

WiFi

Students Network

Username: guest

Password: h1SiPCoT

Aadhaar

SIM Cards

Which SIM Cards do not have
corresponding Aadhaar

Match name/address

For every SIM card

Search for matching entry in
Aadhaar data

How long does this take?

SIM card list size S

Aadhaar list size A

$S \times A$
↓
for each
SIM
card L ? Unordered
 Worst Case
 Scan entire list

$S = ?$ 10^9 1 billion
 100 crore

$A = ?$ 10^9

$S \times A = 10^{18}$

Speed of a desktop computer?

10^9 ops/sec

$$\text{Time for SxA} = \frac{10^{18}}{10^9} = 10^9 \text{ sec}$$

$$60 \text{ sec} \rightarrow \text{min}$$

$$60 \times 60 \times 24 = 3600 \text{ sec hour}$$

$$\approx 100,000 \text{ sec/day } 10^5$$

$$10^8 \text{ sec} \approx 3 \text{ years}$$

$$10^9 \text{ sec} \approx 30 \text{ years}$$

Arrange Aadhaar data systematically

Sort alphabetically by Name, Address

$$\begin{array}{l} 1000 \text{ number} \\ 500 \\ 250 \\ \vdots \end{array}$$

$$(N/2)/2$$

$$2^k = N$$

$$\underline{2^{10} = 1024}$$

$$k = \log_2 N$$

Binary search on A $\approx 10^9$ entries

$$\log_2 10^9 = ? \Rightarrow 30$$

$$1024 = 10^3 \quad \log_2 1000 \approx 10$$

$$10^9 = 10^3 \cdot 10^3 \cdot 10^3$$

$$2^{30} = 2^{10} \cdot 2^{10} \cdot 2^{10}$$

$$S \times \log_2 A$$

$$10^9 \times 30 = 3 \times 10^{10}$$

$$30 \times 10^9 \underline{\underline{\text{sec}}}$$

$$S \approx A \approx N$$

Brute force

$$N^2 \approx 30 \text{ yrs}$$

Clever Version

$$N \log_2 N \approx 30 \text{ sec}$$

10^9 ops/sec

Complexity

$\log N$

N

$N \log N$

N^2

N^3

Feasible Input Size

2^{10^9}

10^9

$10^8 \sim 10^9$

10^4

10^3

Insertion Sort

Choose max $\longrightarrow$ 1st

Choose 2nd max $\longrightarrow$ 2nd

Select

,

,

Selection Sort

Sorting

Graded exam papers

Arrange in order of marks

Start with one

Position second relative to first

$\vdots$

Position k^{th} paper wrt $1, 2, \dots, k-1$

Insert k^{th} paper in sorted $1 \dots k-1$

Time?

Insertion Sort

N items

$N \times$

$$0 + 1 + 2 + \dots + (N-1) \\ N-1 + N-2 + \dots + 0$$

$$= \frac{N \cdot (N-1)}{2} \approx N^2$$

$O(N^2)$. For all $f(N) \leq cN^2$
for some c

Selection Sort

N papers - find the max of $N, N-1, N-2, \dots, 1$

$N + N-1 + \dots + 2 + 1$

$$\frac{N \cdot N+1}{2} \approx O(N^2)$$

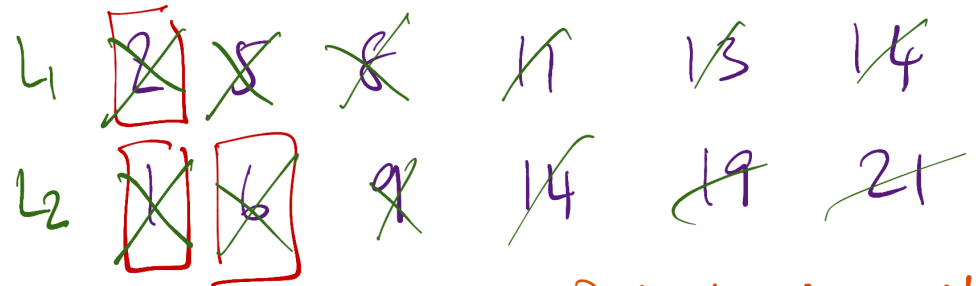
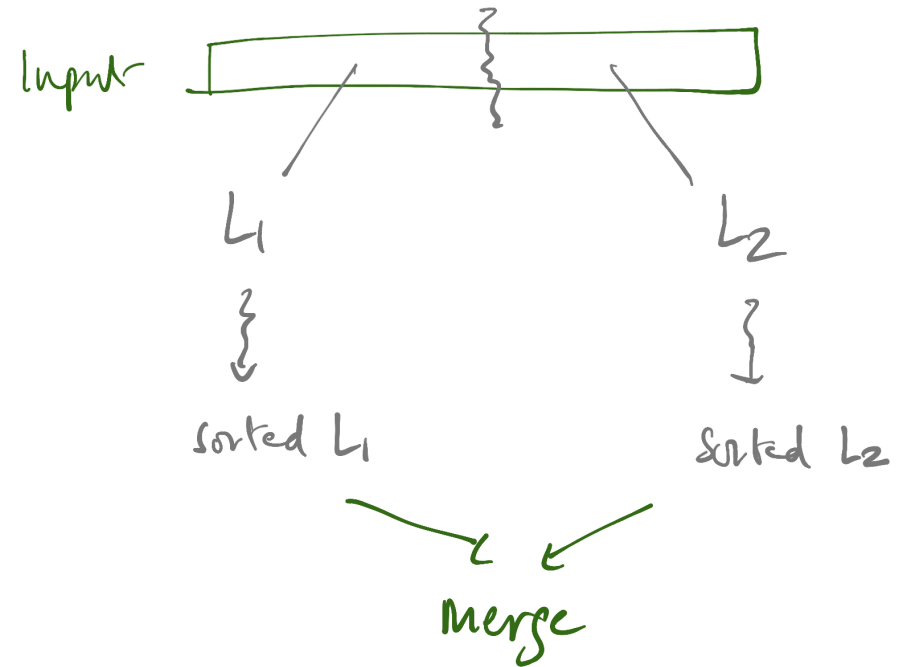
Selection Sort - Always N^2

Insertion Sort - Varies according to input

Ideally - "average case"

$S, A \rightarrow A^2$ insertion sort
 $\downarrow$
 $S \times \log A$ search

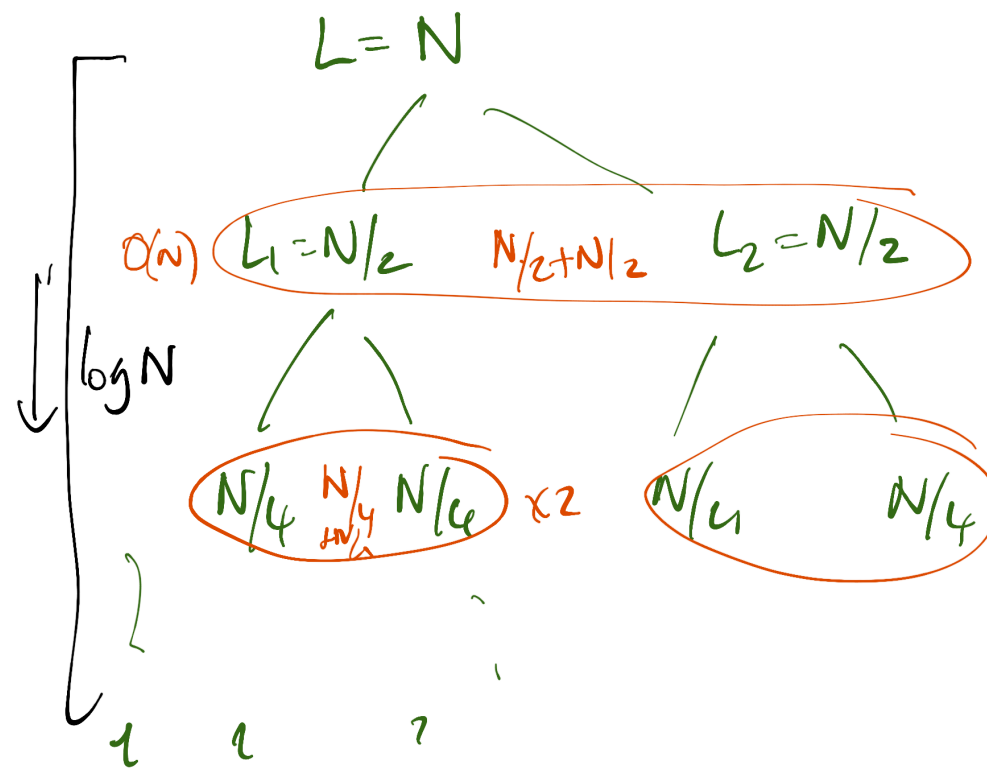
Divide & Conquer



$O(N)$ time - Output is size N
One comparison per output element

1 2 5 6 8 9 11 13

14 14 19 21



Each level requires $O(N)$ work to merge

$$O(N)$$

$$O(N/2) \times 2$$

$$O(N/4) \times 4$$

:

Overall $O(N \log N)$

Binary Search

$$T(N) = 1 + T(N/2)$$

$$T(1) = 1$$

$$T(N) = 1 + T(N/2)$$

$$= 1 + 1 + T(N/4)$$

$$= 1 + 1 + 1 + T(N/8) \dots$$

$$= \underbrace{1 + 1 + 1 + \dots}_{\log N} T(N/N)$$

Insertion Sort

$$T(N) = T(N-1) + N$$

$$T(N-2) + N-1$$

$$T(N-3) + N-2$$

Similarly

Selection Sort

$$T(1) + 2$$

$$T(0) + 1$$

Merge Sort

$$T(N) = 2 \cdot \underline{T(N/2)} + N$$

$$2 \cdot (2 \cdot T(N/4) + N/2) + N$$

$$= 2^2 T(N/2^2) + 2N$$

$$= 2^k T(N/2^k) + kN$$

$$\begin{array}{c} \parallel \\ 1 \\ k = \log N \end{array}$$

$$\frac{2^k T(N/2^k)}{N} + \frac{kN}{(\log N) \cdot N}, \quad k = \log N$$
$$N + N \log N$$