

End-Semester Exam (ToC)

23/11/2016

(Let w^R denote the reverse of a word w . That is, if $w = a_1a_2 \dots a_n$, then $w^R = a_n \dots a_1$.)

1. **[3 marks]** Consider the context-free grammar G_1 given below:

$$S \rightarrow aSb \mid aSbb \mid aaSb \mid aS \mid Sb \mid \epsilon$$

Is $L(G_1)$ regular? If yes, give the *minimal* DFA. If not, argue using Myhill-Nerode theorem or pumping lemma.

2. **[6 marks]** Consider the context-free grammar given by G_2 below.

$$S \rightarrow aS \mid bA \mid aB$$

$$A \rightarrow aS \mid bB$$

$$B \rightarrow bS \mid aA \mid bA \mid \epsilon$$

Is $L(G_2)$ regular? If yes, give the *minimal* DFA. If not, argue using Myhill-Nerode theorem or pumping lemma.

3. **[9 marks]** For each of the following languages, state whether it is context-free or not. Justify your answers.

(a) $\{xyx^Ry^R \mid x, y \in \{a, b\}^*\}$

(b) $\{xyy^Rx^R \mid x, y \in \{a, b\}^*\}$

(c) $\{xx^Rx \mid x \in \{a, b\}^*\}$

The Intersection-non-emptiness problem of context-free grammars is given below.

Problem:	Intersection-Non-emptiness
Input:	G_1, G_2 : two context-free grammars
Question:	Is $L(G_1) \cap L(G_2) \neq \emptyset$?

4. **[4 marks]** Show that Intersection-Non-emptiness is undecidable by a reduction from Post's correspondence problem.
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For a language $L \subseteq \Sigma^*$, let $\text{REV-CLOSURE}(L) = \{w \mid w \in L \text{ or } w^R \in L\}$. Consider the following problem.

Problem:	Reverse-closedness
Input:	a Turing Machine description $\langle M \rangle$
Question:	Is $L(M) = \text{REV-CLOSURE}(L(M))$?

5. **[5 marks]** Prove that Reverse-closedness is undecidable by a reduction from the halting problem for Turing machines.
6. **[5 marks]** Is the set $\{\langle M \rangle \mid L(M) = \text{REV-CLOSURE}(L(M))\}$ recursively enumerable? Is it co-recursively enumerable? Justify.
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