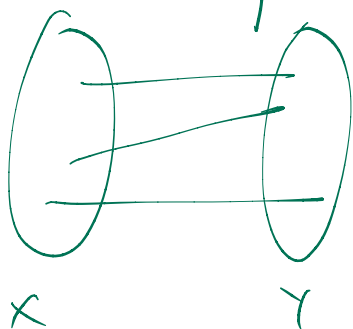


• Bipartite graphs:

$G = (V, E)$ is bipartite if $V = X \cup Y$
and $E \subseteq (X \rightarrow Y)$

(ie) there is a bipartition of vertices



and edges only
go from X to Y.

• Characterization of bipartite graphs:

Prop

G is bipartite iff G has no odd cycle.

Pf:

• Can do this using BFS.

HAMILTONIAN CYCLES:

- In general hard to find algorithmically in polynomial time.

Thm: G be connected. If $\deg(v_i) \geq |V|/2$
Then G has a H. Cycle;

Pf:

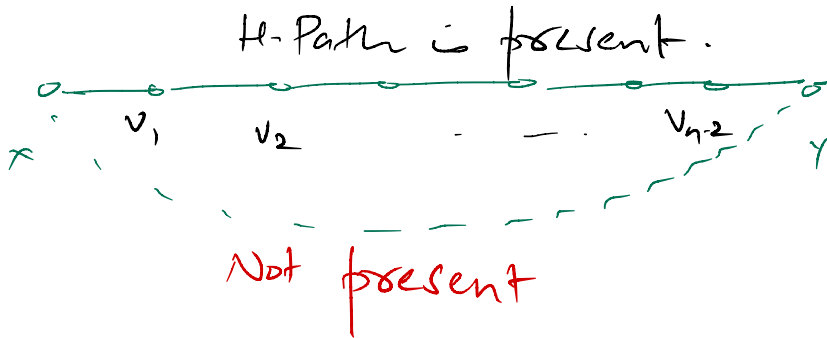
Take a maximal graph (maximal in terms of edges) satisfying the hypothesis and not the conclusion

Show that is a contradiction;

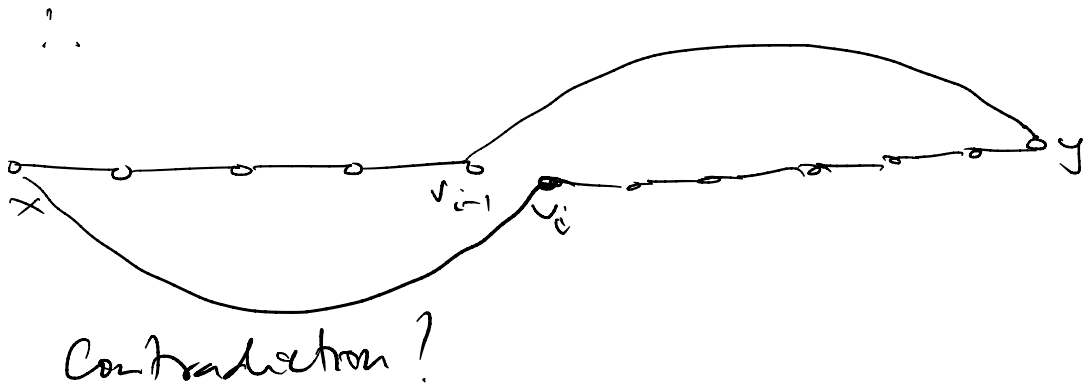
\therefore No graph exists satisfying the hypothesis & not the conclusion (WHY?)

\therefore Theorem is true;

on the maximal graph.



on the H-Path $\exists i$, s.t
 $x - v_i$ is an edge and $v_{i-1} - y$ is
an edge; **WHY?**



- Directed graphs-

- Tournaments.