

# Assignment 1 : Complexity Theory\*

To be submitted by Feb 19th

1. Let  $S_1, S_2, f$  be space-constructible with  $S_2(n) \geq \log(n)$  and  $f(n) \geq n$ . Show that  $\text{NSPACE}(S_1(n)) \subseteq \text{SPACE}(S_2(n))$  implies  $\text{NSPACE}(S_1(f(n))) \subseteq \text{SPACE}(S_2(f(n)))$ .
2. For any real number  $a > 0$  and natural number  $k > 1$ , show that  $\text{DTIME}(n^k) \subset \text{DTIME}(n^k \log^a(n))$ . (Note that one consequence of this result is that within  $\mathbf{P}$ , there can be no “complexity gaps” of size  $(\log n)^{\Omega(1)}$ .)
3. **PRIMES** is in  $\mathbf{NP} \cap \mathbf{co-NP}$ .
4. Show that  $\mathbf{P} = \mathbf{NP} \implies \mathbf{EXP} = \mathbf{NEXP}$ .
5. Prove that the problem of deciding whether a polynomial with integer coefficients has an integer solution is  $\mathbf{NP}$  - hard.
6. Show that 2-SAT is in  $\mathbf{NL}$ .
7. Show that  $\mathbf{NP} \neq \mathbf{SPACE}(n)$  .
8.  $L$  is a sparse set if there is a polynomial  $p$  such that  $|L \cap \{0, 1\}^n| \leq p(n) \forall n$ . Prove that if a sparse set is  $\mathbf{NP}$ -complete, then  $\mathbf{P} = \mathbf{NP}$ .
9. Suppose we pick a random language  $C$  by choosing every string to be in  $C$  with probability  $\frac{1}{2}$ . Prove that, with high probability,  $\mathbf{P}^C \neq \mathbf{NP}^C$ .
10. Show that there is a language  $B \in \mathbf{EXP}$  such that  $\mathbf{NP}^B \neq \mathbf{P}^B$ .
11. Suppose  $L_1, L_2 \in \mathbf{NP} \cap \mathbf{co-NP}$ . Then show that  $L_1 \oplus L_2$  is in  $\mathbf{NP} \cap \mathbf{co-NP}$ , where  $L_1 \oplus L_2 = \{x : x \text{ is in exactly one of } L_1, L_2\}$ .

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\*Questions collected by Abhiroop Sanyal