

Merge Sort: Shortcomings

- * Merging A and B creates a new array C
 - * No obvious way to efficiently merge in place
- * Extra storage can be costly
- * Inherently recursive
 - * Recursive call and return are expensive

Divide and conquer without merging

- * Suppose the median value in A is m
- * Move all values $\leq m$ to left half of A
 - * Right half has values $> m$
 - * This shifting can be done in place, in time $O(n)$
- * Recursively sort left and right halves
- * A is now sorted! No need to merge
 - * $t(n) = 2t(n/2) + n = O(n \log n)$

Alternative approach

- * Extra space is required to merge
- * Merging happens because elements in left half must move right and vice versa
- * Can we divide so that everything to the left is smaller than everything to the right?
 - * No need to merge!

Divide and conquer without merging

- * How do we find the median?
 - * Sort and pick up middle element
 - * But our aim is to sort!
- * Instead, pick up some value in A — **pivot**
 - * Split A with respect to this pivot element

Quicksort

- * Choose a pivot element
 - * Typically the first value in the array
- * Partition A into lower and upper parts with respect to pivot
- * Move pivot between lower and upper partition
- * Recursively sort the two partitions

Quicksort

- * High level view

43	32	22	78	63	57	91	13
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Quicksort

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Quicksort: Partitioning

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Diagram illustrating the initial partitioning step of Quicksort. The array is [43, 32, 22, 78, 63, 57, 91, 13]. The pivot is 22 (orange). The left pointer is at index 2 (orange arrow) and the right pointer is at index 7 (green arrow).

Quicksort: Partitioning

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Diagram illustrating the first swap in the partitioning step. The array is [43, 32, 22, 13, 63, 57, 91, 78]. The pivot is 22 (orange). The left pointer is at index 3 (orange arrow) and the right pointer is at index 7 (green arrow).

Quicksort: Partitioning

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Diagram illustrating the second swap in the partitioning step. The array is [43, 32, 22, 78, 63, 57, 91, 13]. The pivot is 22 (orange). The left pointer is at index 2 (orange arrow) and the right pointer is at index 7 (green arrow).

Quicksort: Partitioning

13	32	22	43	63	57	91	78
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Diagram illustrating the final state of the partitioning step. The array is [13, 32, 22, 43, 63, 57, 91, 78]. The pivot is 22 (orange). The left pointer is at index 2 (orange arrow) and the right pointer is at index 7 (green arrow).

Quicksort: Implementation

```
Quicksort(A,l,r) // Sort A[l..r-1]
```

```
    if (r - l <= 1) return; // Base case
```

```
    // Partition with respect to pivot, a[l]
```

```
    yellow = l+1;
```

```
    for (green = l+1; green < r; green++)
```

```
        if (A[green] <= A[l])
```

```
            swap(A,yellow,green);
```

```
            yellow++;
```

```
    swap(A,l,yellow-1); // Move pivot into place
```

```
    Quicksort(A,l,yellow); // Recursive calls
```

```
    Quicksort(A,yellow+1,r);
```

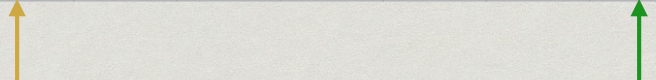
Quicksort: Another Partitioning Strategy

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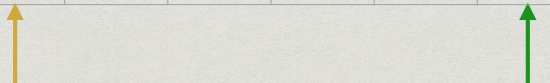
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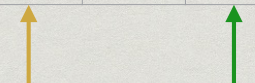
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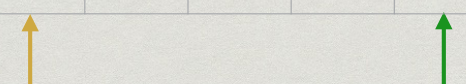
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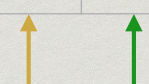
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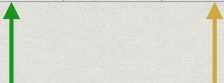
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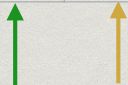
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 - * Typically the first value in the array
- * Partition A into lower and upper parts with respect to pivot
- * Move pivot between lower and upper partition
- * Recursively sort the two partitions

Analysis of Quicksort

- * Partitioning with respect to pivot takes $O(n)$
- * If pivot is median
 - * Each partition is of size $n/2$
 - * $t(n) = 2t(n/2) + n = O(n \log n)$
- * Worst case?

Analysis of Quicksort

But ...

- * Average case is $O(n \log n)$
 - * Sorting is a rare example where average case can be computed
- * What does average case mean?

Analysis of Quicksort

Worst case

- * Pivot is maximum or minimum
 - * One partition is empty
 - * Other is size $n-1$
 - * $t(n) = t(n-1) + n = t(n-2) + (n-1) + n$
 $= \dots = 1 + 2 + \dots + n = O(n^2)$
- * Already sorted array is worst case input!

Quicksort: Average case

- * Assume input is a permutation of $\{1, 2, \dots, n\}$
 - * Actual values not important
 - * Only relative order matters
 - * Each input is equally likely (uniform probability)
- * Calculate running time across all inputs
- * **Expected running time** can be shown $O(n \log n)$

Quicksort: randomization

- * Worst case arises because of fixed choice of pivot
 - * We chose the first element
 - * For any fixed strategy (last element, midpoint), can work backwards to construct $O(n^2)$ worst case
- * Instead, choose pivot **randomly**
 - * Pick any index in $[0..n-1]$ with uniform probability
- * Expected running time is again $O(n \log n)$

Quicksort in practice

- * In practice, Quicksort is very fast
 - * Typically the default algorithm for in-built sort functions
 - * Spreadsheets
 - * Built in sort function in programming languages

Iterative Quicksort

- * Recursive calls work on disjoint segments of array
 - * No recombination of results required
- * Can use an explicit stack to simulate recursion
 - * Stack only needs to store left and right endpoints of interval to be sorted

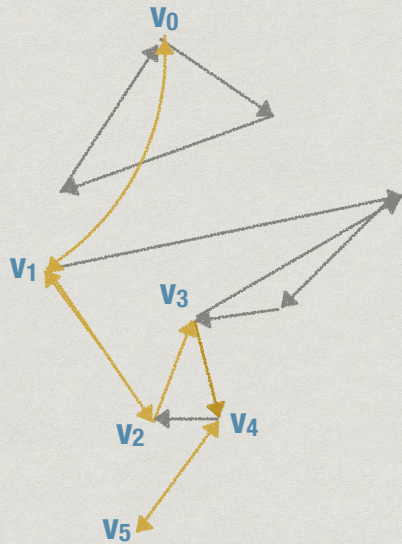
Graphs, formally

$$G = (V, E)$$

- * Set of vertices V
- * Set of edges E
 - * E is a subset of pairs (v, v') : $E \subseteq V \times V$
 - * Undirected graph: (v, v') and (v', v) are the same edge
 - * Directed graph:
 - * (v, v') is an edge from v to v'
 - * Does not guarantee that (v', v) is also an edge

Finding a route

- * Directed graph
- * Find a sequence of vertices v_0, v_1, \dots, v_k such that
 - * v_0 is New Delhi
 - * Each (v_i, v_{i+1}) is an edge in E
 - * v_k is Trivandrum

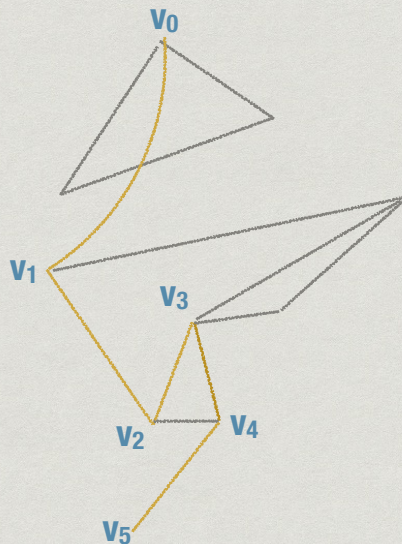


Working with graphs

- * We are given $G = (V, E)$, **undirected**
- * Is there a path from source v_s to target v_t ?
- * Look at the picture and see if v_s and v_t are connected
- * How do we get an algorithm to “look at the picture”?

Finding a route

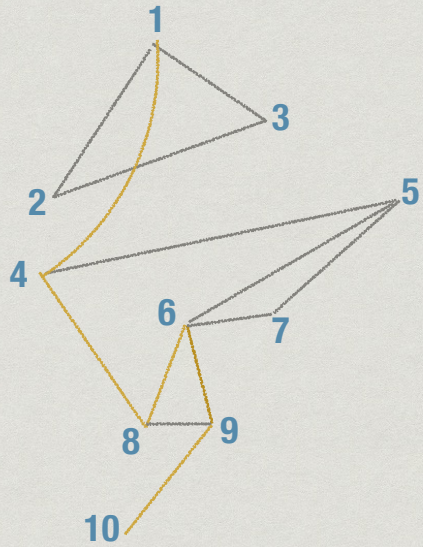
- * Also makes sense for undirected graphs
- * Find a sequence of vertices v_0, v_1, \dots, v_k such that
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Representing graphs

- * Let V have n vertices
 - * We can assume vertices are named $1, 2, \dots, n$
- * Each edge is now a pair (i, j) , where $1 \leq i, j \leq n$
- * Let $A(i, j) = 1$ if (i, j) is an edge and 0 otherwise
- * A is an $n \times n$ matrix describing the graph
 - * **Adjacency matrix**

Adjacency matrix

[illegible]

Adjacency matrix

- * Neighbours of i
 - * Any column j in row i with entry 1
 - * Scan row i from left to right to identify all neighbours
- * Neighbours of 4 are {1,5,8}

[illegible]

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Finding a path

- * Start with v_s
- * New Delhi is 1
- * Mark each neighbour as reachable
- * Explore neighbours of marked vertices
- * Check if target is marked
- * $v_t = 10 = \text{Trivandrum}$

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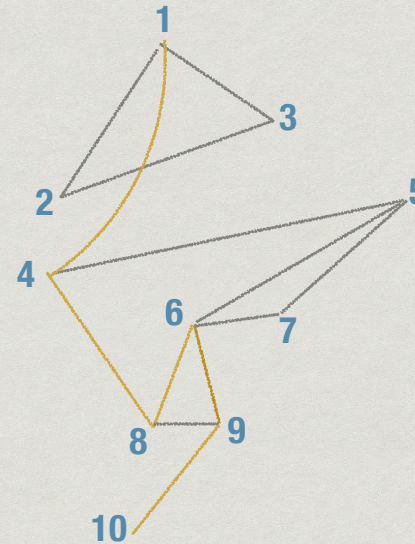
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Exploring graphs

- * Need a systematic algorithm
 - * Mark vertices that have been visited
 - * Keep track of vertices whose neighbours have already been explored
 - * Avoid going round indefinitely in circles
- * Two fundamental strategies: breadth first and depth first

Adjacency list



- * For each vertex, maintain a list of its neighbours

1	2,3,4
2	1,3
3	1,2
4	1,5,8
5	4,6,7
6	5,7,8,9
7	5,6
8	4,6,9
9	6,8,10
10	9

An alternative representation

- * Adjacency matrix has many 0's
- * Size of the matrix is n^2 regardless of number of edges
- * Maximum size of E is $n(n-1)/2$ if we disallow self loops
- * Typically E is much smaller

	1	2	3	4	5	6	7	8	9	10
1	0	1	1	1	0	0	0	0	0	0
2	1	0	1	0	0	0	0	0	0	0
3	1	1	0	0	0	0	0	0	0	0
4	1	0	0	0	1	0	0	1	0	0
5	0	0	0	1	0	1	1	0	0	0
6	0	0	0	0	1	0	1	1	1	0
7	0	0	0	0	1	1	0	0	0	0
8	0	0	0	1	0	1	0	0	1	0
9	0	0	0	0	0	1	0	1	0	1
10	0	0	0	0	0	0	0	0	1	0

Comparing representations

- * Adjacency list typically requires less space
- * **Is j a neighbour of i ?**
 - * Just check if $A[i][j]$ is 1 in adjacency matrix
 - * Need to scan neighbours of i in adjacency list
- * **Which vertices are neighbours of i ?**
 - * Scan all n columns in adjacency matrix
 - * Takes time proportional to neighbours in adjacency list