 0 \end{array}\right] \rightarrow \left[\begin{array}{ccc} \frac{1}{4} & \frac{3}{8} & \frac{3}{8} \end{array}\right] \rightarrow \left[\begin{array}{ccc} \frac{9}{16} & \frac{5}{16} & \frac{1}{8} \end{array}\right]$ 

• Is it the case that P[j] > 0 for all j continuously, after some point?

# Ergodicity

- Markov chain A is ergodic if there is some  $t_0$  such that for every P, for all  $t > t_0$ , for every j,  $(P^T A^t)[j] > 0$ .
  - No matter where we start, after t > t<sub>0</sub> steps, every state has a nonzero probability of being visited in step t
- Properties of ergodic Markov chains
  - There is a stationary distribution  $\pi$  such that  $\pi^{\top} A = \pi^{\top}$

★  $\pi^{\top}$  is a left eigenvector of A

• For any starting distribution *P*,  $\lim_{t\to\infty} P^{\top} A^t = \pi^{\top}$ 

# Ergodicity ...

- How can ergodicity fail?
  - Starting from *i*, we reach a set of states from which there is no path back to *i*
  - We have a cycle i → j → k → i → j → k ···, so we can only visit some states periodically
- Sufficient conditions for ergodicity
  - Irreducibility: When viewed as a directed graph, A is strongly connected
    - \* For all states i, j, there is a path from i to j and a path from j to i
  - Aperiodicity: For any pair of vertices *i*, *j*, the gcd of the lengths of all paths from *i* to *j* is 1
    - ★ In particular, paths (loops) from i to i do not all have lengths that are multiples of some k ≥ 2
    - ★ Prevents bad cycles

## Making the web graph ergodic

- No reason why web graph is irreducible and aperiodic
- Web graph has dead ends terminal documents, no outgoing links
- Solution: Add random jumps between documents teleportation
- Teleportation matrix T: For all i, j, T[i, j] = 1/N, where N is the total number of documents
  - The random surfer ignores all the links in the current document and types a new URL
- Let  $\alpha$  be the probability of teleportation:  $M = \alpha T + (1 \alpha)A$ 
  - Check that *M* is stochastic
- By construction,
  - ► *M* is strongly connected direct edge between each pair of documents
  - M is aperiodic paths of any length exist between i and j
  - M has no dead ends

## Page Rank

- In the modified web graph, stationary distribution is the Page rank,  $\pi^T M = \pi^T$
- Compute using  $\lim_{t\to\infty} P^T M^t$
- Use recursive doubling to accelerate computation of  $\lim_{t\to\infty} P^T M^t$ 
  - Compute  $M, M^2, (M^2)^2 = M^4, \ldots, (M^{2i})^2 = M^{4i}, \ldots$
  - Set a threshold for progress to stop the process
- Some limitations of Page rank
  - Universal property of a webpage, independent of a query
  - Define a topic-sensitive page rank
- Page rank was one the keys to the initial success of Google
  - Constant tweaks to ranking algorithm to keep ahead of search engine optimizers (SEO)

# Summary

- IR on webpages presents new challenges because document content is unreliable
- Hypertext tags can provide better indexing terms
- Hypertext links create a graph of webpages
- Apply techniques from social network analysis, Markov chains
- Page rank computes the prestige of a documents using the graph structure of webpages