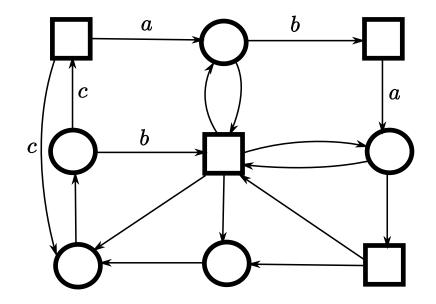
# Positional equilibria in infinite perfect information games

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ACTS



#### **Transition systems**

$$\mathcal{A} = (S, \Delta)$$

S— a finite set of states, C— a set of colors,  $\Delta \subset S \times C \times S$  — a finite set of *actions*.

$$e = (s, c, t) \in \Delta$$

s = source(e) — the source, t = target(e) — the target,  $c = \gamma_C(e)$  — the colour of e.

$$\Delta(s) = \{ e \in \Delta \mid \text{source}(e) = s \}$$

the set of actions *available at* s.

s a sink state if  $\Delta(s) \neq \emptyset$ .

## Paths/Plays

$$p = e_1 e_2 e_3 \dots$$
  $\forall i \ge 0, \operatorname{target}(e_i) = \operatorname{source}(e_{i+1}).$ 

$$\gamma_C(p) = \gamma_C(e_1e_2\ldots) = \gamma_C(e_1)\gamma_C(e_2)\ldots$$

#### Arenas

$$\mathcal{A} = (S, \Delta, \pi)$$

- $(S, \Delta)$  a transition system without sink states,
- $\pi: S \to {\min, \max}$  is a mapping designating for each state s the player  $\pi(s)$  controlling s.

## Outcomes

An outcome of an infinite play p is the  $\gamma_C(p)$ .

The set of outcomes

$$\mathcal{O}(C) = \bigcup_{\substack{B \subset C \\ B \text{ finite nonempty}}} B^{\omega}.$$

#### **Preference relation**

A binary relation  $\supseteq$  over the set  $\mathcal{O}(C)$  of outcomes

• reflexive, i.e.  $u \sqsupseteq u$ , for all  $u \in \mathcal{O}(C)$ ,

- transitive, i.e.  $u \sqsupseteq v$  and  $v \sqsupseteq w$  imply  $u \sqsupseteq w$ , for  $u, v, w \in \mathcal{O}(C)$  and
- total, either  $u \sqsupseteq v$  or  $v \sqsupseteq u$ , for all  $u, v \in \mathcal{O}(C)$ .

A preference relation = a total preorder relation over the set  $\mathcal{O}(C)$  of outcomes.

## Meaning

$$u \sqsupseteq v, \quad u, v \in \mathcal{O}(C).$$

u is no worse than v.

A player strictly prefers u to v,  $u \sqsupset v$ , if

 $u \sqsupseteq v$  but not  $v \sqsupseteq u$ .

If  $u \sqsupseteq v$  and  $v \sqsupseteq u$  — a player is indifferent between u and v.

 $\sqsubseteq$  — the inverse of  $\supseteq$ .

#### Two-person strictly antagonistic game

 $(\mathcal{A}, \sqsupseteq),$ 

where  $\mathcal{A}$  is an arena and  $\supseteq$  is a preference relation for Max the preference relation for player Min is  $\sqsubseteq$ .

The obvious aim of each player is to obtain the most advantageous outcome with respect to his preference relation.

## **Preferences versus payoff mappings**

Payoff mapping

$$f: \mathcal{O}(C) \to \mathbb{R} \cup \{-\infty, +\infty\}$$

induces preference  $\Box_f$ ,

 $u \sqsupseteq_f v$  if  $f(u) \ge f(v)$ .

#### Strategies and equilibria

 $\mathcal{A} = (S, \Delta, \pi)$  – an arena.

$$S_{\text{Max}} = \{ s \in S \mid \pi(s) = \text{Max} \}$$

states controlled by player  ${\rm Max}$ 

$$S_{\mathrm{Min}} = S \setminus S_{\mathrm{Max}}$$

states controlled by player Min.

A strategy for player  $\mu \in {Max, Min}$  is a mapping

 $\sigma_{\mu}: \{ p \in \mathscr{P}(\mathcal{A}) \mid \operatorname{target}(p) \in S_{\mu} \} \to \Delta,$ 

such that  $\sigma_{\mu}(p) \in \Delta(s)$ , where s = target(p).

#### Plays consistent with a strategy

 $p=e_0e_1e_2\ldots$  is consistent with player  $\mu$  's strategy  $\sigma_\mu$  if, for each factorization p=p'p'', such that

• p'' is nonempty

• and  $\operatorname{target}(p') = \operatorname{source}(p'')$  is controlled by player  $\mu$ ,

 $\sigma_{\mu}(p')$  is the first action in p''.

### **Positional strategies**

A *positional* (or memoryless) strategy for player  $\mu$ 

$$\sigma_{\mu}: S_{\mu} \to \Delta$$

such that, for all  $s \in S_{\mu}$ ,

 $\sigma_{\mu}(s) \in \Delta(s)$ 

A strategy profile is a pair  $(\sigma, \tau)$  of strategies.

 $p_{\mathcal{A},s}(\sigma,\tau)$ 

unique play with source s consistent with  $\sigma$  and  $\tau.$ 

#### Equilibria

A Nash equilibrium  $(\sigma^{\#}, \tau^{\#})$  if for all states  $s \in S$  and all strategies  $\sigma$  and  $\tau$ ,

$$\gamma_C(p_{\mathcal{A},s}(\sigma^{\#},\tau)) \sqsupseteq \gamma_C(p_{\mathcal{A},s}(\sigma^{\#},\tau^{\#})) \sqsupseteq \gamma_C(p_{\mathcal{A},s}(\sigma,\tau^{\#}))$$

An equilibrium  $(\sigma^{\#}, \tau^{\#})$  is said to be *positional* if the strategies  $\sigma^{\#}$  and  $\tau^{\#}$  are positional.

## Mean-payoff games

$$C = \mathbb{R} \times \mathbb{R}_+$$

$$(r_1, t_1)(r_2, t_2)(r_3, t_3) \dots \supseteq (r'_1, t'_1)(r'_2, t'_2)(r'_3, t'_3) \dots$$

if

$$\lim_{n} \frac{r_1 t_1 + r_2 t_2 + \dots + r_n t_n}{t_1 + t_2 + \dots + t_n} \ge \lim_{n} \frac{r'_1 t'_1 + r'_2 t'_2 + \dots + r'_n t'_n}{t'_1 + t'_2 + \dots + t'_n}$$

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But then

 $1000, 1000, ..., 1000, 0^{\omega} \approx 0^{\omega}$ 

## Overtaking

$$(r_1, t_1)(r_2, t_2)(r_3, t_3) \dots \sqsupseteq (r'_1, t'_1)(r'_2, t'_2)(r'_3, t'_3) \dots$$
 if

$$\exists N, \forall n > N, \quad \frac{r_1 t_1 + r_2 t_2 + \dots + r_n t_n}{t_1 + t_2 + \dots + t_n} \ge \frac{r'_1 t'_1 + r'_2 t'_2 + \dots + r'_n t'_n}{t'_1 + t'_2 + \dots + t'_n}$$

## Weighted limits

$$C = \mathbb{R}, \quad \alpha \in [0, 1]$$

$$f_{\alpha}(r_1r_2r_3\ldots) = \alpha \cdot \limsup_i r_i + (1-\alpha) \cdot \liminf_i r_i$$

#### Extended preference relation and $\succeq$ - equilibria

The extended preference relation  $\succeq$  is defined as follows:

for  $x, y \in \mathcal{O}(C)$ ,  $x \succeq y$  if  $\forall u \in C^*, ux \sqsupseteq uy$ .

Obviously, if  $x \succeq y$  then  $x \sqsupseteq y$ .

 $\succeq$  is transitive and reflexive, but maybe not total.

## $\succeq$ -equilibria

A strategy profile  $(\sigma^{\#}, \tau^{\#})$  is a  $\succeq$ -equilibrium if for all strategies  $\sigma$ ,  $\tau$ 

$$\gamma_C(p_{\mathcal{A},s}(\sigma^{\#},\tau)) \succeq \gamma_C(p_{\mathcal{A},s}(\sigma^{\#},\tau^{\#})) \succeq \gamma_C(p_{\mathcal{A},s}(\sigma,\tau^{\#})) \quad .$$

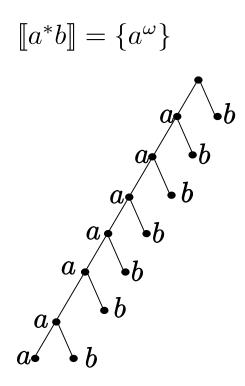
### **Adherence operator**

$$\llbracket \ \rrbracket: 2^{C^*} \to 2^{C^\omega}$$

For  $L \subseteq C^*$ ,

 $\llbracket L \rrbracket = \{ u \in C^{\omega} \mid \operatorname{Pref}(u) \subset \operatorname{Pref}(L) \} .$ 

## Exercise



### Why adherence?

 $\mathcal{A} = (S, \Delta)$  an arena. Then

$$L_s^{\omega}(\mathcal{A}) = \llbracket L_s(\mathcal{A}) \rrbracket$$

If  $L \in \operatorname{Rec}(C^*)$  then  $L_s^{\omega}(\mathcal{A}) = \llbracket L \rrbracket$  for some arena  $\mathcal{A}$ .

### **Properties of the adherence**

**Lemma 1.** Let  $L, M \subseteq C^*$  be finitely generated and  $u \in C^*$ . Then

$$\llbracket uL \rrbracket = u \llbracket L \rrbracket, \tag{1}$$

$$\llbracket \operatorname{Pref}(\llbracket M \rrbracket) \rrbracket = \llbracket M \rrbracket, \tag{2}$$

$$\llbracket L \cup M \rrbracket = \llbracket L \rrbracket \cup \llbracket M \rrbracket, \tag{3}$$

$$\llbracket LM \rrbracket = \llbracket L \rrbracket \cup L \llbracket M \rrbracket, \tag{4}$$

$$\llbracket L^* \rrbracket = (L \setminus \epsilon)^{\omega} \cup L^* \llbracket L \rrbracket \quad .$$
(5)

#### **Conditions for positional equilibria**

Let  $u \in \mathcal{O}(C)$  and  $X \subset \mathcal{O}(C)$ .

Notation.

 $u \succeq X$ 

if, for all  $x \in X$ ,  $u \succeq x$ .

#### Ultimately periodic infinite words

Let  $u, w \in C^*$  and  $v \in C^+$ .

An infinite word of the form

 $uv^{\omega},$ 

is called *ultimately periodic*.

## Simple periodic languages

Let  $u, w \in C^*$  and  $v \in C^+$ .

 $uv^*$ 

Note

 $\llbracket uv^* \rrbracket = \{uv^\omega\}$ 

## Union selection.

 $\succeq$  satisfies *union selection* condition if, for all ultimately periodic words  $u_1u_2^{\omega}$  and  $v_1v_2^{\omega}$ , either

$$u_1 u_2^{\omega} \succeq v_1 v_2^{\omega}$$

or

$$v_1 v_2^{\omega} \succeq u_1 u_2^{\omega}$$

We can rewrite this condition as

$$\exists x \in \{u_1 u_2^{\omega}, v_1 v_2^{\omega}\}, \quad x \succeq \llbracket u_1 u_2^* \cup v_1 v_2^* \rrbracket.$$

#### **Product selection.**

We say that  $\succeq$  satisfies *product selection* condition for player Max if, for all  $u, v, w, z \in C^*$  such that |v| > 0 and |w| > 0,

$$\exists x \in \{uv^{\omega}, uwz^{\omega}\}, \quad x \succeq \llbracket uv^*wz^* \rrbracket.$$

Note that

$$\{uv^{\omega}, uwz^{\omega}\} \subset uv^{\omega} \cup uv^*wz^{\omega} = \llbracket uv^*wz^* \rrbracket.$$

#### Star selection.

 $\succeq$  satisfies star selection condition for player  ${\rm Max}$  if for each nonempty language  $L\in {\rm Rec}(C^+)$ 

$$\exists x \in \llbracket L \rrbracket \cup \{ u^{\omega} \mid u \in L \}, \quad x \succeq \llbracket L^* \rrbracket.$$

Note

$$\llbracket L \rrbracket \cup \{ u^{\omega} \mid u \in L \} \subset \llbracket L \rrbracket \cup L^{\omega} = \llbracket L^* \rrbracket$$

#### Remark

If  $\succeq$  satisfies all three selection conditions then for each  $L \in \text{Rec}(C)$ , if  $[\![L]\!] \neq \emptyset$  then

 $\exists uv^{\omega} \in \llbracket L \rrbracket$  such that  $uv^{\omega} \succeq \llbracket L \rrbracket$ 

### One player $\operatorname{Max}$ games

 $\succeq \text{ satisfies all three selection conditions} \\ \text{if and only if} \\ \text{one-player Max games have optimal positional strategies for player Max.} \\$ 

## **Dual conditions.**

$$\succeq \quad \leftrightarrow \quad \preceq \, .$$

#### Main result

**Theorem 2.** Let  $\supseteq$  be a preference relation over  $\mathcal{O}(C)$  and let  $\succeq$  be the corresponding extended preference relation. The following conditions are equivalent:

- (1) There exist positional equilibria for all games  $(\mathcal{A}, \supseteq)$  over finite arenas.
- (2) There exist positional  $\succeq$ -equilibria for all games  $(\mathcal{A}, \sqsupseteq)$  over finite arenas.
- (3)  $\succeq$  satisfies union selection, product selection and star selection conditions for player Max and player Min.

- (4) For all one-player games  $(\mathcal{A}, \supseteq)$  the player controlling the arena  $\mathcal{A}$  has an optimal positional strategy.
- (5) For all one-player games  $(\mathcal{A}, \supseteq)$  the player controlling the arena  $\mathcal{A}$  has a  $\succeq$ -optimal positional strategy.