# Lecture 08, 07 September 2023

## Mutable and immutable values

- · Lists are dictionaries are mutable
- All other values are immutable (numbers, booleans, strings, tuples)
- immutable value : If x holds an immutable value, y = x copies the value to y
- mutable value : If 11 holds a mutable value, 12 = 11 makes both names point to the same value

## Functions and parameters

- Pass a mutable value, then it can updated in the function
- Immutable values will be copied

In [1]:	<pre>def mycopy(m,n):     m = n</pre>
In [2]:	a = 5 b = 7
	mycopy(a,b)
In [3]:	a, b
Out[3]:	(5, 7)
In [4]:	11 = [1,2] 12 = [3,4] mycopy(11,12)
In [5]:	11, 12
Out[5]:	([1, 2], [3, 4])
In [6]:	<pre>def mycopylist(m,n):     print(type(m))     m[0] = n[-1]</pre>
In [7]:	<pre>l1 = [1,2] l2 = [3,4] mycopylist(l1,12)</pre>
	<class 'list'=""></class>
In [8]:	11, 12
Out[8]:	([4, 2], [3, 4])
In [9]:	<pre>def myappend(l,v):     l.append(v)</pre>
In [10]:	myappend(12,5)
In [11]:	12
Out[11]:	[3, 4, 5]
In [12]:	<pre>def myappend2(1,v):     1 = 1 + [v]</pre>
In [13]:	myappend2(12,6)
In [14]:	12
Out[14]:	[3, 4, 5]

## Mutability and functions

It is useful to be able to update a list inside a function --- e.g. sorting it

- Built in list functions update in place
- 1. append (v) -> in place version of 1 = 1+ [v]
- 1.extend(11) -> in place version of 1 = 1 + 11

## Lists, arrays, dictionaries: implementation details

- What are the salient differences?
- How are they stored?
- What is the impact on performance?

#### Arrays

- · Contiguous block of memory
- · Typically size is declared in advance, all values are uniform
- a[0] points to first memory location in the allocated block
- Locate a [i] in memory using index arithmetic
- Skip i blocks of memory, each block's size determined by value stored in array
- Random access -- accessing the value at a[i] does not depend on i
- Useful for procedures like sorting, where we need to swap out of order values a[i] and a[j]
  - a[i], a[j] = a[j], a[i]
  - Cost of such a swap is constant, independent of where the elements to be swapped are in the array
- Inserting or deleting a value is expensive
- · Need to shift elements right or left, respectively, depending on the location of the modification

## Lists

- Each location is a *cell*, consisiting of a value and a link to the next cell
   Think of a list as a train, made up of a linked sequence of cells
- The name of the list 1 gives us access to 1[0], the first cell
- To reach cell 1[i], we must traverse the links from 1[0] to 1[1] to 1[2] ... to 1[i-1]] to 1[i]
  Takes time proportional to i
- Cost of swapping 1[i] and 1[j] varies, depending on values i and j
- On the other hand, if we are already at 1[i] modifying the list is easy
  - Insert create a new cell and reroute the links
  - Delete bypass the deleted cell by rerouting the links
- · Each insert/delete requires a fixed amount of local "plumbing", independent of where in the list it is performed

## Dictionaries

- Values are stored in a fixed block of size *m*
- Keys are mapped to  $\{0, 1, ..., m 1\}$
- Hash function  $h: K \to S$  maps a *large* set of keys K to a *small* range S
- Simple hash function: interpret  $k \in K$  as a bit sequence representing a number  $n_k$  in binary, and compute  $n_k \mod m$ , where |S| = m
- Mismatch in sizes means that there will be *collisions* --  $k_1 \neq k_2$ , but  $h(k_1) = h(k_2)$
- A good hash function maps keys "randomly" to minimize collisions
- Hash can be used as a *signature* of authenticity
  - Modifying k slightly will drastically alter h(k)
  - No easy way to reverse engineer a k' to map to a given h(k)
  - Use to check that large files have not been tampered with in transit, either due to network errors or malicious intervention
- · Dictionary uses a hash function to map key values to storage locations
- Lookup requires computing h(k) which takes roughly the same time for any k
- Compare with computing the offset a[i] for any index i in an array
- · Collisions are inevitable, different mechanisms to manage this, which we will not discuss now
- · Effectively, a dictionary combines flexibility with random access

#### Lists in Python

- · Flexible size, allow inserting/deleting elements in between
- · However, implementation is an array, rather than a list
- · Initially allocate a block of storage to the list
- When storage runs out, double the allocation
- 1. append (x) is efficient, moves the right end of the list one position forward within the array
- 1.insert(0,x) inserts a value at the start, expensive because it requires shifting all the elements by 1
- We will run experiments to validate these claims

#### Measuring execution time

- Call time.perf counter()
- · Actual return value is meaningless, but difference between two calls measures time in seconds

## In [15]: import time

• 10<sup>7</sup> appends to an empty Python list

· Doubling the work approximately doubles the time, linear

```
In [17]: start = time.perf_counter()
    l = []
    for i in range(20000000):
        l.append(i)
    elapsed = time.perf_counter() - start
    print(elapsed)
```

5.7539322239899775

•  $10^5$  inserts at the beginning of a Python list

5.2932833380036755

- Doubling and tripling the work multiplies the time by  $4 \mbox{ and } 9,$  respectively, so quadratic

44.08497351700498

• Creating  $10^7$  entries in an empty dictionary

```
In [21]: start = time.perf_counter()
d = {}
for i in range(10000000,0,-1):
    d[i] = i
elapsed = time.perf_counter() - start
print(elapsed)
```

4.113557312011835

- Doubling the operations, doubles the time, so linear
- Dictionaries are effectively random access

```
In [22]: start = time.perf_counter()
d = {}
for i in range(20000000,0,-1):
    d[i] = i
elapsed = time.perf_counter() - start
print(elapsed)
```

9.394316827994771