

Conflict-Tolerant Specifications for Hybrid Systems

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Joint work with Madhu Gopinathan, S. Ramesh, and Prahlaad Sampath.

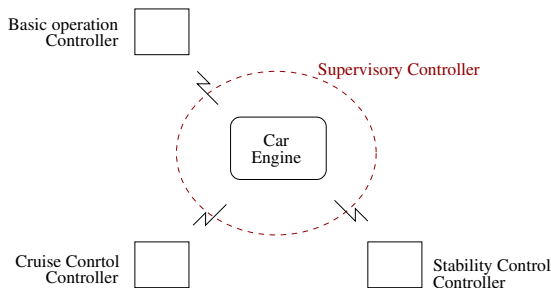
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Outline

- 1 Motivation
- 2 Hybrid automata and control
- 3 Conflict-tolerant specifications
- 4 Cruise Control Example
- 5 Verification for Rectangular Automata

Overview

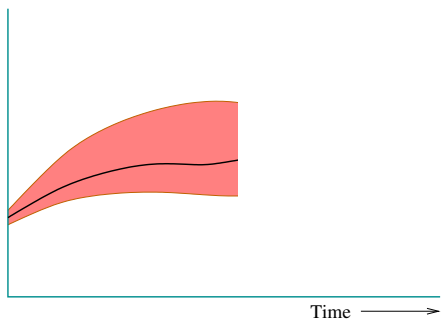
- We consider systems which are composed of a **base system** + multiple **controllers** + a **supervisor**.
- Supervisor chooses when to provide the control input of a controller to the base system (e.g. based on priority).



We propose a way of specifying the behaviour of individual controllers.

Why a classical safety specification is inadequate

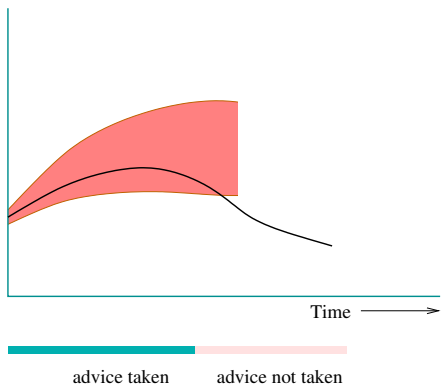
Specifies a **safety “cone”** (prefix-closed set of behaviours) within which the behaviour of the controlled system must lie.



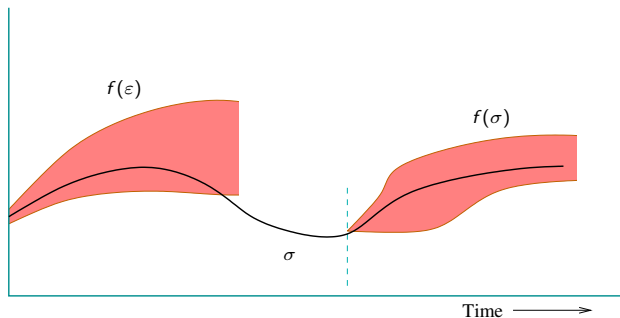
Why a classical safety specification is inadequate

What happens if the controller's input is disregarded by the supervisor (due to a conflict)?

- The resumed controller has no specification to adhere to.



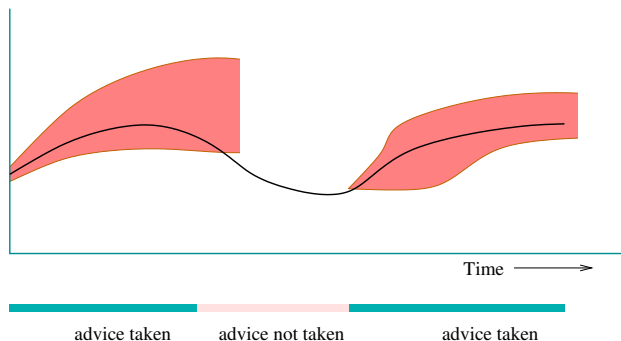
Conflict-tolerant specification



An **advice function** f which specifies a safety cone $f(\sigma)$ after **each** behaviour σ of the base system.

Conflict-tolerant Specification: Guarantee

Suppose a controller \mathcal{C} satisfies its tolerant specification \mathcal{S} wrt a base system \mathcal{B} .



- Then in every period in which it is control, \mathcal{C} does “the right thing” (according to \mathcal{S}).
- This guarantee is regardless of other controllers/supervisors it is composed with

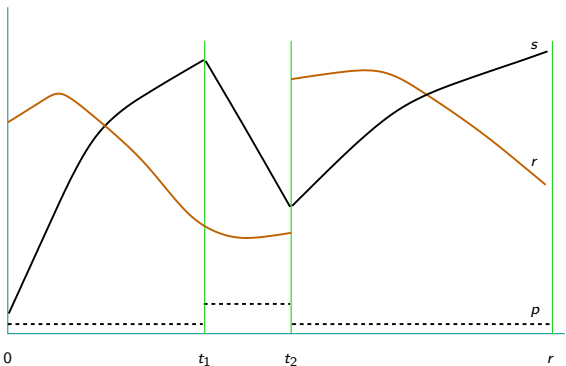
In this talk

- We give a mechanism to describe conflict-tolerant specifications for hybrid systems.
- Solve the verification problem when components are given as initialized rectangular hybrid automata.

Behaviour of a hybrid system

A **signal** over a set of variables W is a function $\sigma : [0, r] \rightarrow \mathbf{W}$ which has only finitely many points of discontinuity.

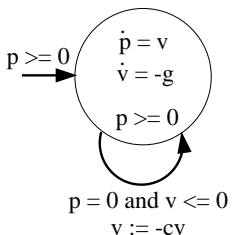
- There is a strictly increasing sequence of time points $t_0 = 0 < t_1 < t_2 < \dots < t_n = r$ such that σ is continuous in the interval $[t_k, t_{k+1})$.



Example hybrid system: Bouncing ball

System variables:

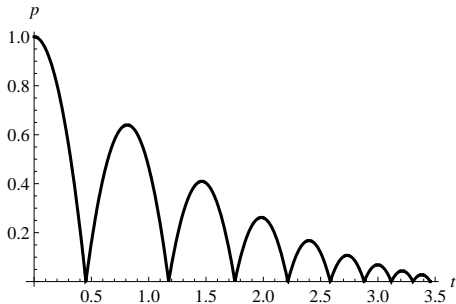
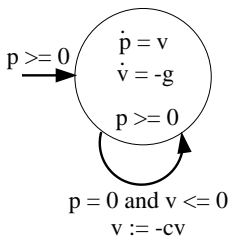
- p : vertical position (height) of the ball.
- v : velocity of ball.



Example hybrid system: Bouncing ball

System variables:

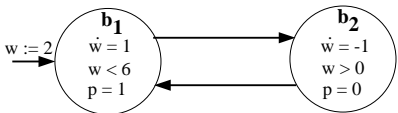
- p : vertical position (height) of the ball.
- v : velocity of ball.



Example hybrid system: Water tank with pump

System variables:

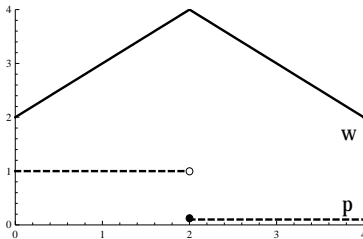
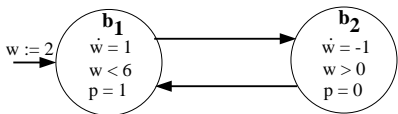
- w : level of water in tank.
- p : On/Off status of pump (1=“on”).



Example hybrid system: Water tank with pump

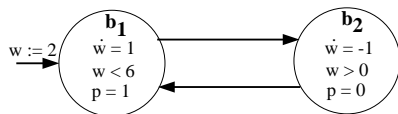
System variables:

- w : level of water in tank.
- p : On/Off status of pump (1=“on”).

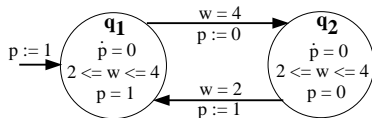


Controller for water tank

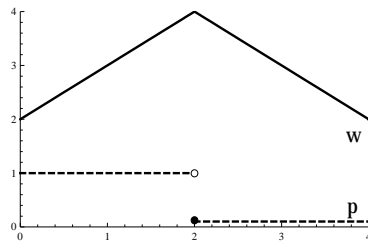
- System variables: $X = \{w\}$
- Control variables: $U = \{p\}$.



Water Tank

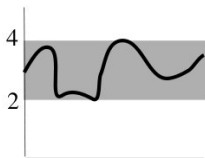
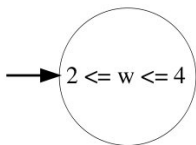


Controller



A classical specification for water level controller

Classical safety specification = prefix-closed set of signals.

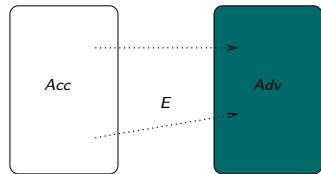


Conflict-tolerant specification

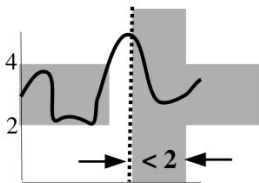
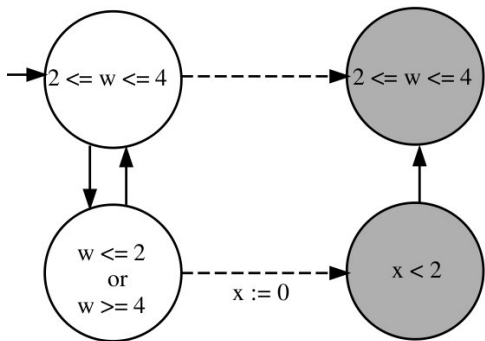
An **advice function** over a set of variables W is function $f : \text{Signals} \rightarrow 2^{\text{Signals}}$ such that each $f(\sigma)$ is prefix-closed.

Mechanism to specify such advice functions: $\mathcal{S} = (Acc, Adv, E)$ where

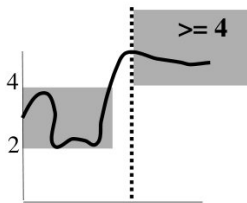
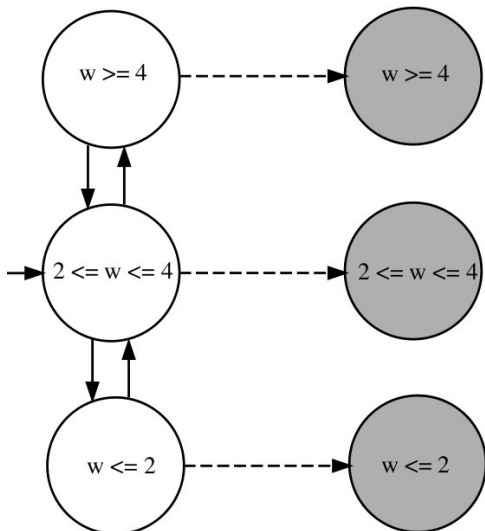
- Acc and Adv are hybrid automata over W
- E is a set of edges between Acc and Adv called an advice relation.



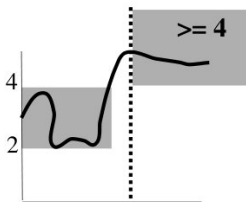
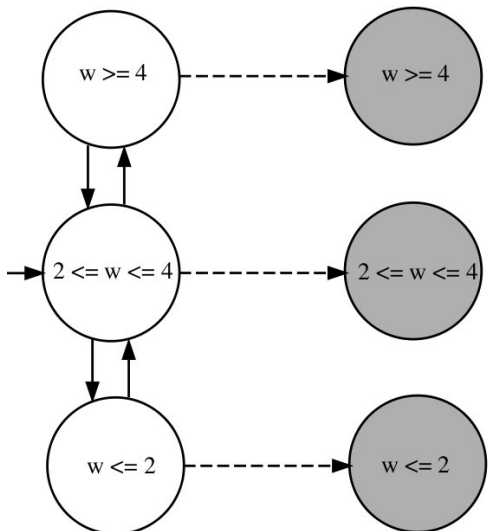
Example tolerant specifications for water level controller I



Example tolerant specifications for water level controller II



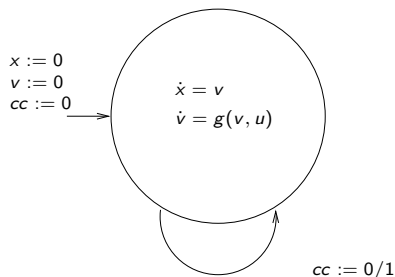
Example tolerant specifications for water level controller II



Note: Both tolerant specs induce the same classical spec but are quite different as tolerant specs.

Car motor base system

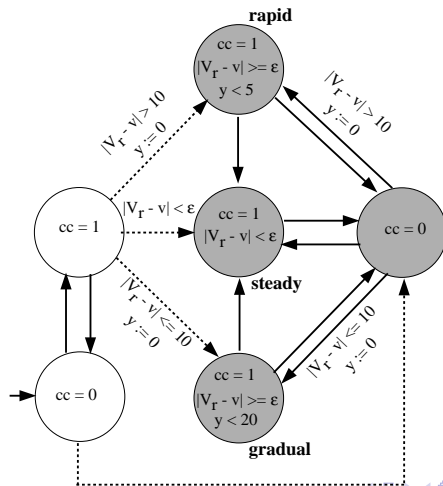
System variables: $\{x, v, cc\}$, Input variables: $\{u\}$.



Consider two controllers: for “cruise-control” and “stability-control”.

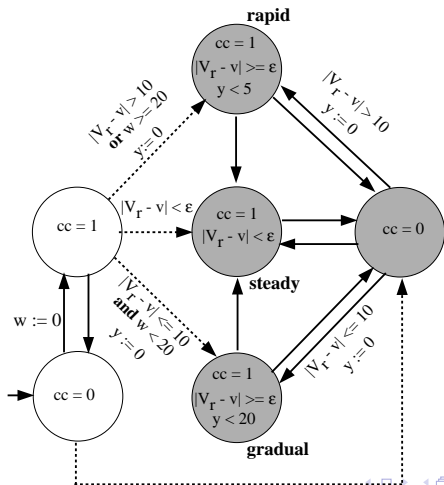
Cruise control tolerant specification I

“Reach set point within 20 sec if already close to it, else reach within 5 sec.”



Cruise control tolerant specification II

“Reach set point within 5 sec if far from it or if cc has been on for more than 20 sec; else reach within 20 sec.”



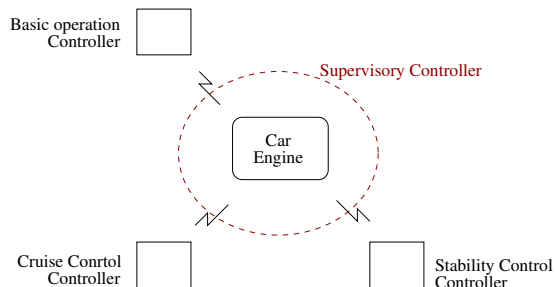
Verification for rectangular hybrid automata

We can solve the verification problem for such specs: Given

- a base system \mathcal{B} modelled as an initialized rectangular automaton over (X, Y) ,
- a controller \mathcal{C} modelled as an initialized rectangular automaton over (X, Y) ,
- a tolerant spec $\mathcal{S} = (Acc, Adv, E)$ whose components are IRHA:

we can check whether \mathcal{C} satisfies \mathcal{S} wrt \mathcal{B} .

Conclusion



- Conflict-tolerant specifications more richly capture a controller's specification.
- A modular or "compositional" way of developing and reasoning about systems with multiple controllers.
- A mechanism to specify them via hybrid automata.
- Decision procedure for the verification problem when components are given as initialised rectangular hybrid