

Unit-7: Linear Temporal Logic

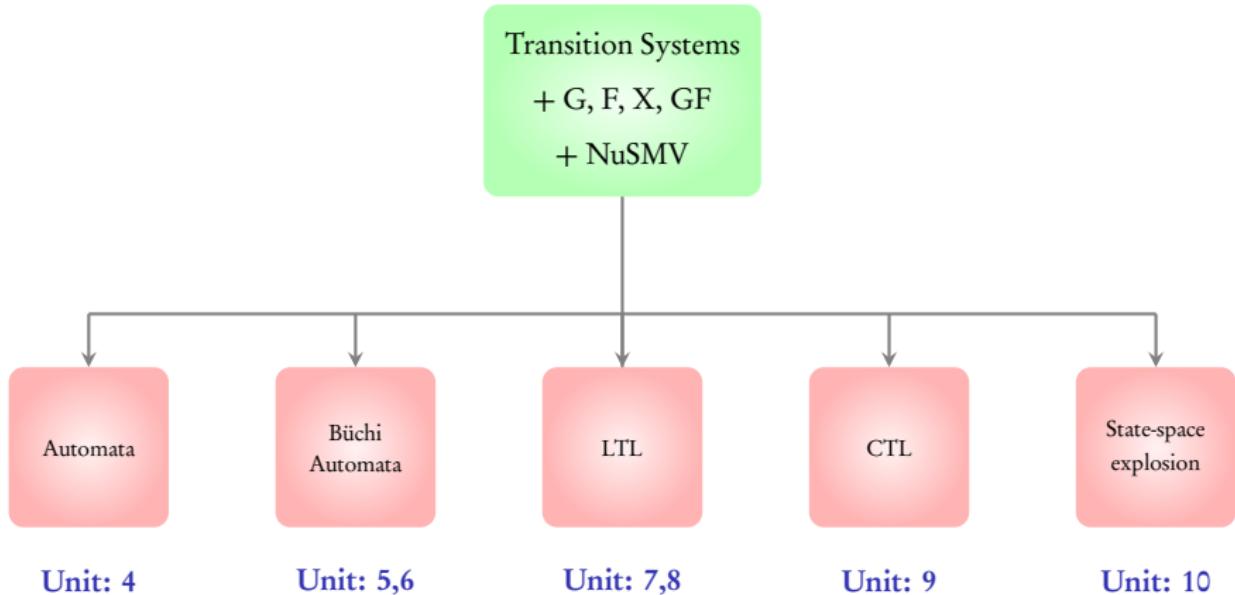
B. Srivathsan

Chennai Mathematical Institute

NPTEL-course

July - November 2015

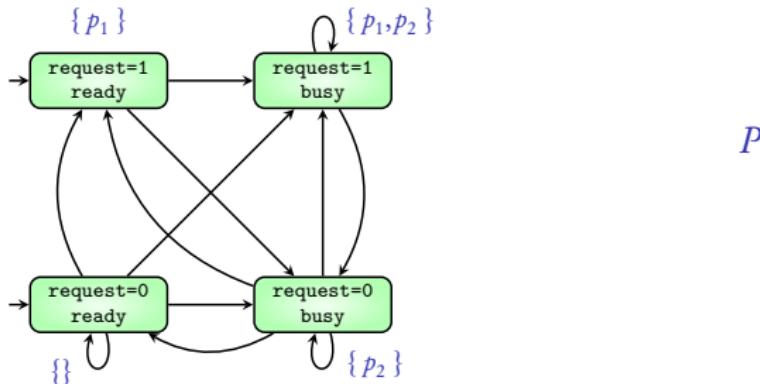
Module 1: Introduction to LTL



$$\text{AP} = \{ p_1, p_2 \}$$

Transition System

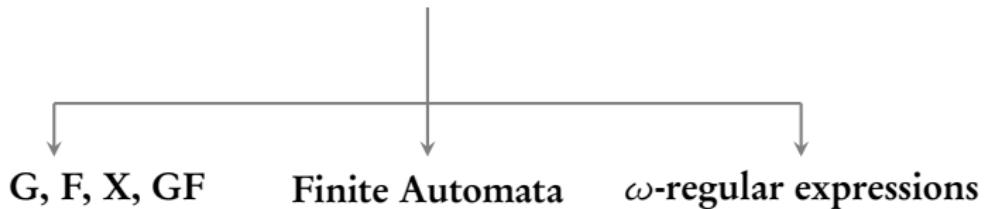
Property



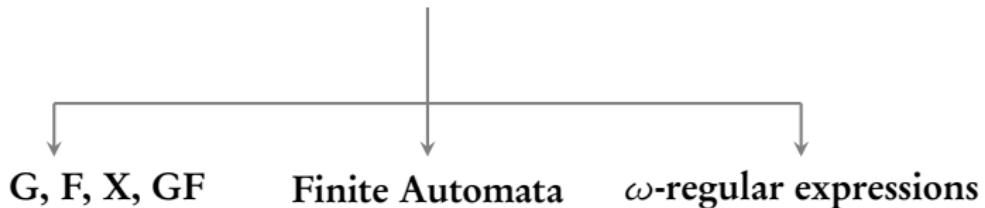
Transition system TS satisfies property P if

$\text{Traces}(TS) \subseteq P$

Specifying properties



Specifying properties



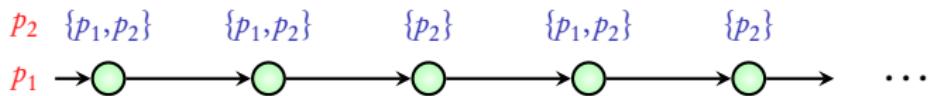
Here: Another formalism - Linear Temporal Logic



$\phi :=$

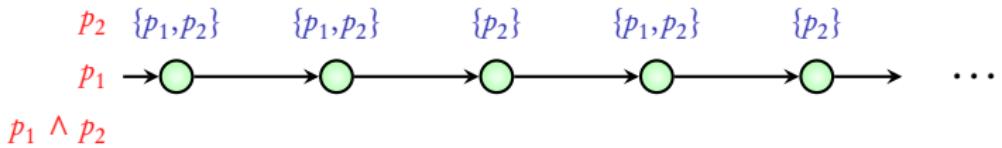


$\phi := \text{true} |$



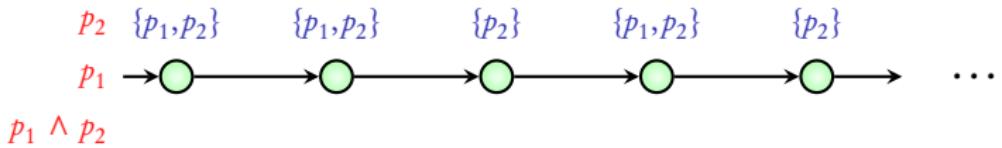
$\phi := \text{true} \mid p_i \mid$

$p_i \in AP$



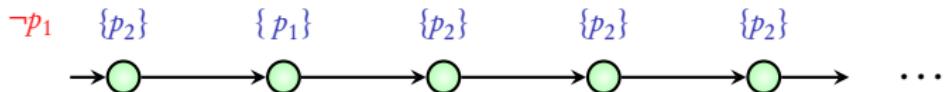
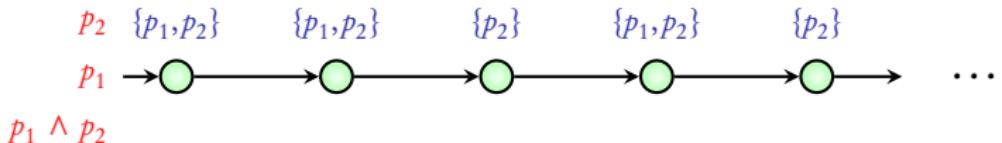
$\phi := \text{true} \mid p_i \mid \phi_1 \wedge \phi_2 \mid$

$p_i \in AP \quad \phi_1, \phi_2 : \text{LTL formulas}$



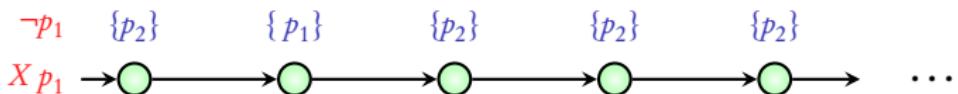
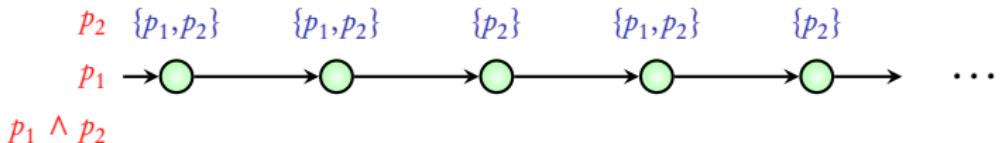
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$p_i \in AP$ ϕ_1, ϕ_2 : LTL formulas



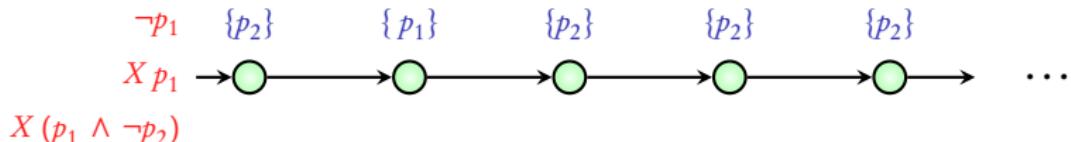
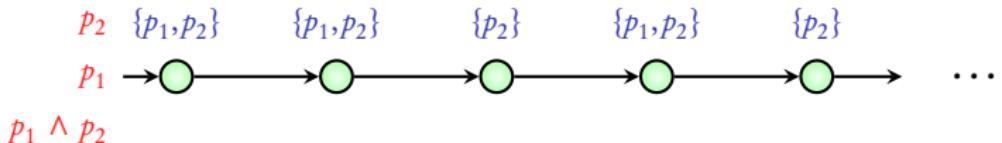
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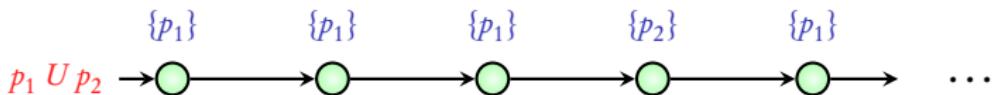
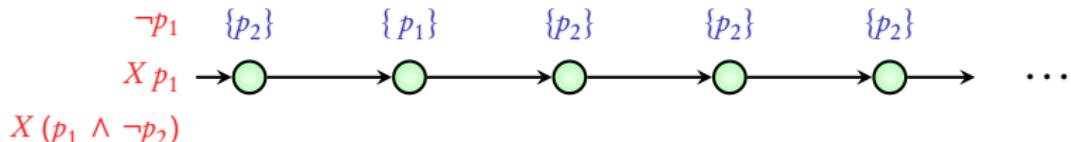
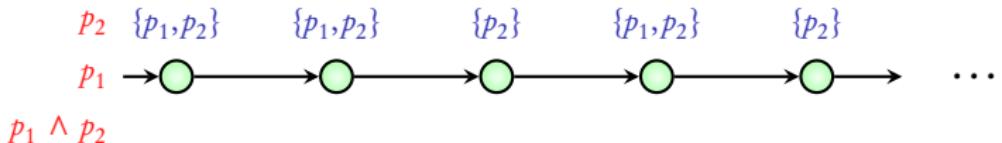
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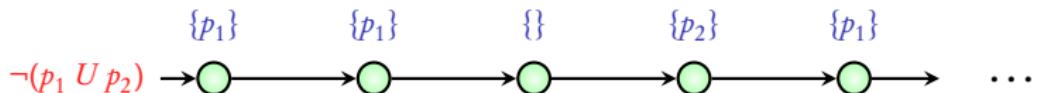
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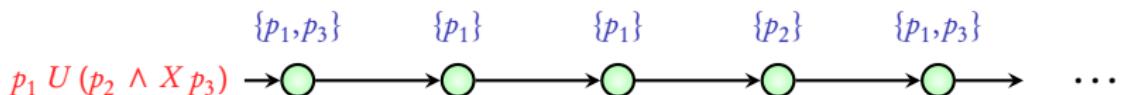
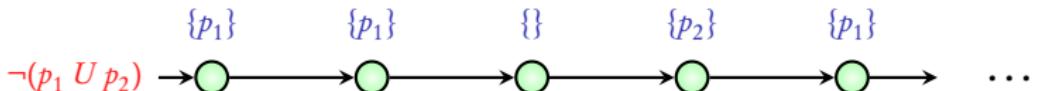


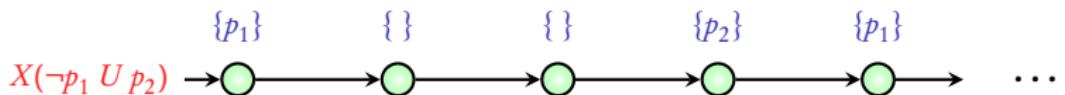
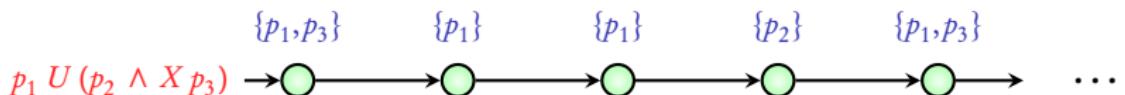
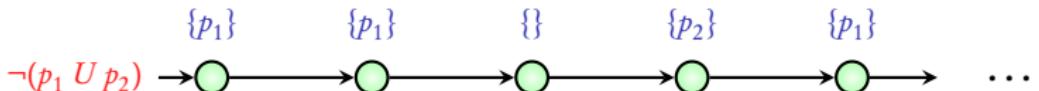
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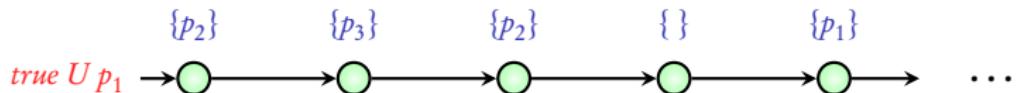
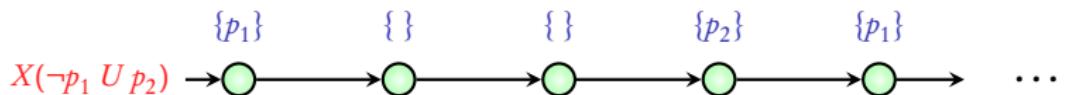
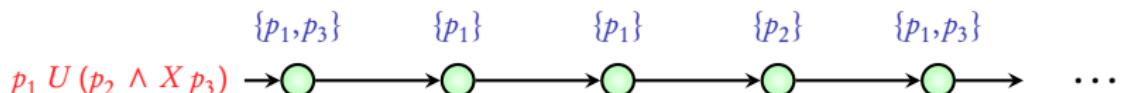
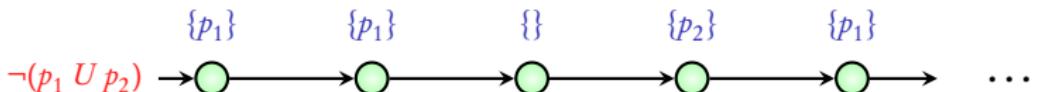
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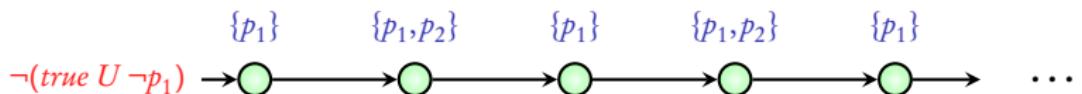
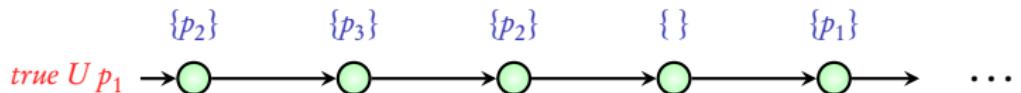
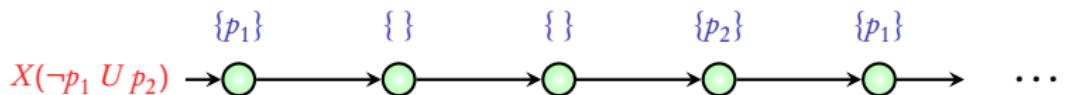
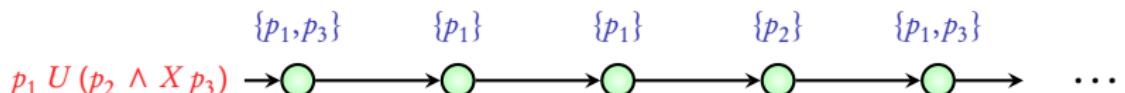
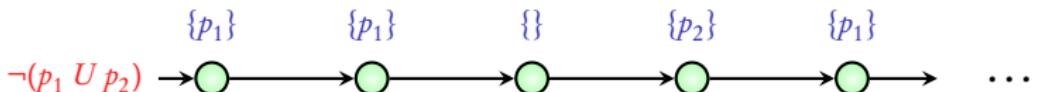
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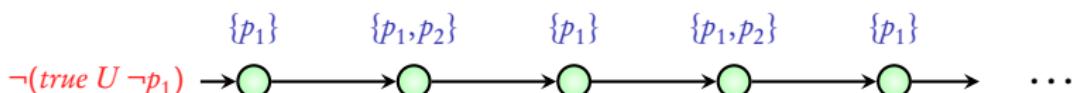
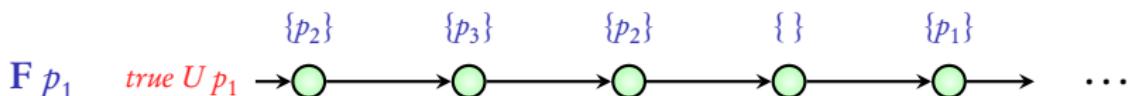
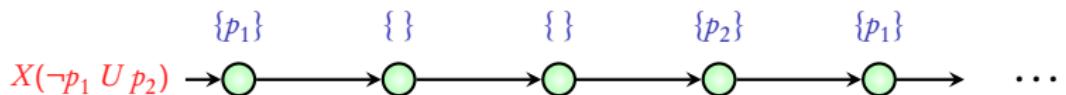
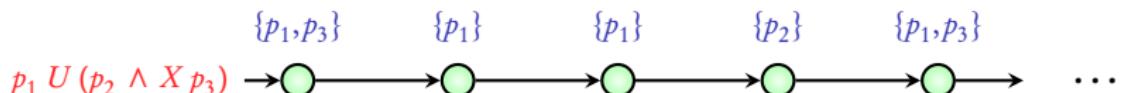
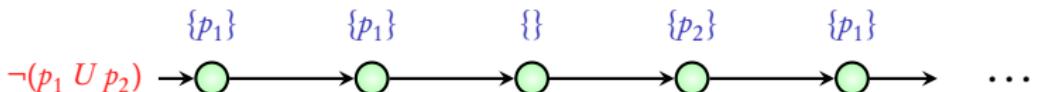
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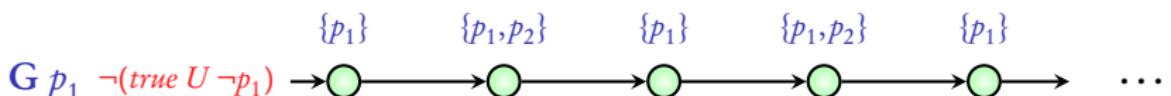
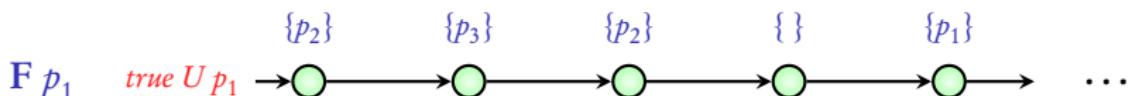
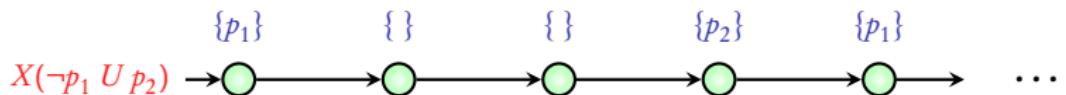
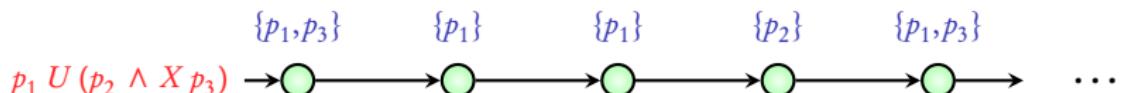
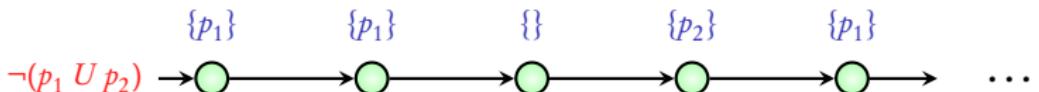
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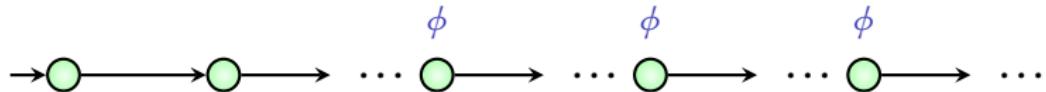
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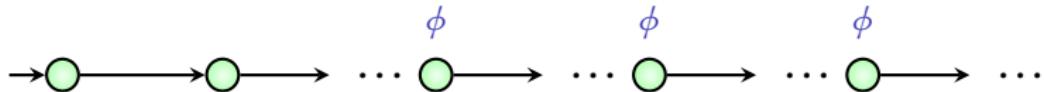
Derived operators

- $\phi_1 \vee \phi_2$: $\neg(\neg\phi_1 \wedge \neg\phi_2)$ (Or)
- $\phi_1 \rightarrow \phi_2$: $\neg\phi_1 \vee \phi_2$ (Implies)
- $F \phi$: true U ϕ (Eventually)
- $G \phi$: $\neg F \neg\phi$ (Always)

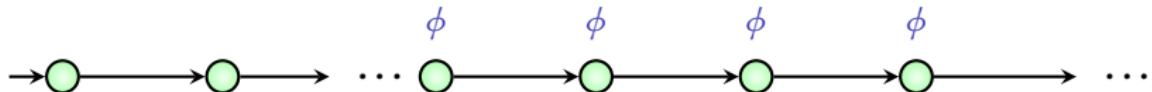
$\text{G F } \phi$ (Infinitely often)



G F ϕ (Infinitely often)



$\text{F G } \phi$ (Eventually forever)



Coming next: More examples

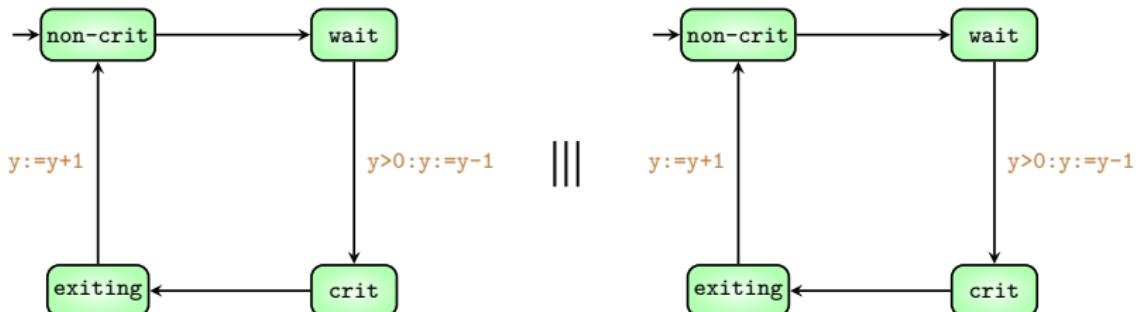
Atomic propositions $AP = \{ crit_1, wait_1, crit_2, wait_2 \}$

$crit_1$: $pr1.location=crit$

$wait_1$: $pr1.location=wait$

$crit_2$: $pr2.location=crit$

$wait_2$: $pr2.location=wait$



- ▶ **Safety:** both processes cannot be in critical section simultaneously

$$G (\neg crit_1 \vee \neg crit_2)$$

- ▶ **Liveness:** each process visits critical section infinitely often

$$G F crit_1 \wedge G F crit_2$$

Summary

$$\phi := \text{true} \mid p_i \mid \phi_1 \wedge \phi_2 \mid \neg \phi_1 \mid X \phi \mid \phi_1 U \phi_2$$

$F \phi$: true U ϕ

(Eventually)

$G \phi$: $\neg F \neg \phi$

(Always)