

Unit-2: Model-checker NuSMV

B. Srivathsan

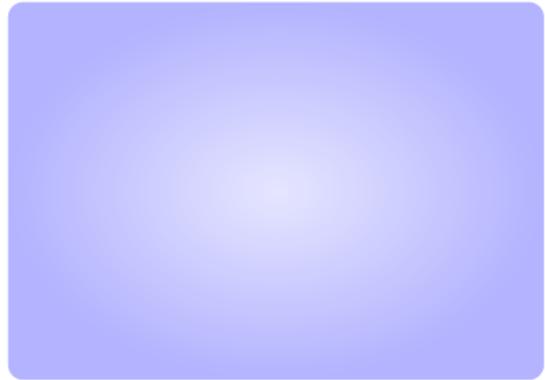
Chennai Mathematical Institute

NPTEL-course

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Module 4:
Modeling concurrent systems

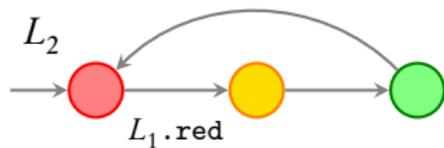
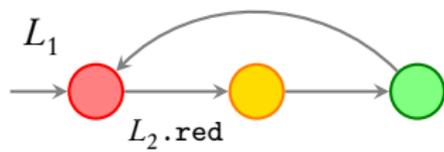
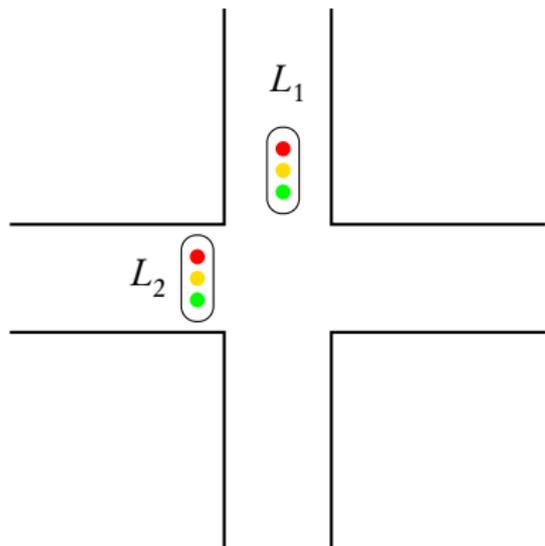
**Synchronous
vs.
Asynchronous
systems**

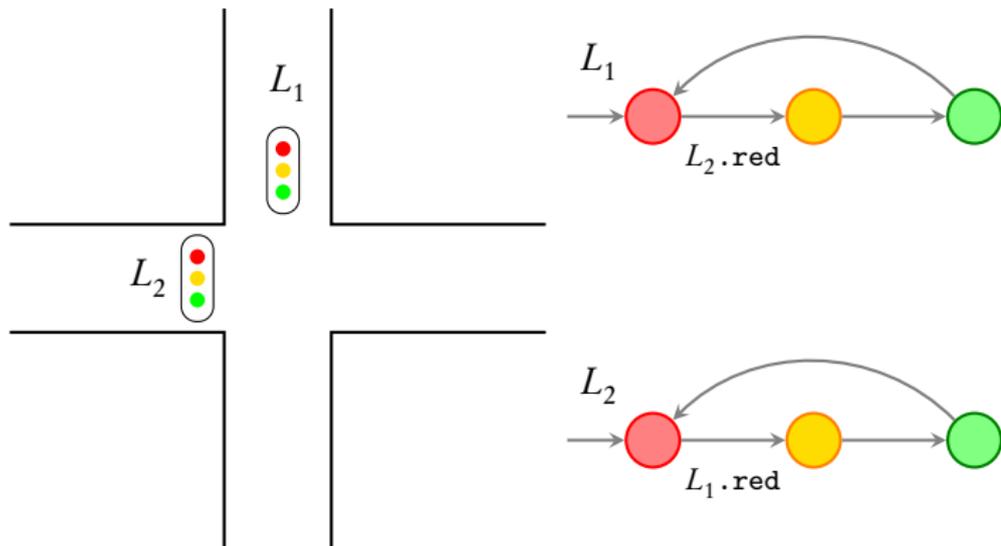


Acknowledgements:

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Prof. Supratik Chakraborty, IIT Bombay





If a light is **red**, it can **stay red** for an **arbitrary period**

If it goes **yellow**, it should become **green within one cycle**

If it is **green**, it can **stay green** for an **arbitrary period**

```
MODULE light(other)
VAR
    state: {r,y,g};
ASSIGN
    init(state) := r;
    next(state) := case
        state=r & other=r : {r, y};
        state=y : g;
        state=g : {g, r};
        TRUE : state;
    esac;

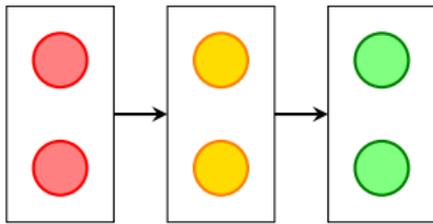
MODULE main
VAR
    t1: light(t12.state);  t12: light(t11.state);
```

Synchronous composition

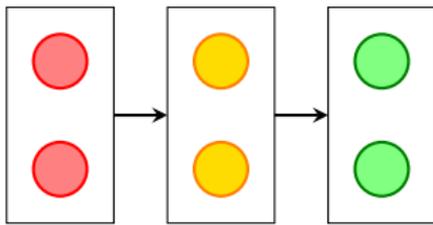
```
MODULE light(other)
VAR
    state: {r,y,g};
ASSIGN
    init(state) := r;
    next(state) := case
        state=r & other=r : {r, y};
        state=y : g;
        state=g : {g, r};
        TRUE : state;
    esac;

MODULE main
VAR
    t1: light(t2.state);  t2: light(t1.state);
```

Synchronous composition



Synchronous composition



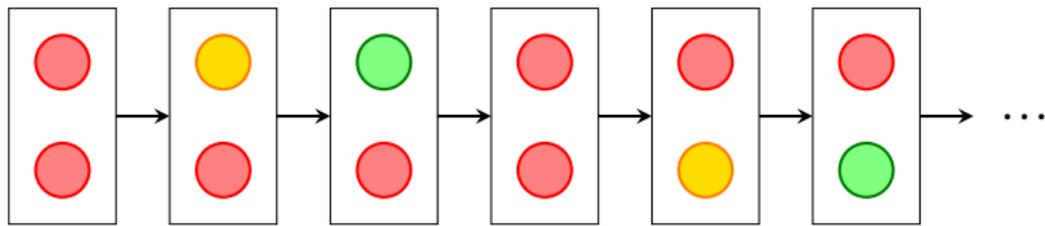
Both lights can **simultaneously** become green!

Asynchronous composition

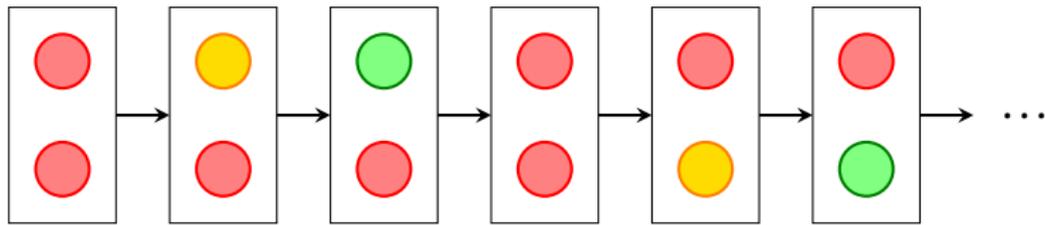
```
MODULE light(other)
VAR
    state: {r,y,g};
ASSIGN
    init(state) := r;
    next(state) := case
        state=r & other=r : {r, y};
        state=y : g;
        state=g : {g, r};
        TRUE : state;
    esac;

MODULE main
VAR
    t1: process light(t1.state);
    t2: process light(t2.state);
```

Asynchronous composition



Asynchronous composition



Only one light can become green at a time

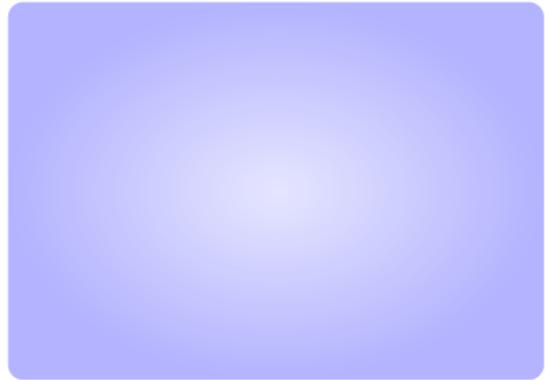
▶ **Synchronous:**

- ▶ all assignments to all modules made **simultaneously**
- ▶ suitable when all modules are synchronized to a **global clock**

▶ **Asynchronous:**

- ▶ execution of modules is **interleaved**
- ▶ at a time, **only one** module executes
- ▶ choice of next module to be executed is **non-deterministic**
- ▶ suitable when **no assumptions** can be made **about communication delay** between modules

**Synchronous
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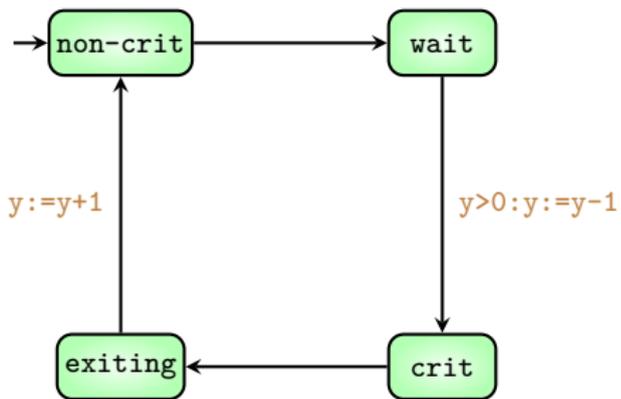


**Synchronous
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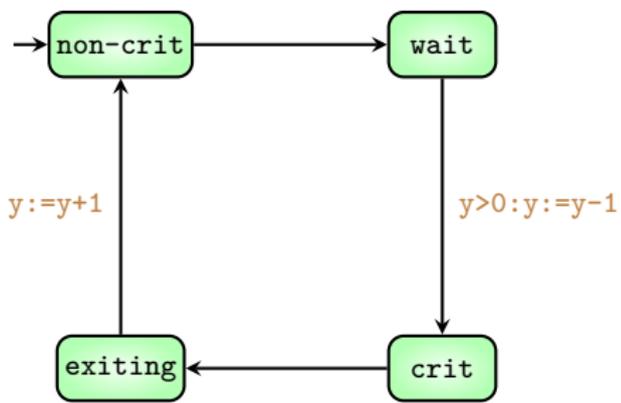
Mutual Exclusion



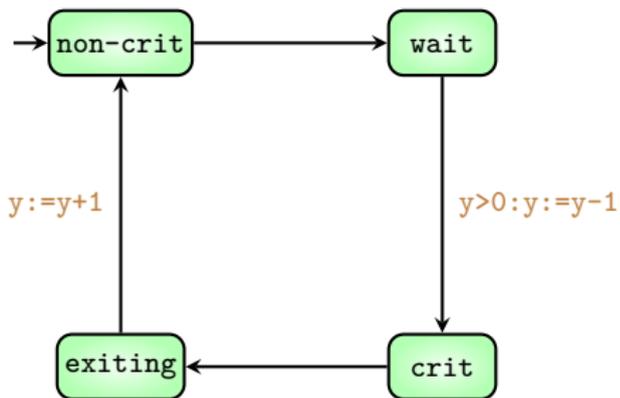
PG₁



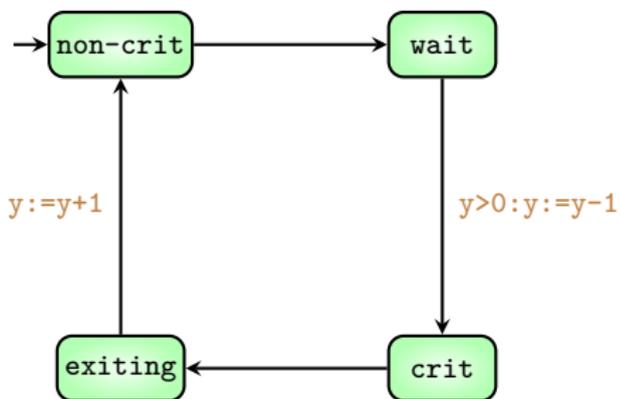
PG₂



PG₁



PG₂



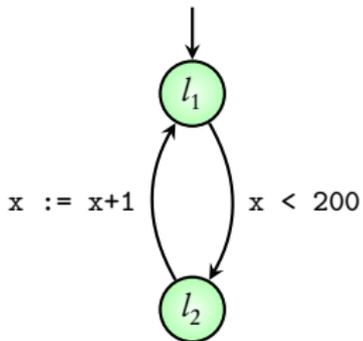
Synchronous
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systems

Mutual Exclusion



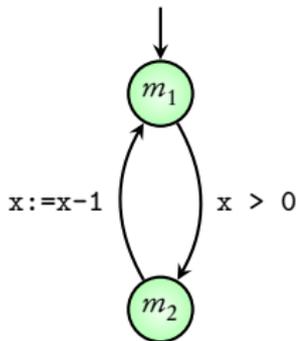
```
while x < 200
```

```
x := x+1
```



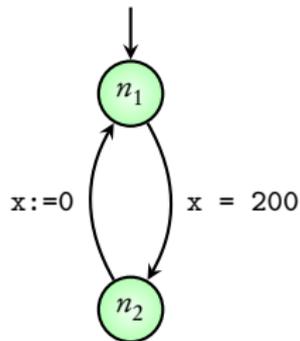
```
while x > 0
```

```
x := x-1
```



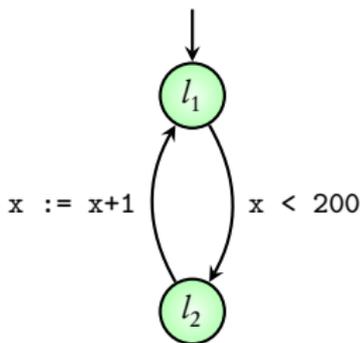
```
while x=200
```

```
x := 0
```



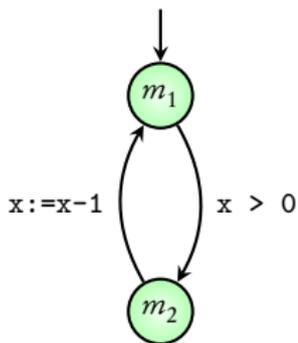
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while x < 200
```

```
x := x+1
```



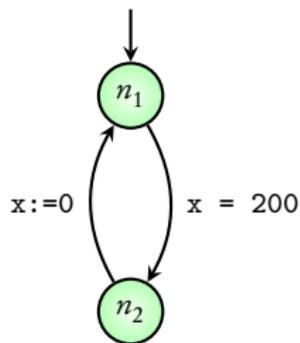
```
while x > 0
```

```
x := x-1
```



```
while x=200
```

```
x := 0
```



**Synchronous
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Asynchronous
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Mutual Exclusion

**Concurrent programs
example**