

Backtracking

Madhavan Mukund

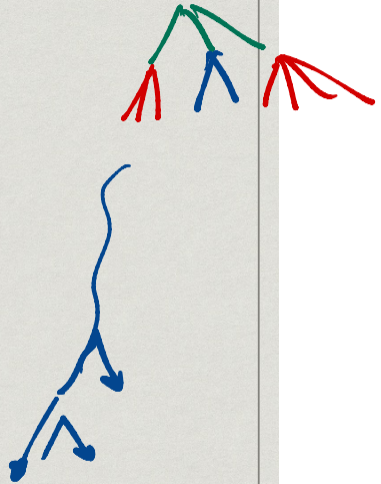
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Programming and Data Structures with Python

Lecture 24, 16 Nov 2023

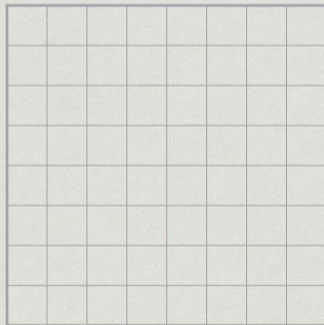
Backtracking

- * Systematically search for a solution
- * Build the solution one step at a time
- * If we hit a dead-end
 - * Undo the last step
 - * Try the next option



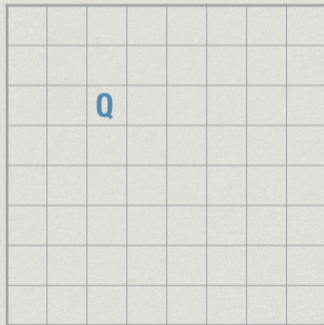
Eight queens

- * Place 8 queens on a chess board so that none of them attack each other
- * In chess, a queen can move any number of squares along a row column or diagonal



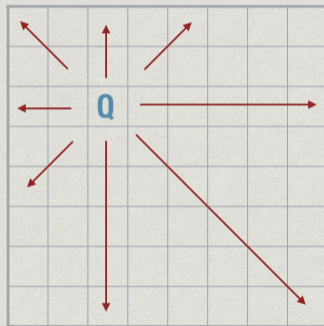
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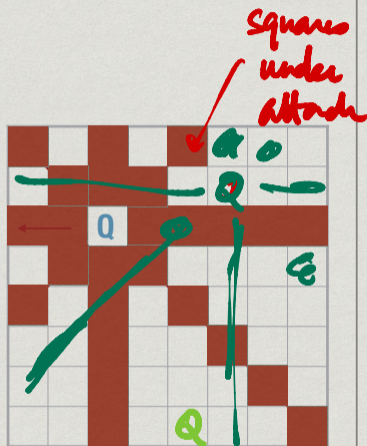
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N queens

- * Place N queens on an N x N chess board so that none attack each other
- * N = 2, 3 impossible



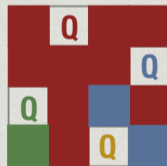
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N queens

- * Place N queens on an N x N chess board so that none attack each other
- * N = 2, 3 impossible
- * N = 4 is possible



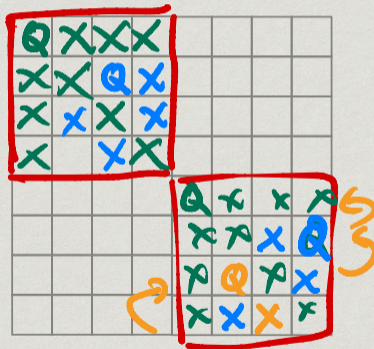
N queens

- * Place N queens on an N x N chess board so that none attack each other
- * N = 2, 3 impossible
- * N = 4 is possible
- * And all bigger N as well



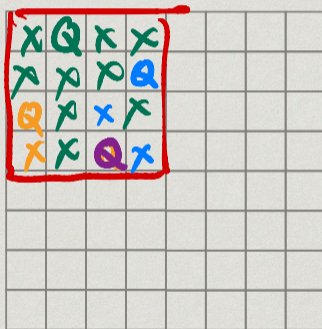
8 queens

- * Clearly, exactly one queen in each row, column
- * Place queens row by row
- * In each row, place a queen in the first available column



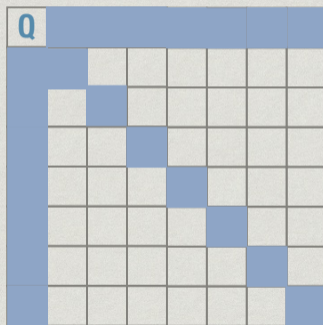
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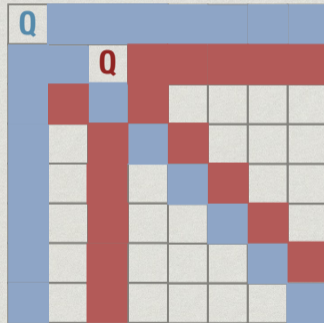
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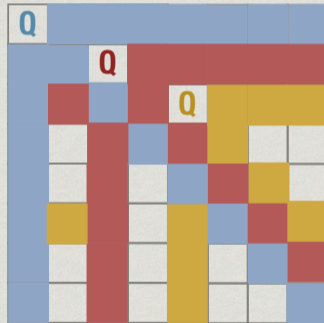
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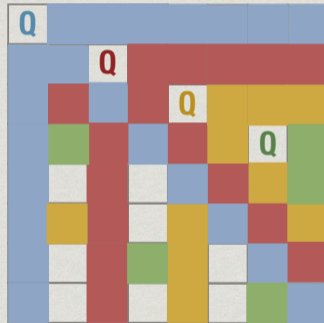
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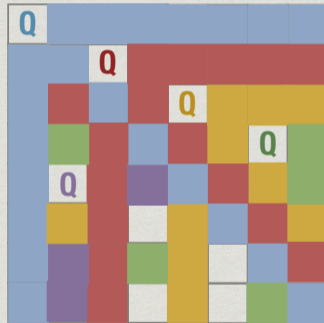
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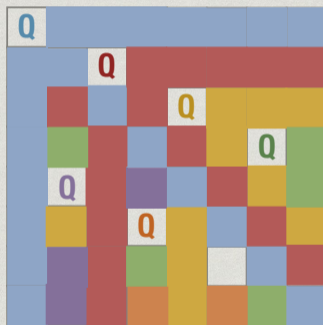
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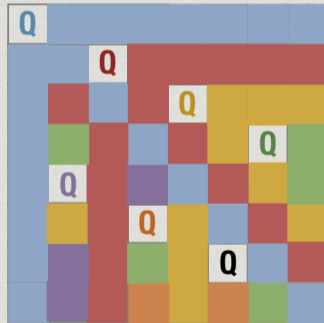
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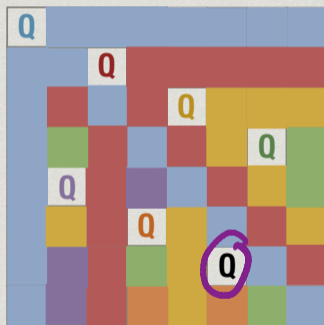
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8 queens

- * Clearly, exactly one queen in each row, column
- * Place queens row by row
- * In each row, place a queen in the first available column
- * Can't place a queen in the 8th row!



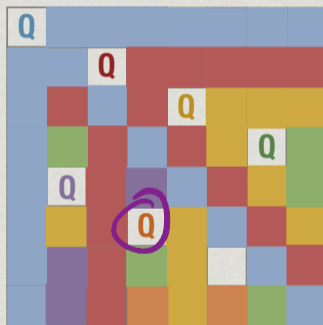
8 queens

- * Can't place the a queen in the 8th row!



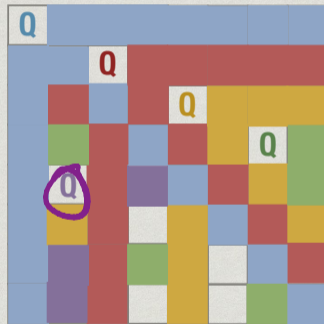
8 queens

- * Can't place the a queen in the 8th row!
- * Undo 7th queen, no other choice



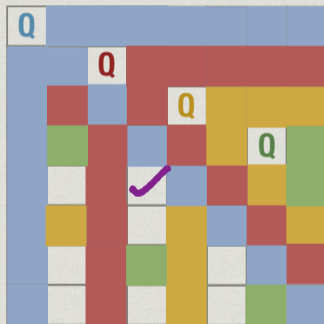
8 queens

- * Can't place the a queen in the 8th row!
- * Undo 7th queen, no other choice
- * Undo 6th queen, no other choice



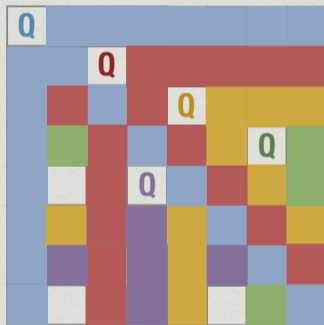
8 queens

- * Can't place the a queen in the 8th row!
- * Undo 7th queen, no other choice
- * Undo 6th queen, no other choice
- * Undo 5th queen, try next



8 queens

- * Can't place the a queen in the 8th row!
- * Undo 7th queen, no other choice
- * Undo 6th queen, no other choice
- * Undo 5th queen, try next



Backtracking

- * Keep trying to extend the next solution
- * If we cannot, undo previous move and try again
- * Exhaustively search through all possibilities
- * ... but systematically!

Coding the solution

- * How do we represent the board?
- * $n \times n$ grid, number rows and columns from 0 to $n-1$
 - * `board[i][j] == 1` indicates queen at (i, j)
 - * `board[i][j] == 0` indicates no queen
- * We know there is only one queen per row
- * Single list `board` of length n with entries 0 to $n-1$
 - * `board[i] == j` : queen in row i , column j , i.e. (i, j)

Overall structure

next queen to place ← current board

```
def placequeen(i,board): # Trying row i
    for each c such that (i,c) is available:
        place queen at (i,c) and update board
        if i == n-1:
            return(True) # Last queen has been placed
        else:
            extendsoln = placequeen(i+1,board)
            if extendsoln:
                return(True) # This solution extends fully
            else:
                undo this move and update board
        else:
            return(False) # Row i failed
```

Updating the board

- * Our 1-D and 2-D representations keep track of the queens
- * Need an efficient way to compute which squares are free to place the next queen
- * $n \times n$ `attack` grid
 - * `attack[i][j] == 1` if (i, j) is attacked by a queen
 - * `attack[i][j] == 0` if (i, j) is currently available
- * How do we undo the effect of placing a queen?
 - * Which `attack[i][j]` should be reset to 0?

Updating the board

-1 = free

k in $0..n-1$ = first attacked by k

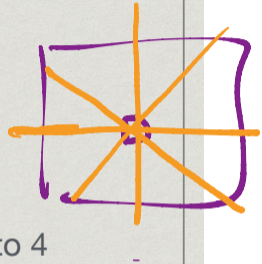
- * Queens are added row by row
- * Number the queens 0 to $n-1$
- * Record earliest queen that attacks each square
 - * $\text{attack}[i][j] == k$ if (i, j) was first attacked by queen k
 - * $\text{attack}[i][j] == -1$ if (i, j) is free
- * Remove queen k — reset $\text{attack}[i][j] == k$ to -1
- * All other squares still attacked by earlier queens

Updating the board

- * `attack` requires n^2 space
 - * Each update only requires $O(n)$ time
 - * Only need to scan row, column, two diagonals
- * Can we improve our representation to use only $O(n)$ space?

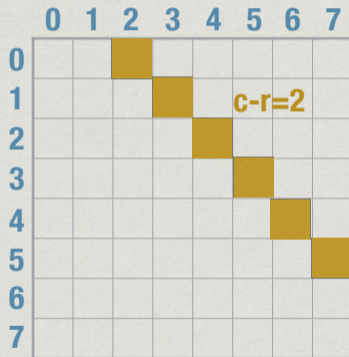
A better representation

- * How many queens attack row i ?
- * How many queens attack row j ?
- * An individual square (i,j) is attacked by upto 4 queens
 - * Queen on row i and on column j
 - * One queen on each diagonal through (i,j)



Numbering diagonals

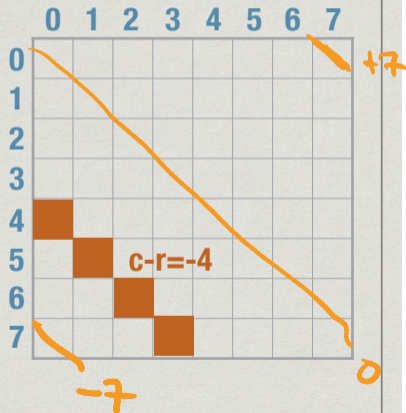
- * Decreasing diagonal:
column - row is invariant



Numbering diagonals

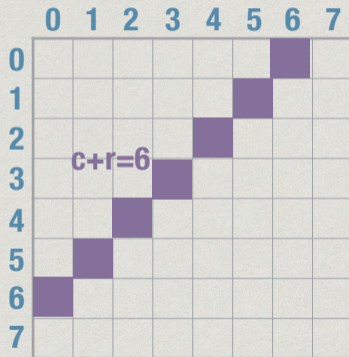
$$2n-1$$

- * Decreasing diagonal:
column - row is invariant



Numbering diagonals

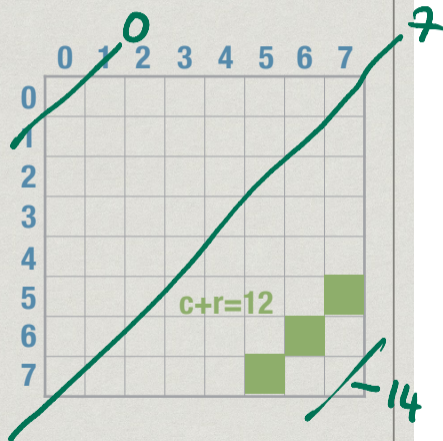
- * Decreasing diagonal:
column - row is invariant
- * Increasing diagonal:
column + row is invariant



Numbering diagonals

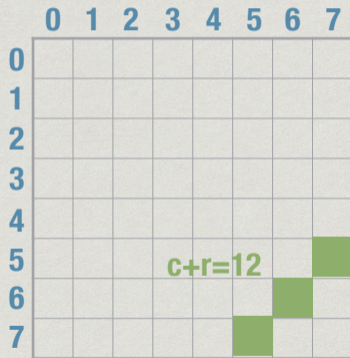
- * Decreasing diagonal:
column - row is invariant
- * Increasing diagonal:
column + row is invariant

$n \times n$ board
 $O(n)$ rows + cols
+ diagonals



Numbering diagonals

- * Decreasing diagonal:
column - row is invariant
- * Increasing diagonal:
column + row is invariant
- * (i,j) is attacked if
 - * row i is attacked
 - * column j is attacked
 - * diagonal $j-i$ is attacked
 - * diagonal $j+i$ is attacked



O(n) representation

- * $row[i] == 1$ if row i is attacked, $0..N-1$
- * $col[i] == 1$ if column i is attacked, $0..N-1$
- * $NWtoSE[i] == 1$ if NW to SE diagonal i is attacked, $-(N-1)$ to $(N-1)$
- * $SWtoNW[i] == 1$ if SW to NE diagonal i is attacked, 0 to $2(N-1)$

Updating the board

- * (i, j) is free if

$row[i] == col[j] == NWtoSE[j-i] == SWtoNE[j+i] == 0$

- * Add queen at (i, j)

$board[i] = j$

$(row[i], col[j], NWtoSE[j-i], SWtoNE[j+i]) =$
 $(1, 1, 1, 1)$

- * Remove queen at (i, j)

$board[i] = -1$

$(row[i], col[j], NWtoSE[j-i], SWtoNE[j+i]) =$
 $(0, 0, 0, 0)$

Implementation details

- * Maintain `board` as nested dictionary
 - * `board['queen'][i] = j` : Queen located at (i, j)
 - * `board['row'][i] = 1` : Row i attacked
 - * `board['col'][i] = 1` : Column i attacked
 - * `board['nwtose'][i] = 1` : NWtoSW diagonal i attacked
 $\hookrightarrow (n-1)$ to $n-1$
 - * `board['swtone'][i] = 1` : SWtoNE diagonal i attacked
 0 to $2(n-1)$

Overall structure

```
def placequeen(i,board): # Trying row i
    for each c such that (i,c) is available:
        place queen at (i,c) and update board
        if i == n-1:
            return(True) # Last queen has been placed
        else:
            extendsoln = placequeen(i+1,board)
            if extendsoln:
                return(True) # This solution extends fully
            else:
                undo this move and update board
    else:
        return(False) # Row i failed
```

All solutions?

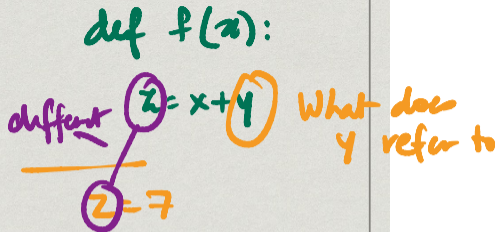


```
def placequeen(i,board): # Try row i
    for each c such that (i,c) is available:
        place queen at (i,c) and update board
    if i == n-1:
        record solution # Last queen placed
    else:
        extendsoln = placequeen(i+1,board)
    undo this move and update board
```

Global variables

- * Can we avoid passing `board` explicitly to each function?
- * Can we have a single `global` copy of `board` that all functions can update?

Scope of name



- * Scope of name is the portion of code where it is available to read and update
- * By default, in Python, scope is local to functions
- * But actually, only if we update the name inside the function

Two examples

```
def f():  
    y = x  
    print(y)
```

— implicitly
global

```
x = 7  
f()
```

} "main" code

Fine!

Two examples

```
def f():  
    y = x  
    print(y)
```

```
x = 7  
f()
```

Fine!

```
def f():  
    y = x  
    print(y)
```

```
x = 22
```

```
x = 7  
f()
```

Error!

Error - undefined value

x becomes local
in all of f()

Two examples

```
def f():  
    y = x  
    print(y)
```

```
x = 7  
f()
```

7

Fine!

```
def f():  
    y = x  
    print(y)
```

```
x = 22
```

```
x = 7
```

```
f()
```

✗

Error!

- * If `x` is not found in `f()`, Python looks at enclosing function for **global** `x`
- * If `x` is updated in `f()`, it becomes a **local** name!

Global variables

- * Actually, this applies only to immutable values
- * Global names that point to mutable values can be updated within a function

```
def f():  
    y = x[0]  
    print(y)  
    x[0] = 22
```

```
x = [7]  
f()
```

$x = [22] + x[1:]$
makes x
local

✓ OK


Fine!

Global immutable values

- * What if we want a global integer
- * Count the number of times a function is called
- * Declare a name to be `global`

```
def f():  
    global x  
    y = x  
    print(y)  
x = 22  
x = 7  
f()  
print(x)
```

```
countf = 0  
def f():  
    global countf  
    countf = countf + 1
```

 Without this line
x is implicitly
global

Global immutable values

- * What if we want a global integer
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```
def f():  
    global x  
    y = x  
    print(y)  
    x = 22
```

```
x = 7  
f()  
print(x)      22
```

Nest function definitions

- * Can define local "helper" functions
- * `g()` and `h()` are only visible to `f()`
- * Cannot be called directly from outside

```
def f():  
    def g(a):  
        return(a+1)  
  
    def h(b):  
        return(2*b+2)  
  
    global x  
    y = g(x) + h(x)  
    print(y)  
    x = 22  
  
x = 7  
f()
```

`g() =`
`h() -`

def f()

def g(a):

global x

x = x + a ← local to g

def h(b)

2 * b + x

global x

$y = g(x) + h(x)$
no change
x changes

Nest function definitions

- * If we look up `x`, `y` inside `g()` or `h()` it will first look in `f()`, then outside
- * Can also declare names global inside `g()`, `h()`
- * Intermediate scope declaration: `nonlocal`
- * See Python documentation

```
def f():  
    def g(a):  
        return(a+1)
```

```
    def h(b):  
        return(2*b)
```

```
    global x
```

```
22 y = g(x) + h(x) → 7 → 8 → 7 → 14  
    print(y)  
    x = 22
```

```
x = 7
```

```
f() 22
```