

Week-1: Introduction to model checking

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Chennai Mathematical Institute

NPTEL-course

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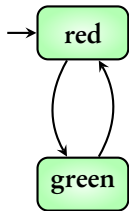
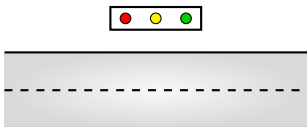
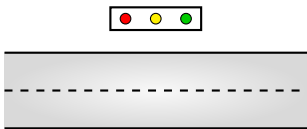
Module 4:
Modeling concurrent systems

Concurrent systems

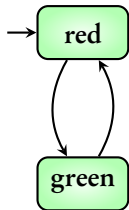
Independent

Shared variables

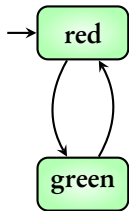
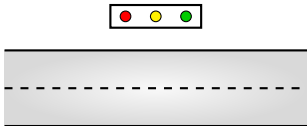
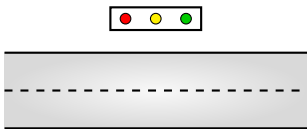
Shared actions



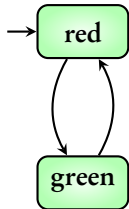
$TrLight_1$



$TrLight_2$

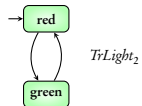
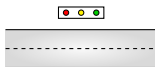
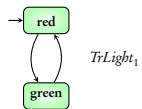


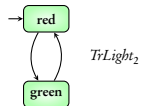
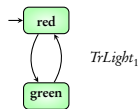
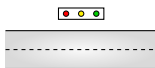
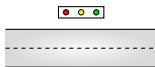
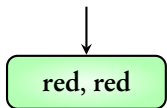
$TrLight_1$

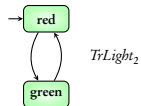
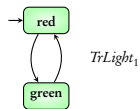
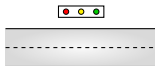
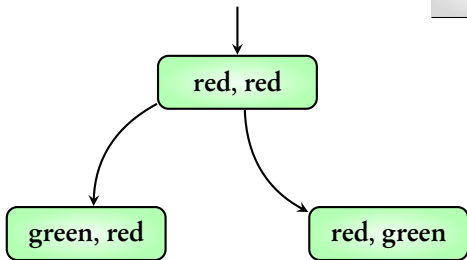


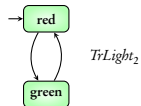
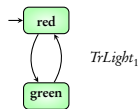
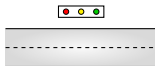
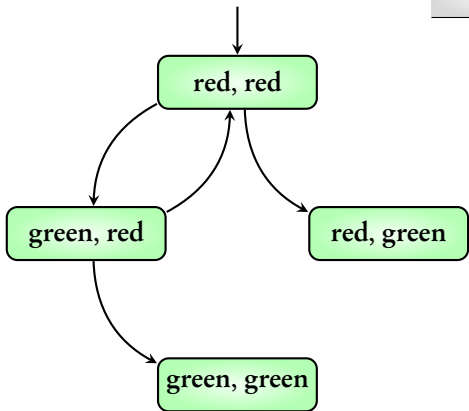
$TrLight_2$

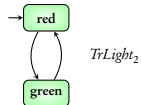
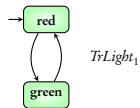
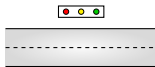
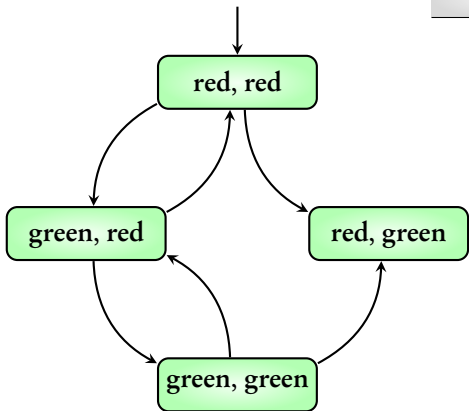
What is the transition system for the **joint behaviour**?

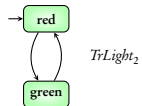
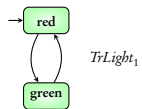
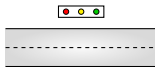
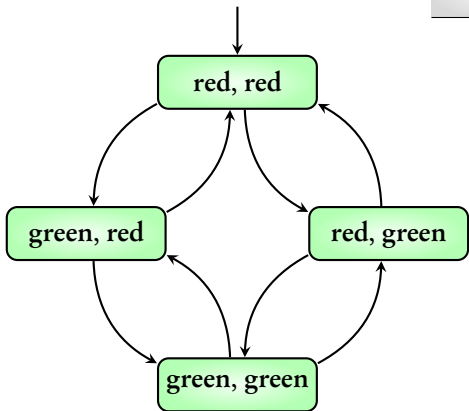


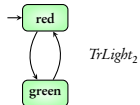
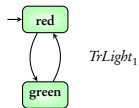
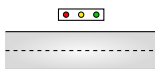
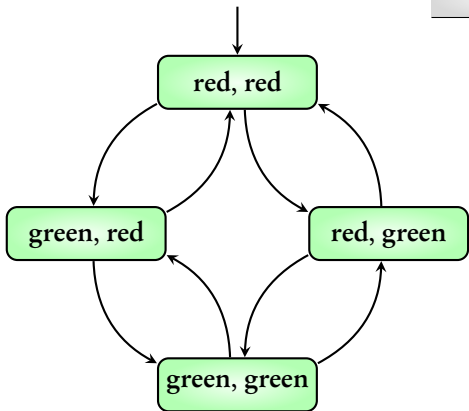




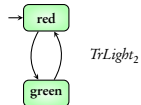
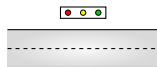
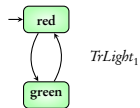
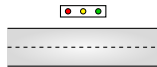




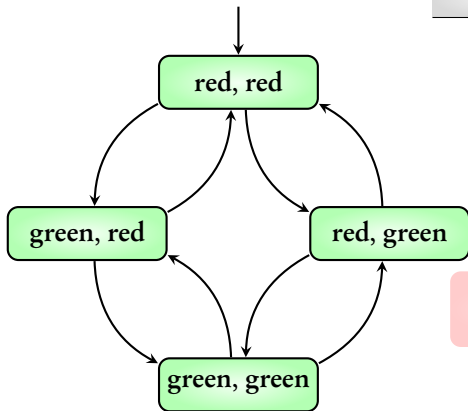




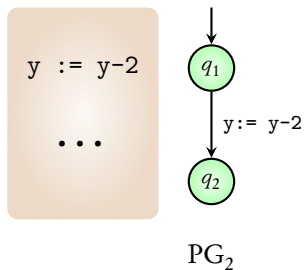
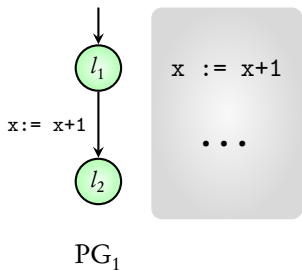
$TrLight_1 \equiv TrLight_2$

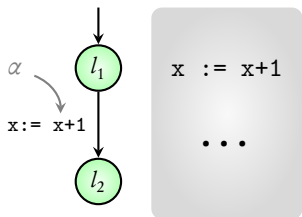


$TrLight_1 \parallel TrLight_2$

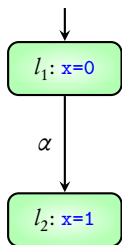


\parallel : Interleaving operator



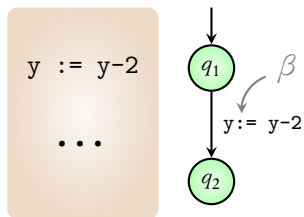


PG₁

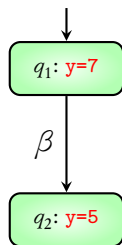


TS₁

(initially x=0)

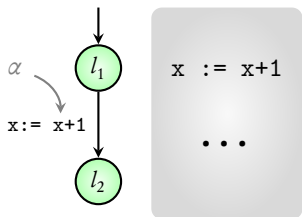


PG₂

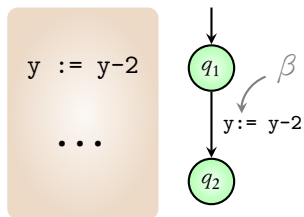


TS₂

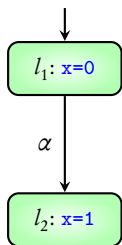
(initially y=7)



PG₁

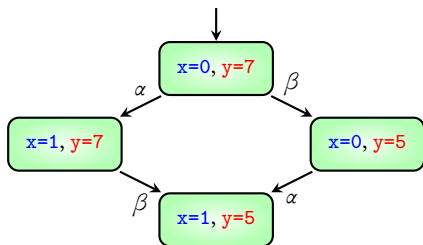


PG₂

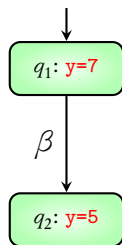


TS₁

(initially x=0)

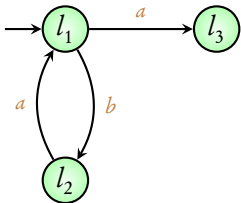


TS₁ ||| TS₂

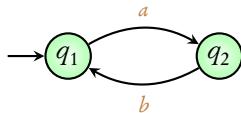


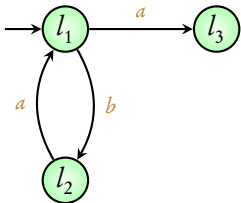
TS₂

(initially y=7)

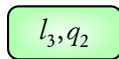
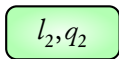
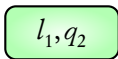
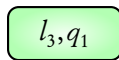
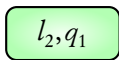
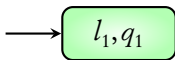
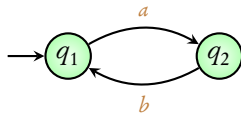


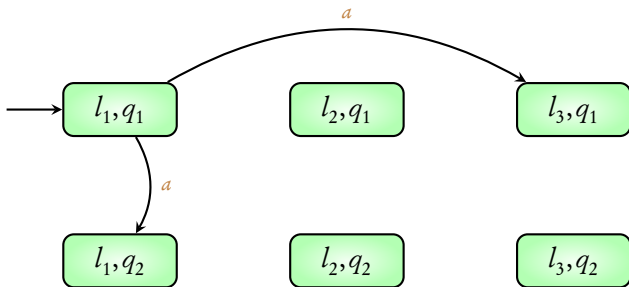
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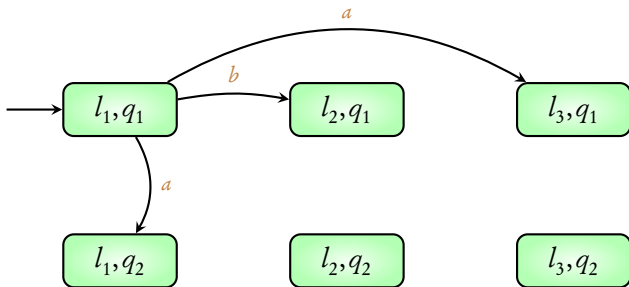


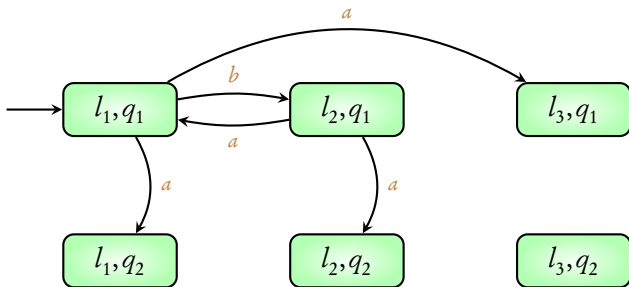
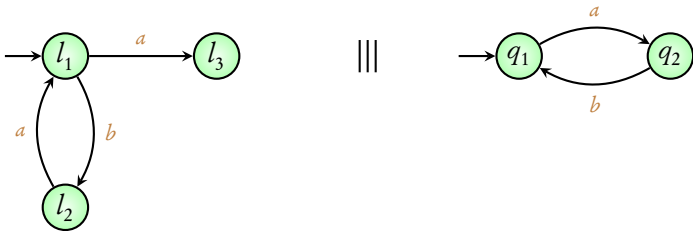


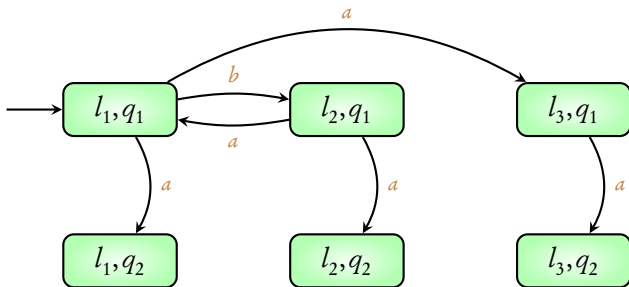
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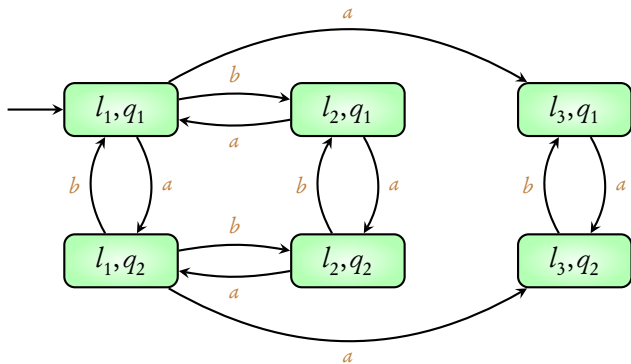












Multiple systems

$TS_1 \parallel TS_2 \parallel \dots \parallel TS_n$

Multiple systems

$TS_1 \parallel TS_2 \parallel \dots \parallel TS_n$

Exercise: Try out an example of interleaving **three** systems

Concurrent systems

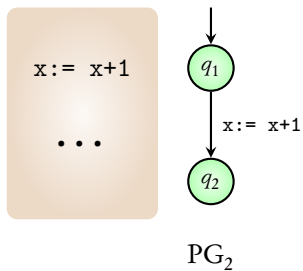
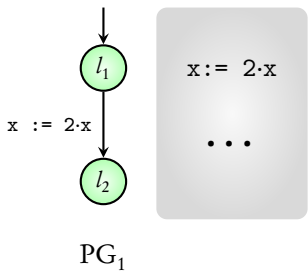
Independent

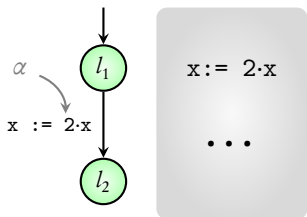
Interleaving

$TS_1 \parallel TS_2 \parallel \dots \parallel TS_n$

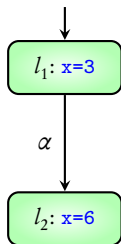
Shared variables

Shared actions



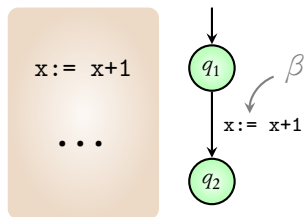


PG₁

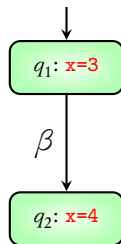


TS₁

(initially x=3)

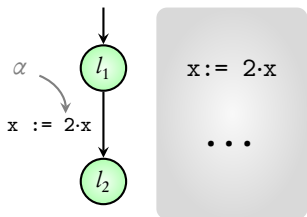


PG₂

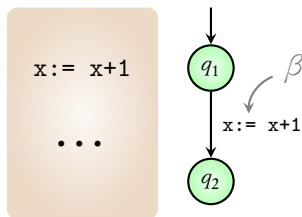


TS₂

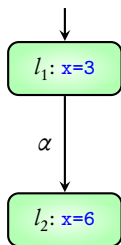
(initially x=3)



PG₁

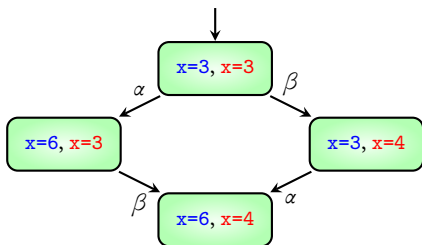


PG₂

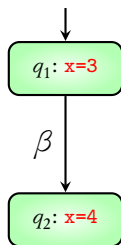


TS₁

(initially x=3)

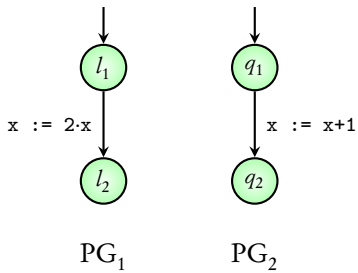


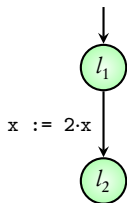
TS₁ ||| TS₂



TS₂

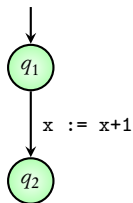
(initially x=3)





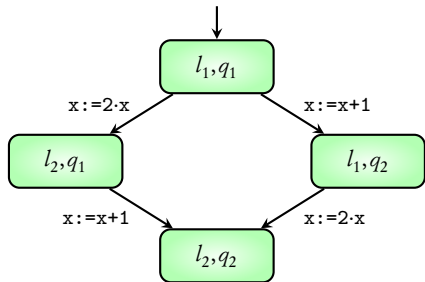
PG₁

|||

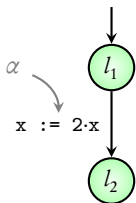


PG₂

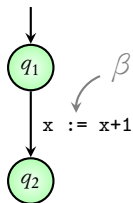
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PG₁ ||| PG₂



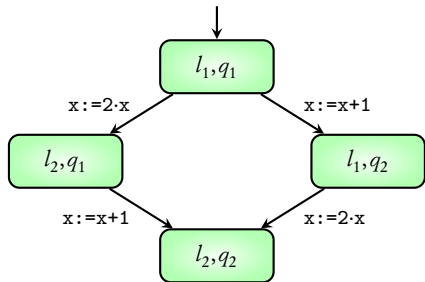
PG_1



PG_2

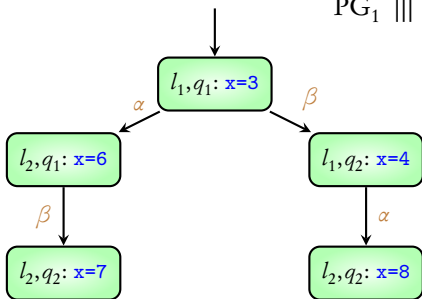
\equiv

$=$



$PG_1 \parallel PG_2$

$TS(PG_1 \parallel PG_2):$



Concurrent systems

Independent

Interleaving

$TS_1 \parallel TS_2 \parallel \dots \parallel TS_n$

Shared variables

$TS(PG_1 \parallel PG_2 \parallel \dots \parallel PG_n)$

Shared actions

Coming next: Another example

```
while x < 200
```

```
  x := x+1
```

```
while x>0
```

```
  x := x-1
```

```
while x=200
```

```
  x := 0
```

```
while x < 200
```

```
x := x+1
```

```
while x>0
```

```
x := x-1
```

```
while x=200
```

```
x := 0
```

Is the value of x **always** between 0 and 200?

while $x < 200$

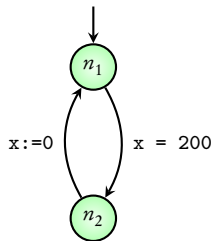
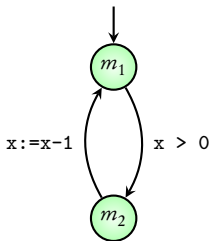
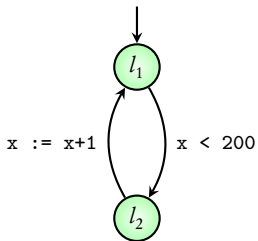
$x := x+1$

while $x > 0$

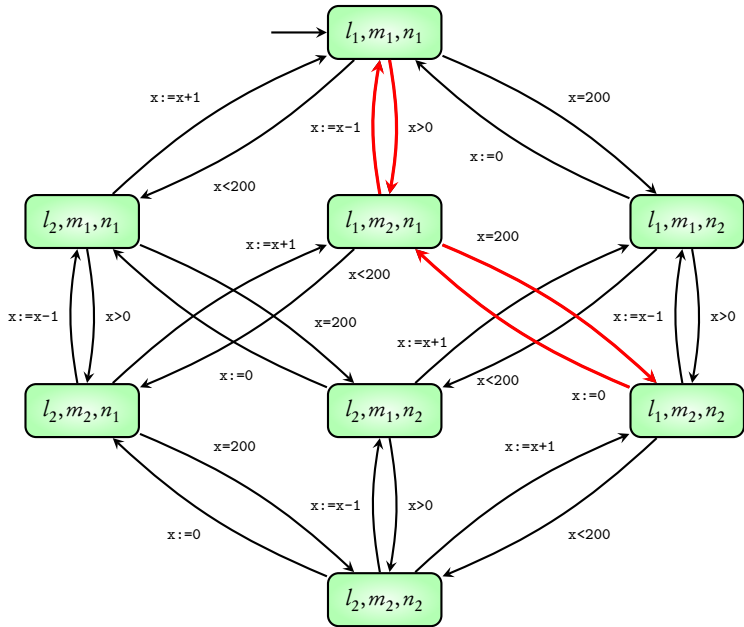
$x := x-1$

while $x=200$

$x := 0$

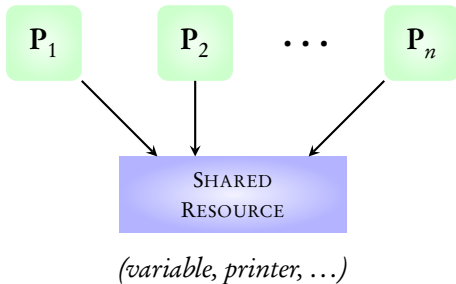


Is the value of x **always** between 0 and 200?



Is the value of x always between 0 and 200? **No**

Coming next: Mutual exclusion



Mutual Exclusion: No two processes can access the resource simultaneously

Goal: Modeling the **protocols** used for mutual exclusion

P₁

loop forever

⋮ *non-critical actions*

request

critical section

release

⋮ *non-critical actions*

end loop

P₂

loop forever

⋮ *non-critical actions*

request

critical section

release

⋮ *non-critical actions*

end loop

P_1

loop forever

 \vdots *non-critical actions**request*

critical section

release \vdots *non-critical actions*

end loop

 P_2

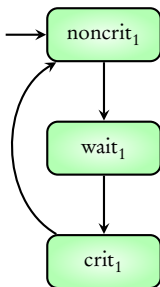
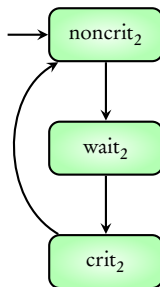
loop forever

 \vdots *non-critical actions**request*

critical section

release \vdots *non-critical actions*

end loop

 PG_1  PG_2 

P_1

loop forever

```

⋮          *non-critical actions*

```

```

⟨ if y>0:  y:=y-1 ⟩  *request*

```

critical section

```

y:=y+1          *release*

```

```

⋮          *non-critical actions*

```

end loop

 P_2

loop forever

```

⋮          *non-critical actions*

```

```

⟨ if y>0:  y:=y-1 ⟩  *request*

```

critical section

```

y:=y+1          *release*

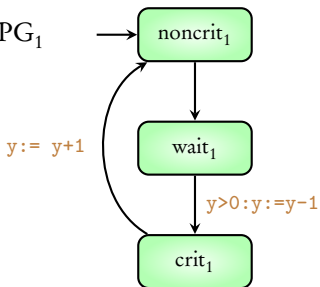
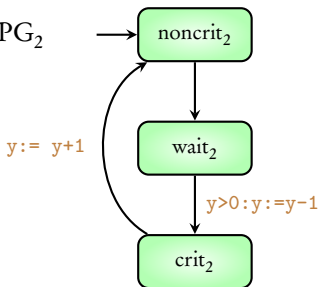
```

```

⋮          *non-critical actions*

```

end loop

PG₁PG₂

P_1

loop forever

```

:      *non-critical actions*
< if y>0: y:=y-1 > *request*
critical section
y:=y+1      *release*
:      *non-critical actions*
end loop

```

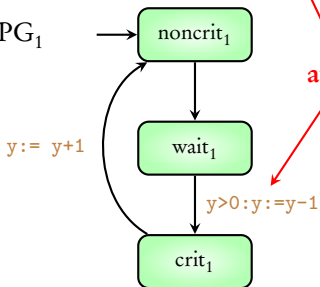
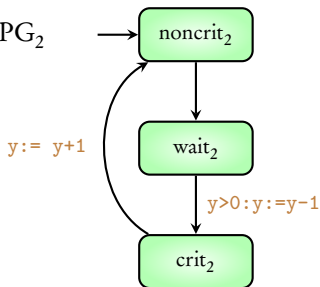
 P_2

loop forever

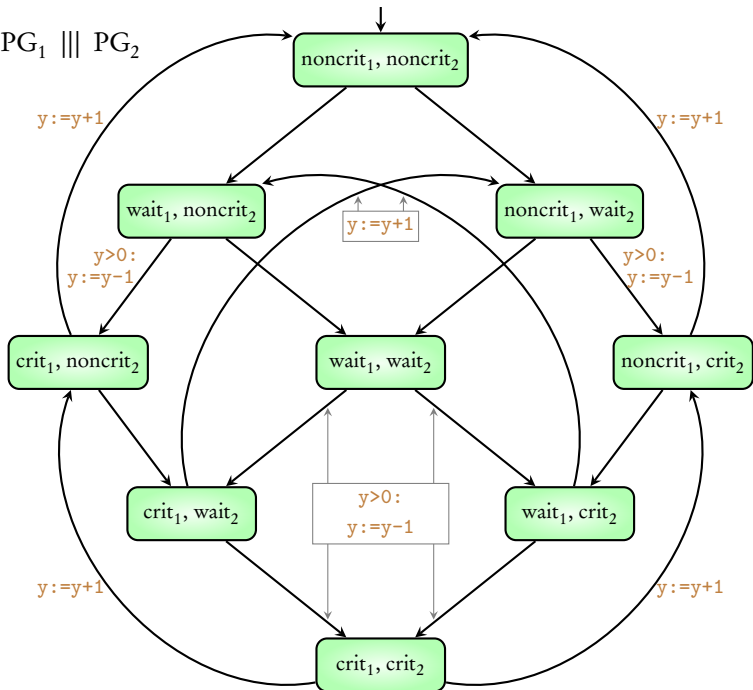
```

:      *non-critical actions*
< if y>0: y:=y-1 > *request*
critical section
y:=y+1      *release*
:      *non-critical actions*
end loop

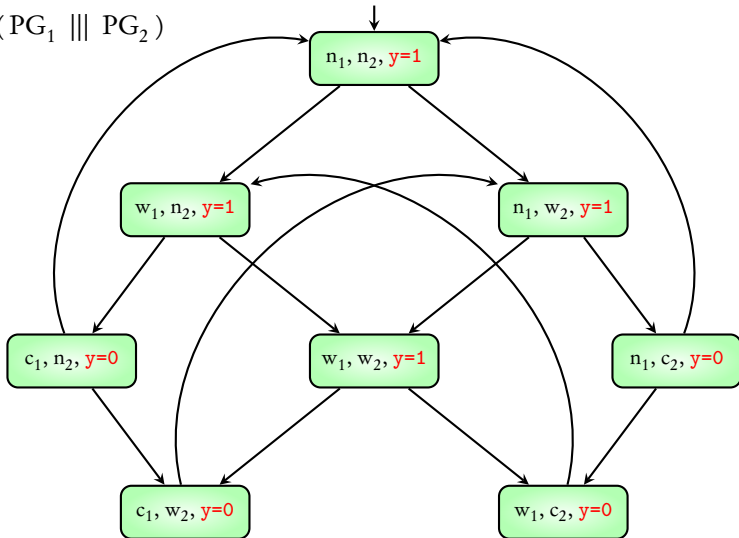
```

PG₁PG₂**atomic**

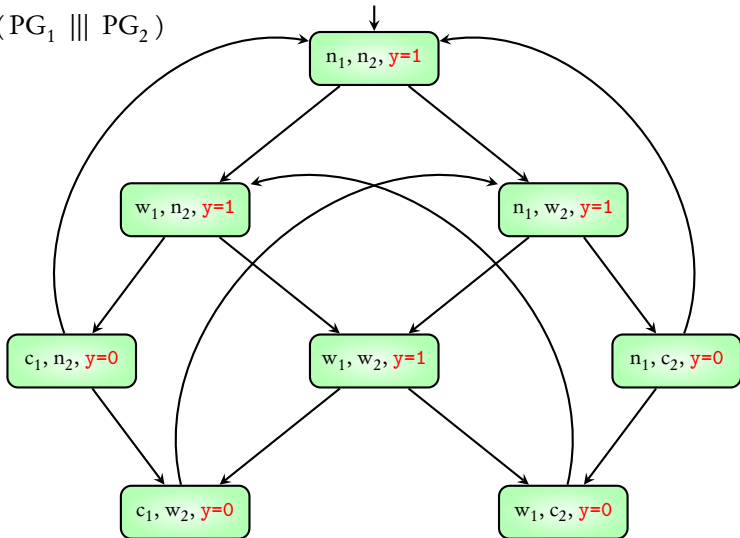
$PG_1 \parallel PG_2$



TS($PG_1 \parallel PG_2$)



TS($PG_1 \parallel PG_2$)



Both processes **cannot be** in critical section **simultaneously**

Concurrent systems

Independent

Interleaving

$TS_1 \parallel TS_2 \parallel \dots \parallel TS_n$

Shared variables

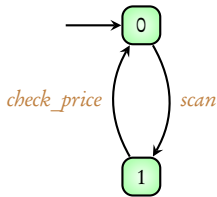
$TS(PG_1 \parallel PG_2 \parallel \dots \parallel PG_n)$

Mutual Exclusion

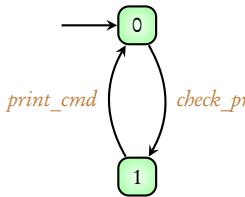
Shared actions

Coming next: Book-keeping system in a supermarket

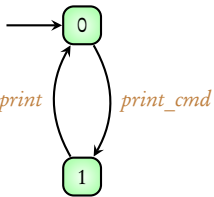
Bar-Code Reader (BCR)



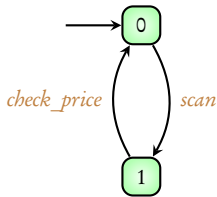
Booking Program (BP)



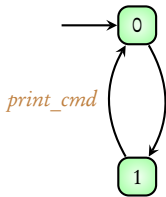
Printer (P)



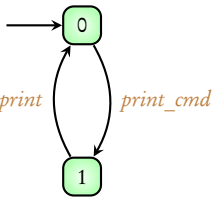
Bar-Code Reader (BCR)



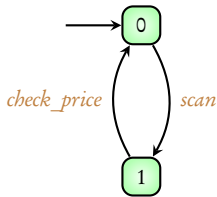
Booking Program (BP)



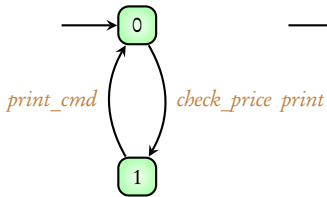
Printer (P)



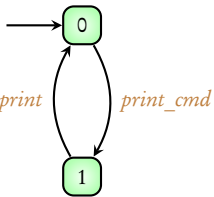
Bar-Code Reader (BCR)



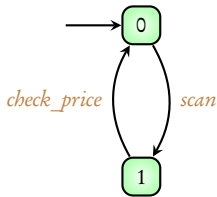
Booking Program (BP)



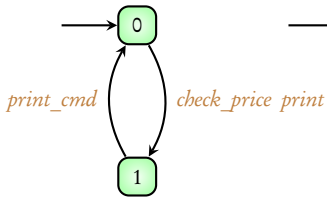
Printer (P)



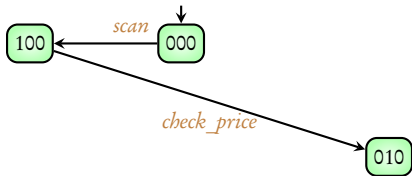
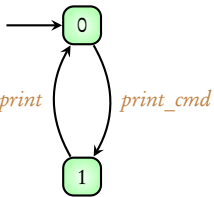
Bar-Code Reader (BCR)



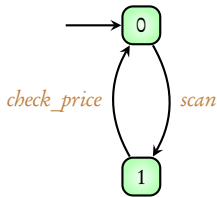
Booking Program (BP)



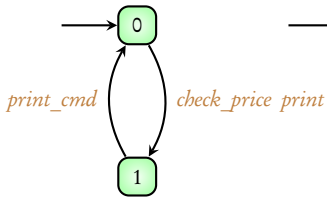
Printer (P)



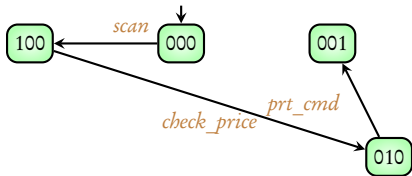
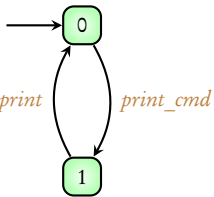
Bar-Code Reader (BCR)



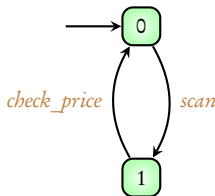
Booking Program (BP)



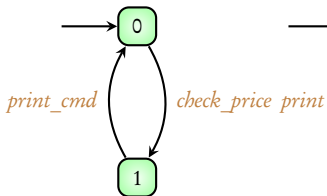
Printer (P)



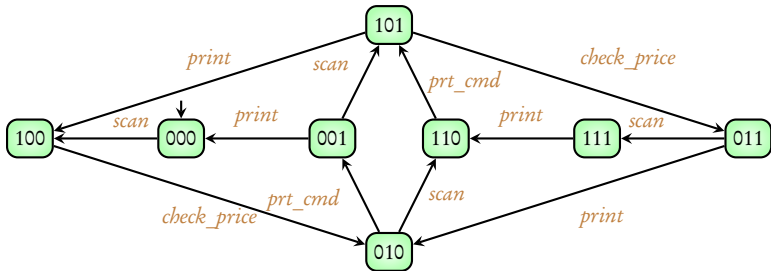
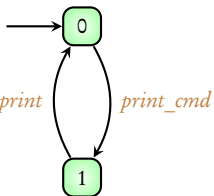
Bar-Code Reader (BCR)



Booking Program (BP)



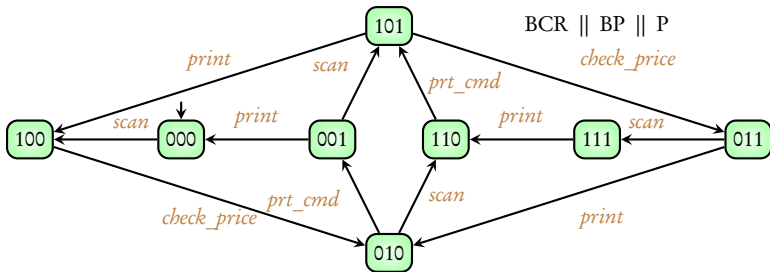
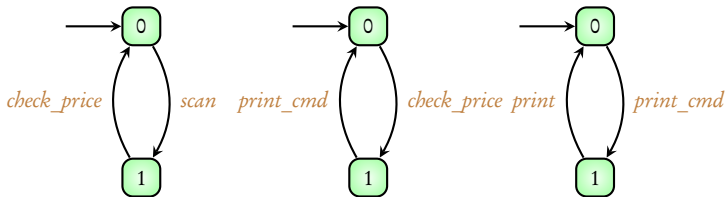
Printer (P)



Bar-Code Reader (BCR)

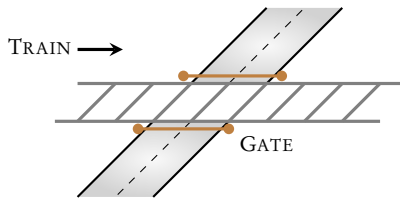
Booking Program (BP)

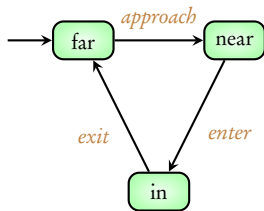
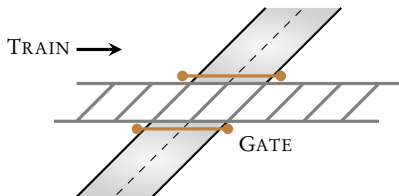
Printer (P)



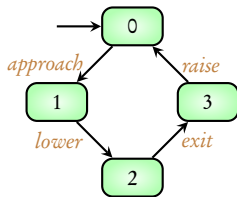
check_price, print_cmd: **Shared actions** (also called **handshaking actions**)

Next example: Train-Gate-Controller

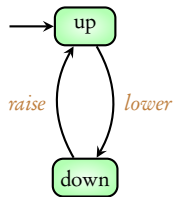




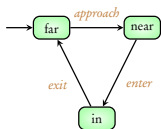
Train



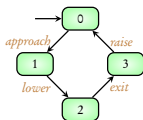
Controller



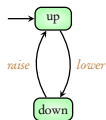
Gate



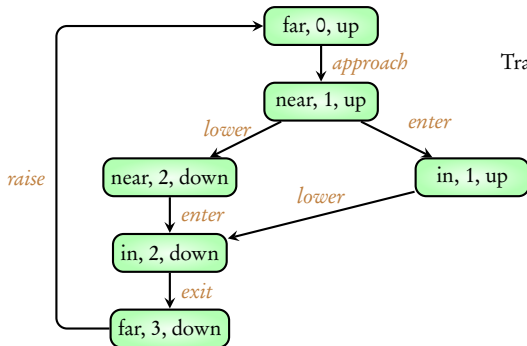
Train



Controller



Gate



Train || Controller || Gate

|| : Handshake operator

Independent

Interleaving

$TS_1 \parallel TS_2 \parallel \dots \parallel TS_n$

Shared variables

$TS(PG_1 \parallel PG_2 \parallel \dots \parallel PG_n)$

Mutual Exclusion

Shared actions

$TS_1 \parallel TS_2$

Reference: Principles of Model Checking, *Baier and Katoen*, MIT Press (2008)
Pages 35 - 53