You shall receive feedback on these problems only if:

- 1. You submit to one of the TAs by 2359 hrs on Thursday, August 15, 2019, and
- 2. Submit each problem in a separate sheet with your name on each sheet. This is essential because the TAs divide correction duties by problem.

"Appearances are not reality; but they often can be a convincing alternative to it." – Dean Koontz

1. For a word  $w \in \Sigma^*$ ,  $w = a_1 \cdots a_n, a_1, \ldots, a_n \in \Sigma$ , define the *reverse* word  $w^R = a_n \cdots a_1$ . For a language  $L \subseteq \Sigma^*$ , define the reverse language  $L^R$  of L as

$$L^R = \{ w^R \mid w \in L \}$$

Consider some L that is recognizable. Show that  $L^R$  is recognizable.

2. We shall define the *Kleene star* operation. For a language L,

$$L^* = \{ w_1 \cdots w_n \mid n \ge 0, w_1, \dots, w_n \in L \}$$

If L is recognizable, show that  $L^*$  is recognizable.

- 3. Show that recognizable languages are closed under :
  - (a) two-stutter, where for a language L,

two-stutter(L) = 
$$\{a_1a_1\cdots a_na_n \mid a_1\cdots a_n \in L\}$$

(b) stutter, where for a language L,

stutter(L) = 
$$\{a_1^{m_1} \cdots a_n^{m_n} \mid a_1 \cdots a_n \in L, m_1, \dots, m_n > 0\}$$

4. Show by giving an example that, if A is an NFA that recognizes language L, swapping the accept and non-accept states in A doesn't necessarily yield an NFA that recognizes the complement of L.