

# Computational Complexity

## Assignment 1

February 13, 2013

Marks: 50

1. Prove that  $P \neq \text{SPACE}(n)$ . 7
2. The language 2SAT is defined as follows:  
2SAT = { $\phi : \phi$  is a satisfiable CNF formula where each clause has exactly two literals}.  
Prove that 2SAT is NL-complete. 7  
(Hint: Reduce 2SAT to directed s-t connectivity. For each clause  $x \vee y$ , draw a directed edge from  $\bar{x}$  to  $y$  and from  $\bar{y}$  to  $x$ . Show that the reduction is correct and can be done in logspace.)
3. Prove that  $\Sigma_k^p = \Pi_k^p$  implies PH =  $\Sigma_k^p$ . 7
4. *Cook reduction*: A language  $L$  is polynomial-time *Cook reducible* (also known as *Turing reducible*) to a language  $L'$  if there is a polynomial-time TM  $M$  that, given an oracle for deciding  $L'$ , can decide  $L$ .  
Show that Cook reducibility is transitive. Show that the usual notion of reductions (also called *Karp reduction* or *many-one reduction*) is a special case of Cook reduction. 7
5. Define UCYCLE = { $G \mid G$  is an undirected graph with a simple cycle }. Show that UCYCLE  $\in$  L. 7  
(Hint: You may assume that, for each vertex  $v$ , you are given a function  $next_v$  which gives a cyclic ordering of neighbours of  $v$ . Thus, if  $u$  is a neighbour of  $v$ , then  $next_v(u)$  is the neighbour of  $v$  which is next to  $u$  in that cyclic ordering.)
6. Define CYCLE = { $G \mid G$  is a directed graph that has a directed cycle}. Prove that CYCLE is NL-complete. 7  
(Hint: You may consider the notion of a layered graph  $G^{(k)}$  corresponding to the given graph  $G$ . The graph  $G^{(k)}$  has vertex set  $V \times [k]$ . If  $(u, v) \in E(G)$ , then  $((u, i), (v, i + 1)) \in E(G^{(k)})$  for  $1 \leq i < k$ .)
7. For a list  $S$  of integers, let  $mode(S)$  be the element that appears the maximum number of times in  $S$ . Show that there is a deterministic logspace algorithm that prints 'no' if  $mode(S)$  appears at most  $\lfloor n/2 \rfloor$  times, and prints 'yes' and  $mode(S)$  if  $mode(S)$  appears at least  $\lfloor n/2 \rfloor + 1$  times. Show that this can be done *simultaneously* in logspace and linear time. 8