

# Lecture 9: Algorithms for LTL

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*Model-Checking and Systems Verification*

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Module 1:

**Automata-based LTL  
model-checking**

Does **Transition system** satisfy **LTL formula  $\phi$** ?

Does **Transition system** satisfy **LTl formula**  $\phi$  ?

**Negation**  $\neg \phi$

Does **Transition system** satisfy **LTL formula**  $\phi$ ?

**Negation**  $\neg \phi$



**NBA**  $\mathcal{A}_{\neg \phi}$

Does **Transition system** satisfy **LTL formula  $\phi$** ?



**NBA**  $\mathcal{A}_{T.S}$

**Negation**  $\neg \phi$



**NBA**  $\mathcal{A}_{\neg\phi}$

Does **Transition system** satisfy **LTL formula  $\phi$** ?



**NBA**  $\mathcal{A}_{T.S.}$

**Negation**  $\neg \phi$



**NBA**  $\mathcal{A}_{\neg\phi}$

Is  $L(\mathcal{A}_{T.S.}) \cap L(\mathcal{A}_{\neg\phi})$  empty?

Does **Transition system** satisfy **LTL formula  $\phi$** ?



**NBA**  $\mathcal{A}_{T.S.}$

**Negation**  $\neg \phi$



**NBA**  $\mathcal{A}_{\neg\phi}$

Is  $L(\mathcal{A}_{T.S.}) \cap L(\mathcal{A}_{\neg\phi})$  empty?

Is  $L(\mathcal{A}_{T.S.} \times \mathcal{A}_{\neg\phi})$  empty?



**Here:** Converting LTL formulas to NBA

**Here:** Converting LTL formulas to NBA

**Coming next:** Examples

Atomic propositions  $\mathbf{AP} = \{ p_1, p_2 \}$

*Alphabet:*

$\{ \{ \}, \{ p_1 \}, \{ p_2 \}, \{ p_1, p_2 \} \}$

**F**  $p_1$  Words where  $p_1$  occurs sometime

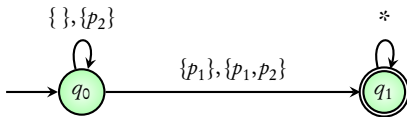
$\{p_2\} \{ \} \{ \} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$

$\{p_1, p_2\} \{ \} \{ \} \{ \} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$

$\vdots$

**F**  $p_1$  Words where  $p_1$  occurs sometime

$\{p_2\} \{ \} \{ \} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$   
 $\{p_1, p_2\} \{ \} \{ \} \{ \} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$   
 $\vdots$



**G**  $p_1$  Words where  $p_1$  occurs always

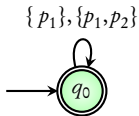
$\{p_1\} \{p_1, p_2\} \{p_1\} \{p_1, p_2\} \{p_1\} \{p_1\} \{p_1\} \dots$

$\{p_1, p_2\} \{p_1, p_2\} \{p_1\} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$

$\vdots$

**G**  $p_1$  Words where  $p_1$  occurs always

$\{p_1\} \{p_1, p_2\} \{p_1\} \{p_1, p_2\} \{p_1\} \{p_1\} \{p_1\} \dots$   
 $\{p_1, p_2\} \{p_1, p_2\} \{p_1\} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$   
 $\vdots$



$p_1 \wedge \neg p_2$  Words starting with  $\{p_1\}$

$\{p_1\} \{\} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$

$\{p_1\} \{\} \{\} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$

$\vdots$

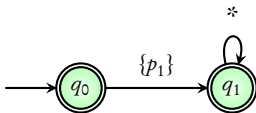


$p_1 \wedge \neg p_2$  Words starting with  $\{p_1\}$

$\{p_1\} \{\} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$

$\{p_1\} \{\} \{\} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$

$\vdots$



$$p_1 \wedge \mathbf{X} \neg p_2$$

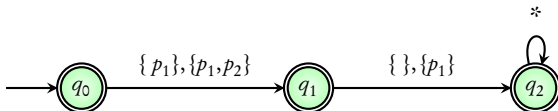
$$\{p_1\} \{\}\{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$$

$$\{p_1, p_2\} \{p_1\} \{\}\{p_1\} \{p_1\} \{p_1, p_2\} \dots$$

⋮

$$p_1 \wedge \mathbf{X} \neg p_2$$

$\{p_1\} \{\} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$   
 $\{p_1, p_2\} \{p_1\} \{\} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$   
 $\vdots$

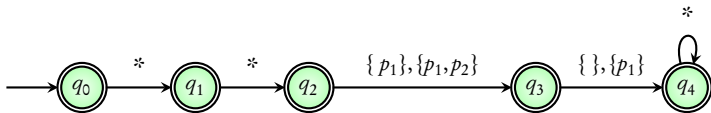


$$\mathbf{XX}(p_1 \wedge \mathbf{X}\neg p_2)$$

$$\begin{aligned} & \{ \} \{ \} \{ p_1 \} \{ \} \{ p_2 \} \{ p_1, p_2 \} \{ p_2 \} \{ p_2 \} \{ p_2 \} \dots \\ & \{ p_2 \} \{ p_1 \} \{ p_1, p_2 \} \{ p_1 \} \{ \} \{ p_1 \} \{ p_1 \} \{ p_1, p_2 \} \dots \\ & \quad \vdots \end{aligned}$$

$$\mathbf{XX}(p_1 \wedge \mathbf{X}\neg p_2)$$

$\{\}\{\}\{p_1\}\{\}\{p_2\}\{p_1,p_2\}\{p_2\}\{p_2\}\{p_2\}\dots$   
 $\{p_2\}\{p_1\}\{p_1,p_2\}\{p_1\}\{\}\{p_1\}\{p_1\}\{p_1,p_2\}\dots$   
 $\vdots$



$$p_1 \cup p_2$$

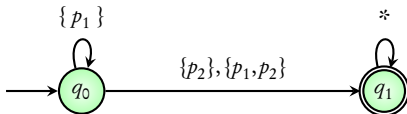
$$\{p_1\} \{p_1\} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$$

$$\{p_1, p_2\} \{ \} \{ \} \{ \} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$$

⋮

$p_1 \cup p_2$

$\{p_1\} \{p_1\} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$   
 $\{p_1, p_2\} \{ \} \{ \} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$   
 $\vdots$



$(X \ p_1) \ U \ p_2$

$\{p_2\} \{\} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$

$\{\} \{p_1\} \{p_1\} \{p_1\} \{p_1, p_2\} \{p_1, p_2\} \dots$

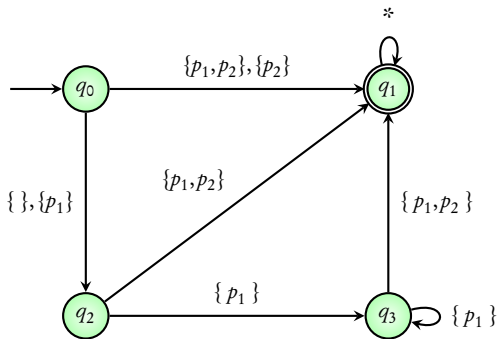
$\{\} \{p_1, p_2\} \{\} \{\} \{p_2\} \{p_1, p_2\} \dots$

$\vdots$



# $(X p_1) U p_2$

$\{p_2\} \{ \} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$   
 $\{ \} \{p_1\} \{p_1\} \{p_1\} \{p_1, p_2\} \{p_1, p_2\} \dots$   
 $\{ \} \{p_1, p_2\} \{ \} \{ \} \{ \} \{p_2\} \{p_1, p_2\} \dots$   
 $\vdots$



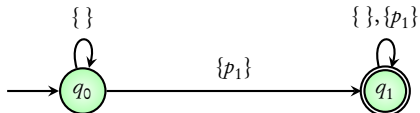
$$\mathbf{F} p_1 \wedge \neg \mathbf{F} p_2$$

$$\{ \} \{ \} \{ \} \{ p_1 \} \{ p_1 \} \{ \} \{ \} \{ p_1 \} \{ p_1 \} \dots$$
$$\{ p_1 \} \{ \} \{ \} \{ \} \{ \} \{ \} \{ \} \{ \} \{ \} \{ \} \dots$$

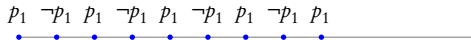
⋮

$$\mathbf{F} p_1 \wedge \neg \mathbf{F} p_2$$

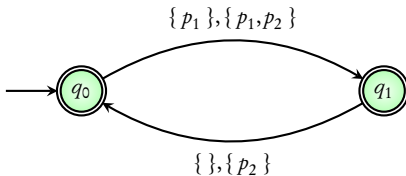
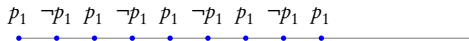
$\{\} \{\} \{\} \{p_1\} \{p_1\} \{\} \{\} \{p_1\} \{p_1\} \dots$   
 $\{p_1\} \{\} \{\} \{\} \{\} \{\} \{\} \{\} \{\} \dots$   
 $\vdots$



$$p_1 \wedge \mathbf{X} \neg p_1 \wedge \mathbf{G} ( p_1 \leftrightarrow \mathbf{X} \mathbf{X} p_1 )$$



$$p_1 \wedge \mathbf{X} \neg p_1 \wedge \mathbf{G} ( p_1 \leftrightarrow \mathbf{X} \mathbf{X} p_1 )$$



**GF**  $p_1$  Words where  $p_1$  occurs infinitely often

$\{ \} \{ p_1 \} \{ p_2 \} \{ p_1, p_2 \} \{ p_2 \} \{ p_1 \} \{ p_2 \} \dots$

$\{ \} \{ \} \{ \} \{ \} \{ p_1 \} \{ p_1 \} \{ p_1 \} \dots$

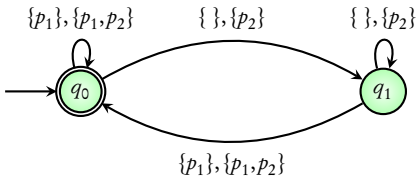
$\vdots$

**GF**  $p_1$  Words where  $p_1$  occurs infinitely often

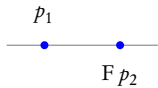
$\{ \} \{ p_1 \} \{ p_2 \} \{ p_1, p_2 \} \{ p_2 \} \{ p_1 \} \{ p_2 \} \dots$

$\{ \} \{ \} \{ \} \{ \} \{ p_1 \} \{ p_1 \} \{ p_1 \} \dots$

$\vdots$

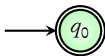
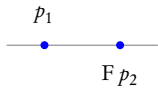


$G(p_1 \rightarrow XF p_2)$

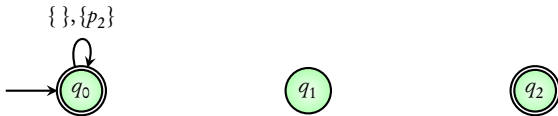
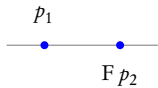




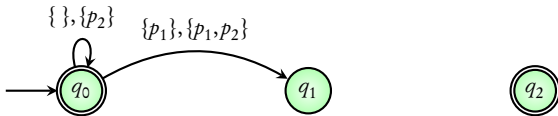
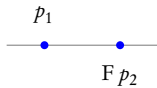
$G(p_1 \rightarrow \text{XF } p_2)$



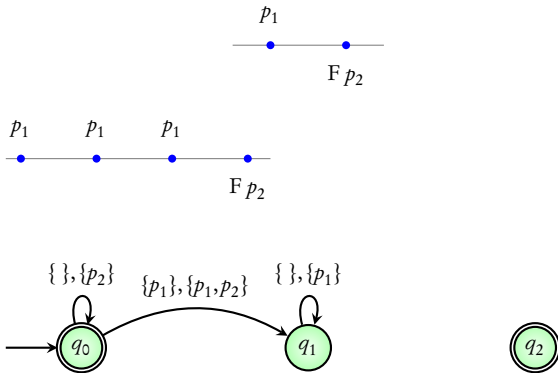
$G(p_1 \rightarrow \text{XF } p_2)$



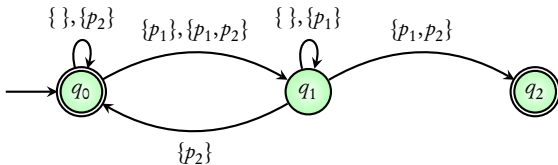
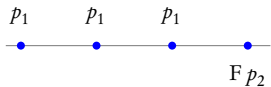
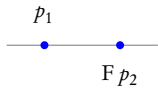
$G(p_1 \rightarrow \text{XF } p_2)$



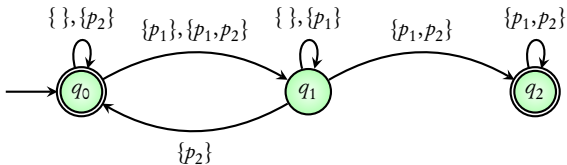
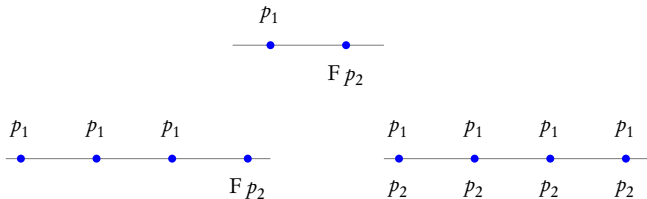
$G(p_1 \rightarrow \text{XF } p_2)$



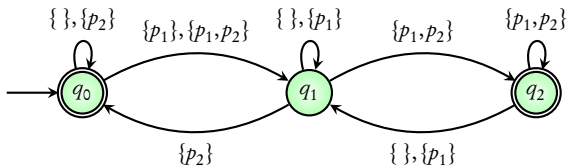
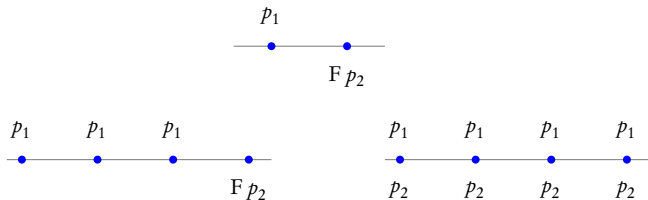
$G(p_1 \rightarrow \text{XF } p_2)$



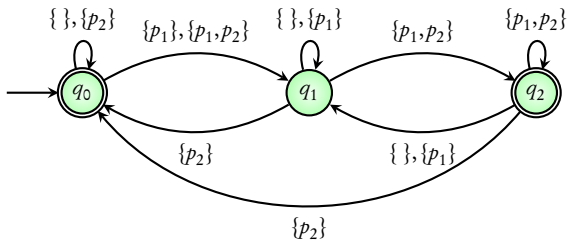
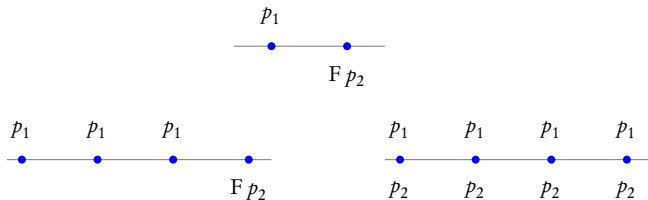
$G(p_1 \rightarrow \text{XF } p_2)$



$G(p_1 \rightarrow \text{XF } p_2)$



$G(p_1 \rightarrow \text{XF } p_2)$





# Summary

## LTL model-checking

Method

LTL to NBA examples

Module 2:  
**LTL to NBA**

**Goal:** Understand the **evaluation** of an LTL formula on an infinite word

$$p_1 \cup p_2$$

$$p_1 \cup p_2$$

$$\{p_1\} \quad \{p_1\} \quad \{p_1\} \quad \{p_1\} \quad \{p_2\} \quad \{p_1\} \quad \{p_1\} \quad \{p_1\} \quad \{p_1, p_2\} \quad \dots$$

$$p_1 \cup p_2$$

$\{p_1\}$   $\{p_1\}$   $\{p_1\}$   $\{p_1\}$   $\{p_2\}$   $\{p_1\}$   $\{p_1\}$   $\{p_1\}$   $\{p_1, p_2\}$   $\dots$

$p_1$

$p_2$

$p_1 \cup p_2$

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\}$ ...
$p_1$									
$p_2$									
$p_1 \cup p_2$									

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\} \dots$
$p_1$	1	1	1	1	0	1	1	1	1
$p_2$	0	0	0	0	1	0	0	0	1
$p_1 \cup p_2$									



$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\} \dots$
$p_1$	1	1	1	1	0	1	1	1	1
$p_2$	0	0	0	0	1	0	0	0	1
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\} \dots$
$p_1$	1	1	1	1	0	1	1	1	1
$p_2$	0	0	0	0	1	0	0	0	1
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1

$p_1$									
$p_2$									
$p_1 \cup p_2$									

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\} \dots$
$p_1$	1	1	1	1	0	1	1	1	1
$p_2$	0	0	0	0	1	0	0	0	1
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1

	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\} \dots$
$p_1$									
$p_2$									
$p_1 \cup p_2$									

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\} \dots$
$p_1$	1	1	1	1	0	1	1	1	1
$p_2$	0	0	0	0	1	0	0	0	1
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1

	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\} \dots$
$p_1$	1	0	1	0	1	0	1	0	1
$p_2$	0	0	0	0	0	0	0	0	0
$p_1 \cup p_2$									

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\} \dots$
$p_1$	1	1	1	1	0	1	1	1	1
$p_2$	0	0	0	0	1	0	0	0	1
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1

	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\} \dots$
$p_1$	1	0	1	0	1	0	1	0	1
$p_2$	0	0	0	0	0	0	0	0	0
$p_1 \cup p_2$	0	0	0	0	0	0	0	0	0

**GF**  $p_1$

**GF**  $p_1$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\mathbf{GF} p_1$$

recall that  $\mathbf{F} \phi = true \mathbf{U} \phi$  and  $\mathbf{G} \phi = \neg true \mathbf{U} \neg \phi$

$$\neg true \mathbf{U} \neg (true \mathbf{U} p_1)$$



$$\mathbf{GF} p_1$$

recall that  $\mathbf{F} \phi = true \mathbf{U} \phi$  and  $\mathbf{G} \phi = \neg true \mathbf{U} \neg \phi$

$$\neg true \mathbf{U} \neg (true \mathbf{U} p_1)$$

$\{\}$   $\{\}$   $\{p_1\}$   $\{\}$   $\{\}$   $\{p_1\}$   $\{\}$   $\{\}$   $\{p_1\}$

$$GF p_1$$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }
<i>p<sub>1</sub></i>									
<i>true</i>									
<i>true</i> $\cup$ <i>p<sub>1</sub></i>									
$\neg true \cup p_1$									
<i>true</i> $\cup$ $\neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$\mathbf{GF} p_1$$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }
<i>p<sub>1</sub></i>	0	0	1	0	0	1	0	0	1
<i>true</i>									
<i>true</i> $\cup$ <i>p<sub>1</sub></i>									
$\neg true \cup p_1$									
<i>true</i> $\cup$ $\neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$\mathbf{GF} p_1$$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }
<i>p<sub>1</sub></i>	0	0	1	0	0	1	0	0	1
<i>true</i>	1	1	1	1	1	1	1	1	1
<i>true</i> $\cup$ <i>p<sub>1</sub></i>									
$\neg true \cup p_1$									
<i>true</i> $\cup$ $\neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }
<i>p<sub>1</sub></i>	0	0	1	0	0	1	0	0	1
<i>true</i>	1	1	1	1	1	1	1	1	1
<i>true</i> $\cup$ <i>p<sub>1</sub></i>	1	1	1	1	1	1	1	1	1
$\neg true \cup p_1$									
<i>true</i> $\cup$ $\neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }
<i>p<sub>1</sub></i>	0	0	1	0	0	1	0	0	1
<i>true</i>	1	1	1	1	1	1	1	1	1
<i>true</i> $\cup$ <i>p<sub>1</sub></i>	1	1	1	1	1	1	1	1	1
$\neg true \cup p_1$	0	0	0	0	0	0	0	0	0
<i>true</i> $\cup$ $\neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }
<i>p<sub>1</sub></i>	0	0	1	0	0	1	0	0	1
<i>true</i>	1	1	1	1	1	1	1	1	1
<i>true</i> $\cup$ <i>p<sub>1</sub></i>	1	1	1	1	1	1	1	1	1
$\neg true \cup p_1$	0	0	0	0	0	0	0	0	0
<i>true</i> $\cup$ $\neg(true \cup p_1)$	0	0	0	0	0	0	0	0	0
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }
<i>p<sub>1</sub></i>	0	0	1	0	0	1	0	0	1
<i>true</i>	1	1	1	1	1	1	1	1	1
<i>true</i> $\cup$ <i>p<sub>1</sub></i>	1	1	1	1	1	1	1	1	1
$\neg true \cup p_1$	0	0	0	0	0	0	0	0	0
<i>true</i> $\cup$ $\neg(true \cup p_1)$	0	0	0	0	0	0	0	0	0
$\neg true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1	1



**GF**  $p_1$

**GF**  $p_1$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\mathbf{GF} p_1$$

recall that  $\mathbf{F} \phi = \mathit{true} \mathbf{U} \phi$  and  $\mathbf{G} \phi = \neg \mathit{true} \mathbf{U} \neg \phi$

$$\neg \mathit{true} \mathbf{U} \neg(\mathit{true} \mathbf{U} p_1)$$

$$\mathbf{GF} p_1$$

recall that  $\mathbf{F} \phi = true \mathbf{U} \phi$  and  $\mathbf{G} \phi = \neg true \mathbf{U} \neg \phi$

$$\neg true \mathbf{U} \neg (true \mathbf{U} p_1)$$

{p<sub>1</sub>} {p<sub>1</sub>} {} {} {} {} {} {} {}

$$GF p_1$$

recall that  $F \phi = true \cup \phi$  and  $G \phi = \neg true \cup \neg \phi$

$$\neg true \cup \neg(true \cup p_1)$$

	{p <sub>1</sub> }	{p <sub>1</sub> }	{}	{}	{}	{}	{}	{}	{}
<i>p<sub>1</sub></i>									
<i>true</i>									
<i>true</i> ∪ <i>p<sub>1</sub></i>									
$\neg true \cup p_1$									
<i>true</i> ∪ $\neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
$p_1$	1	1	0	0	0	0	0	0	0
$true$									
$true \cup p_1$									
$\neg true \cup p_1$									
$true \cup \neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
$p_1$	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$									
$\neg true \cup p_1$									
$true \cup \neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
$p_1$	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$									
$true \cup \neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									



$$GF p_1$$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
$p_1$	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$	0	0	1	1	1	1	1	1	1
$true \cup \neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
$p_1$	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$	0	0	1	1	1	1	1	1	1
$true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1	1
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that  $F\phi = true \cup \phi$  and  $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
$p_1$	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$	0	0	1	1	1	1	1	1	1
$true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1	1
$\neg true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0	0

$$\mathbf{F} (\neg p_1 \wedge \mathbf{X} (\neg p_2 \mathbf{U} p_1))$$

$$\mathbf{F} (\neg p_1 \wedge \mathbf{X} (\neg p_2 \mathbf{U} p_1))$$

recall that  $\mathbf{F} \phi = \text{true} \mathbf{U} \phi$

$$\mathbf{F} (\neg p_1 \wedge \mathbf{X} (\neg p_2 \mathbf{U} p_1))$$

recall that  $\mathbf{F} \phi = \mathit{true} \mathbf{U} \phi$

$$\mathit{true} \mathbf{U} (\neg p_1 \wedge \mathbf{X} (\neg p_2 \mathbf{U} p_1))$$

$$\mathbf{F} (\neg p_1 \wedge \mathbf{X} (\neg p_2 \mathbf{U} p_1))$$

recall that  $\mathbf{F} \phi = \mathit{true} \mathbf{U} \phi$

$$\mathit{true} \mathbf{U} (\neg p_1 \wedge \mathbf{X} (\neg p_2 \mathbf{U} p_1))$$

$$\{\} \quad \{p_2\} \quad \{\} \quad \{\} \quad \{p_1\} \quad \{p_1, p_2\} \quad \{p_1, p_2\} \quad \dots$$

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$								
$p_2$								
$\neg p_1$								
$\neg p_2$								
$\neg p_2 U p_1$								
$X (\neg p_2 U p_1)$								
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								



$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	{}	{p <sub>2</sub> }	{}	{}	{p <sub>1</sub> }	{p <sub>1</sub> ,p <sub>2</sub> }	{p <sub>1</sub> ,p <sub>2</sub> }	...
<i>p</i> <sub>1</sub>	0	0	0	0	1	1	1	
<i>p</i> <sub>2</sub>	0	1	0	0	0	1	1	
$\neg p_1$								
$\neg p_2$								
$\neg p_2 U p_1$								
$X (\neg p_2 U p_1)$								
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$								
$X (\neg p_2 U p_1)$								
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	{}	{p <sub>2</sub> }	{}	{}	{p <sub>1</sub> }	{p <sub>1</sub> ,p <sub>2</sub> }	{p <sub>1</sub> ,p <sub>2</sub> }	...
<i>p</i> <sub>1</sub>	0	0	0	0	1	1	1	
<i>p</i> <sub>2</sub>	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$					1	1	1	
$X (\neg p_2 U p_1)$								
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	{}	{p <sub>2</sub> }	{}	{}	{p <sub>1</sub> }	{p <sub>1</sub> , p <sub>2</sub> }	{p <sub>1</sub> , p <sub>2</sub> }	...
<i>p</i> <sub>1</sub>	0	0	0	0	1	1	1	
<i>p</i> <sub>2</sub>	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$			1	1	1	1	1	
$X (\neg p_2 U p_1)$								
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$								
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	{}	{p <sub>2</sub> }	{}	{}	{p <sub>1</sub> }	{p <sub>1</sub> ,p <sub>2</sub> }	{p <sub>1</sub> ,p <sub>2</sub> }	...
<i>p<sub>1</sub></i>	0	0	0	0	1	1	1	
<i>p<sub>2</sub></i>	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0							
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	{}	{p <sub>2</sub> }	{}	{}	{p <sub>1</sub> }	{p <sub>1</sub> , p <sub>2</sub> }	{p <sub>1</sub> , p <sub>2</sub> }	...
<i>p<sub>1</sub></i>	0	0	0	0	1	1	1	
<i>p<sub>2</sub></i>	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1						
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1					
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								



$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1				
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1			
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1		
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0							
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1						
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1	1					
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1	1	1				
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								



$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1	1	1	0			
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1	1	1	0	0		
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1	1	1	0	0	0	
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1	1	1	0	0	0	
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$	1	1	1	1				

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that  $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1	1	1	0	0	0	
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$	1	1	1	1	0	0	0	

$$p_1 \cup p_2$$

	{p <sub>1</sub> }	{p <sub>1</sub> }	{p <sub>1</sub> }	{p <sub>1</sub> }	{p <sub>2</sub> }	{p <sub>1</sub> }	{p <sub>1</sub> }	{p <sub>1</sub> }	{p <sub>1</sub> ,p <sub>2</sub> }	...
<i>p<sub>1</sub></i>	1	1	1	1	0	1	1	1	1	
<i>p<sub>2</sub></i>	0	0	0	0	1	0	0	0	0	1
<i>p<sub>1</sub> ∪ p<sub>2</sub></i>	1	1	1	1	1	1	1	1	1	

	{p <sub>1</sub> }	{}	{p <sub>1</sub> }	{}	{p <sub>1</sub> }	{}	{p <sub>1</sub> }	{}	{p <sub>1</sub> }	...
<i>p<sub>1</sub></i>	1	0	1	0	1	0	1	0	1	
<i>p<sub>2</sub></i>	0	0	0	0	0	0	0	0	0	
<i>p<sub>1</sub> ∪ p<sub>2</sub></i>	0	0	0	0	0	0	0	0	0	

$$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$$

	{}	{p <sub>2</sub> }	{}	{}	{p <sub>1</sub> }	{p <sub>1</sub> ,p <sub>2</sub> }	{p <sub>1</sub> ,p <sub>2</sub> }	...
<i>p<sub>1</sub></i>	0	0	0	0	1	1	1	
<i>p<sub>2</sub></i>	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 \cup p_1$	0	0	1	1	1	1	1	
$X(\neg p_2 \cup p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X(\neg p_2 \cup p_1)$	0	1	1	1	0	0	0	
$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$	1	1	1	1	1	1	0	

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p <sub>1</sub> }	{}	{}	{p <sub>1</sub> }	{}	{}
<i>p<sub>1</sub></i>	0	0	1	0	0	1	0	0
<i>true</i>	1	1	1	1	1	1	1	1
<i>true ∪ p<sub>1</sub></i>	1	1	1	1	1	1	1	1
$\neg true \cup p_1$	0	0	0	0	0	0	0	0
$true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0
$\neg true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1

	{p <sub>1</sub> }	{p <sub>1</sub> }	{}	{}	{}	{}	{}	{}	{}
<i>p<sub>1</sub></i>	1	1	0	0	0	0	0	0	0
<i>true</i>	1	1	1	1	1	1	1	1	1
<i>true ∪ p<sub>1</sub></i>	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$	0	0	1	1	1	1	1	1	1
$true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1	1
$\neg true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0	0

# Formula expansions

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\}$	...
$p_1$	1	1	1	1	0	1	1	1	1	
$p_2$	0	0	0	0	1	0	0	0	1	
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1	

	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	...
$p_1$	1	0	1	0	1	0	1	0	1	
$p_2$	0	0	0	0	0	0	0	0	0	
$p_1 \cup p_2$	0	0	0	0	0	0	0	0	0	

$$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 \cup p_1$	0	0	1	1	1	1	1	
$X(\neg p_2 \cup p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X(\neg p_2 \cup p_1)$	0	1	1	1	0	0	0	
$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$	1	1	1	1	1	1	0	

$$\neg true \cup \neg(true \cup p_1)$$

	$\{\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{\}$
$p_1$	0	0	1	0	0	1	0	0
$true$	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	1	1	1	1	1	1
$\neg true \cup p_1$	0	0	0	0	0	0	0	0
$true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0
$\neg true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
$p_1$	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$	0	0	1	1	1	1	1	1	1
$true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1	1
$\neg true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0	0

**Key idea:** Construct automata whose **states are columns of the formula expansion**



**Key idea:** Construct automata whose **states are columns of the formula expansion**

**Next in this module:** understand **properties** of formula expansions

## Word compatibility

$p_1$								
$p_2$								

## Word compatibility

	{ }							
$p_1$	0							
$p_2$	0							

## Word compatibility

	{}		{ $p_1$ }					
$p_1$	0		1					
$p_2$	0		0					

## Word compatibility

	{ }		{ $p_1$ }		{ $p_2$ }			
$p_1$	0		1		0			
$p_2$	0		0		1			

## Word compatibility

	$\{\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1, p_2\}$
$p_1$	0	1	0	1
$p_2$	0	0	1	1

## AND-NOT-compatibility

$\phi$

	0		1	
--	---	--	---	--

$\neg\phi$

	1		0	
--	---	--	---	--

## AND-NOT-compatibility

 $\phi$ 

	0		1	
--	---	--	---	--

 $\neg\phi$ 

	1		0	
--	---	--	---	--

 $\phi_1$ 

1		0		1		0
---	--	---	--	---	--	---

 $\phi_2$ 

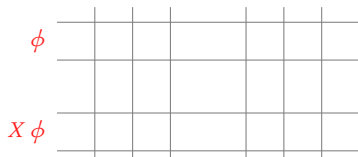
1		1		0		0
---	--	---	--	---	--	---

 $\phi_1 \wedge \phi_2$ 

1		0		0		0
---	--	---	--	---	--	---



## X-compatibility



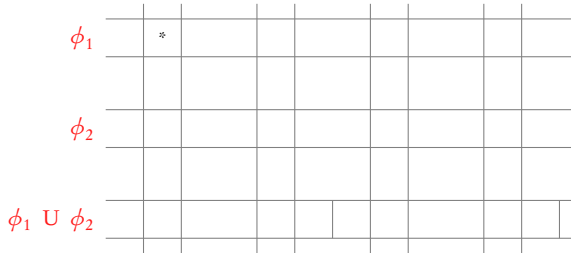
## X-compatibility

$\phi$		0				
$X\phi$	0					

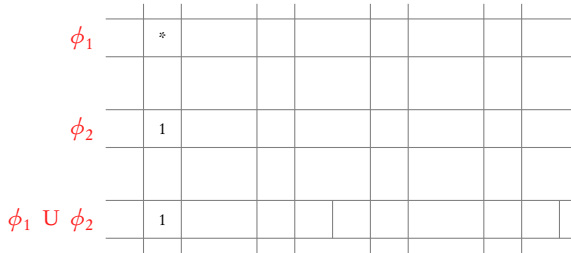
## X-compatibility

$\phi$		0			1	
$X\phi$	0				1	

# Until-compatibility



## Until-compatibility



# Until-compatibility

$\phi_1$	*						
$\phi_2$	1		0				
$\phi_1 \text{ U } \phi_2$	1		1				

## Until-compatibility

$\phi_1$	*	1					
$\phi_2$	1	0					
$\phi_1 \text{ U } \phi_2$	1	1	1				

## Until-compatibility

$\phi_1$	*	1					
$\phi_2$	1	0		0			
$\phi_1 \text{ U } \phi_2$	1	1	1		0		



## Until-compatibility

$\phi_1$	*		1		0			
$\phi_2$	1		0		0			
$\phi_1 \text{ U } \phi_2$	1		1	1	0			

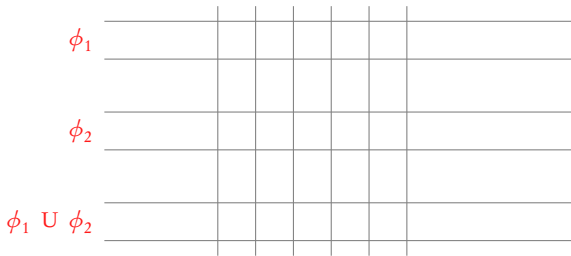
## Until-compatibility

$\phi_1$	*		1		0		1	
$\phi_2$	1		0		0		0	
$\phi_1 \text{ U } \phi_2$	1		1	1		0		0

## Until-compatibility

$\phi_1$	*		1		0		1	
$\phi_2$	1		0		0		0	
$\phi_1 \text{ U } \phi_2$	1		1	1		0		0

## Until-compatibility: eventuality condition



## Until-compatibility: eventuality condition

$\phi_1$	1					
$\phi_2$	0					
$\phi_1 \text{ U } \phi_2$	1	1				

## Until-compatibility: eventuality condition

$\phi_1$		1	1				
$\phi_2$		0	0				
$\phi_1 \text{ U } \phi_2$		1	1	1			

## Until-compatibility: eventuality condition

$\phi_1$		1	1	1			
$\phi_2$		0	0	0			
$\phi_1 \text{ U } \phi_2$		1	1	1	1		

## Until-compatibility: eventuality condition

$\phi_1$		1	1	1	1		
$\phi_2$		0	0	0	0		
$\phi_1 \text{ U } \phi_2$		1	1	1	1	1	



## Until-compatibility: eventuality condition

$\phi_1$		1	1	1	1	1	
$\phi_2$		0	0	0	0	0	...
$\phi_1 \text{ U } \phi_2$		1	1	1	1	1	

## Until-compatibility: eventuality condition

$\phi_1$		1	1	1	1	1	
$\phi_2$		0	0	0	0	0	...
$\phi_1 \cup \phi_2$		1	1	1	1	1	

Cannot happen forever that  $\phi_1 \cup \phi_2 = 1$ ,  $\phi_1 = 1$  but  $\phi_2 = 0$

# Accepting expansions

$$p_1 \cup p_2$$

	$p_1$	$p_2$	$p_1$	$p_2$	$p_1$	$p_2$	$p_1$	$p_2$	$p_1$	$p_2$	$p_1 \cup p_2$	...
$p_1$	1	1	1	1	0	1	1	1	1	1	1	
$p_2$	0	0	0	0	1	0	0	0	0	1		
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1	1		

	$p_1$	$\perp$	$p_2$	$\perp$	$p_1$	$\perp$	$p_2$	$\perp$	$p_1$	$\perp$	$p_1 \cup p_2$	...
$p_1$	1	0	1	0	1	0	1	0	1	0	1	
$p_2$	0	0	0	0	0	0	0	0	0	0	0	
$p_1 \cup p_2$	0	0	0	0	0	0	0	0	0	0	0	

$$\text{true} \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$$

	$\perp$	$p_2$	$\perp$	$\perp$	$p_1$	$p_1 \cup p_2$	$p_1 \cup p_2$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 \cup p_1$	0	0	1	1	1	1	1	
$X(\neg p_2 \cup p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X(\neg p_2 \cup p_1)$	0	1	1	1	0	0	0	
$\text{true} \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$	1	1	1	1	1	1	1	

$$\neg \text{true} \cup \neg(\text{true} \cup p_1)$$

	$\perp$	$\perp$	$p_1$	$\perp$	$\perp$	$p_1$	$\perp$	$\perp$
$p_1$	0	0	1	0	0	1	0	0
$\text{true}$	1	1	1	1	1	1	1	1
$\text{true} \cup p_1$	1	1	1	1	1	1	1	1
$\neg \text{true} \cup p_1$	0	0	0	0	0	0	0	0
$\text{true} \cup \neg(\text{true} \cup p_1)$	0	0	0	0	0	0	0	0
$\neg \text{true} \cup \neg(\text{true} \cup p_1)$	1	1	1	1	1	1	1	1

	$p_1$	$p_1$	$\perp$	$\perp$	$\perp$	$\perp$	$\perp$	$\perp$	$\perp$
$p_1$	1	1	0	0	0	0	0	0	0
$\text{true}$	1	1	1	1	1	1	1	1	1
$\text{true} \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg \text{true} \cup p_1$	0	0	1	1	1	1	1	1	1
$\text{true} \cup \neg(\text{true} \cup p_1)$	1	1	1	1	1	1	1	1	1
$\neg \text{true} \cup \neg(\text{true} \cup p_1)$	0	0	0	0	0	0	0	0	0

Entry in **first column** of **last row** (corresponding to final formula) is 1

# Summary

## LTL to NBA

Formula expansions

Properties

Columns as states of NBA

Module 3:  
**Automaton construction**

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\}$	...
$p_1$	1	1	1	1	0	1	1	1	1	
$p_2$	0	0	0	0	1	0	0	0	1	
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1	

	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	...
$p_1$	1	0	1	0	1	0	1	0	1	
$p_2$	0	0	0	0	0	0	0	0	0	
$p_1 \cup p_2$	0	0	0	0	0	0	0	0	0	

$$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 \cup p_1$	0	0	1	1	1	1	1	
$X(\neg p_2 \cup p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X(\neg p_2 \cup p_1)$	0	1	1	1	0	0	0	
$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$	1	1	1	1	1	1	0	

$$\neg true \cup \neg(true \cup p_1)$$

	$\{\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{\}$
$p_1$	0	0	1	0	0	1	0	0
$true$	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	1	1	1	1	1	1
$\neg true \cup p_1$	0	0	0	0	0	0	0	0
$true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0
$\neg true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
$p_1$	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$	0	0	1	1	1	1	1	1	1
$true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1	1
$\neg true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0	0

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\}$	...
$p_1$	1	1	1	1	0	1	1	1	1	
$p_2$	0	0	0	0	1	0	0	0	1	
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1	

	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	...
$p_1$	1	0	1	0	1	0	1	0	1	
$p_2$	0	0	0	0	0	0	0	0	0	
$p_1 \cup p_2$	0	0	0	0	0	0	0	0	0	

$$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	...
$p_1$	0	0	0	0	1	1	1	
$p_2$	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 \cup p_1$	0	0	1	1	1	1	1	
$X(\neg p_2 \cup p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X(\neg p_2 \cup p_1)$	0	1	1	1	0	0	0	
$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$	1	1	1	1	1	1	0	

$$\neg true \cup \neg(true \cup p_1)$$

	$\{\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{\}$
$p_1$	0	0	1	0	0	1	0	0
$true$	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	1	1	1	1	1	1
$\neg true \cup p_1$	0	0	0	0	0	0	0	0
$true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0
$\neg true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
$p_1$	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$	0	0	1	1	1	1	1	1	1
$true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1	1
$\neg true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0	0

Construct an automaton with states as column vectors that can guess accepting expansions

**Example 1:**  $p_1 \cup p_2$



$p_1$	0	0	0	0	1	1	1	1
$p_2$	0	0	1	1	0	0	1	1
$p_1 \cup p_2$	0	1	0	1	0	1	0	1

$p_1$	0	0	0	0	1	1	1	1
$p_2$	0	0	1	1	0	0	1	1
$p_1 \cup p_2$	0	1	0	1	0	1	0	1

$\phi_1$	*		1		0		1	
$\phi_2$	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0

**Recall Until-compatibility**

$p_1$	0	0	0	1	1	1	1
$p_2$	0	1	1	0	0	1	1
$p_1 \cup p_2$	0	0	1	0	1	0	1

$\phi_1$	*	1	0	1	
$\phi_2$	1	0	0	0	
$\phi_1 \cup \phi_2$	1	1	1	0	0

**Recall Until-compatibility**

$p_1$	0		0	1	1	1	1	1
$p_2$	0		1	0	0	1	1	1
$p_1 \cup p_2$	0		1	0	1	0	1	1

$\phi_1$	*		1		0		1		
$\phi_2$	1		0		0		0		
$\phi_1 \cup \phi_2$	1		1	1		0		0	0

## Recall Until-compatibility

$p_1$	0	0	1	1	1
$p_2$	0	1	0	0	1
$p_1 \cup p_2$	0	1	0	1	1

$\phi_1$	*		1		0		1	
$\phi_2$	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1		0		0 0

## Recall Until-compatibility

$p_1$	0	0	1	1	1
$p_2$	0	1	0	0	1
$p_1 \cup p_2$	0	1	0	1	1

## Compatible states

$\phi_1$	*	1	0	1		
$\phi_2$	1	0	0	0		
$\phi_1 \cup \phi_2$	1	1	1	0	0	0

## Recall Until-compatibility

$p_1$	0	0	1	1	1
$p_2$	0	1	0	0	1
$p_1 \cup p_2$	0	1	0	1	1

$q_0$

$p_1$	0
$p_2$	0
$p_1 \cup p_2$	0

0
1
1

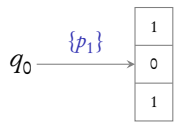
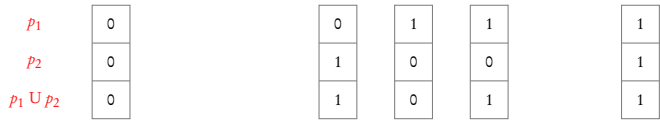
1
0
0

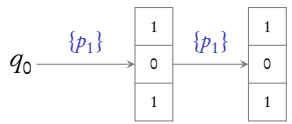
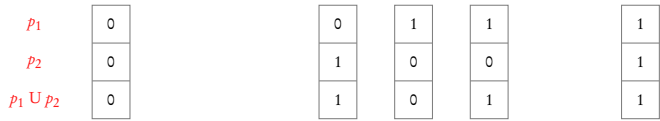
1
0
1

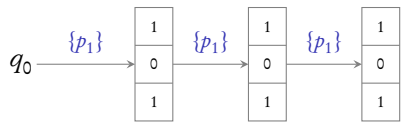
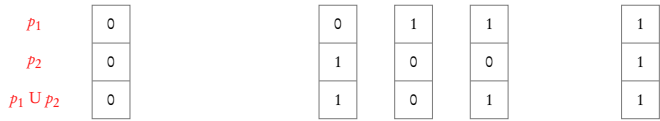
1
1
1

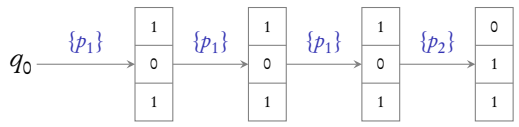
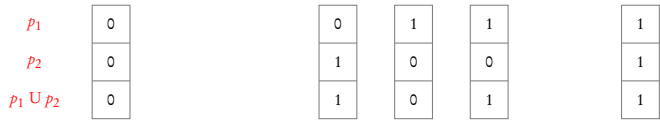
$q_0 \xrightarrow{\{p_1\}}$

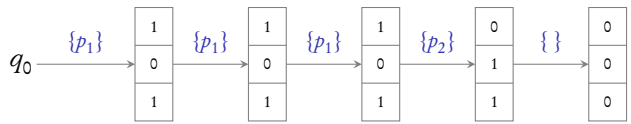
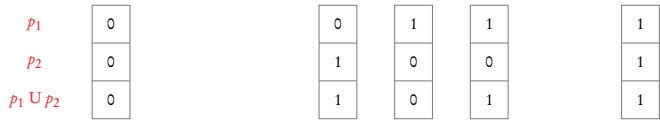


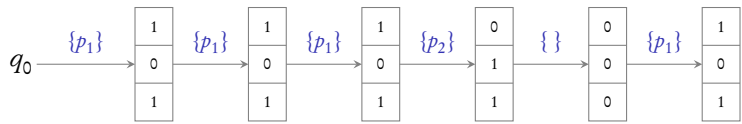
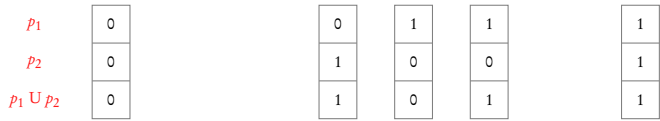


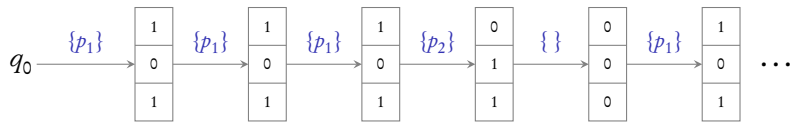
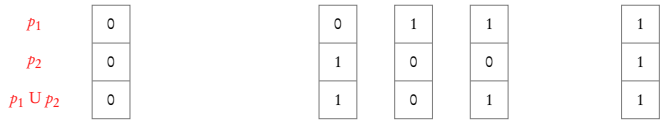


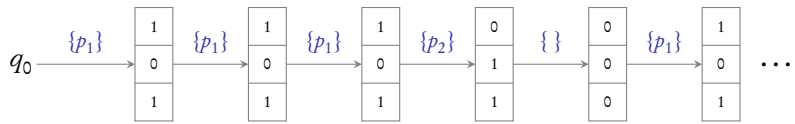
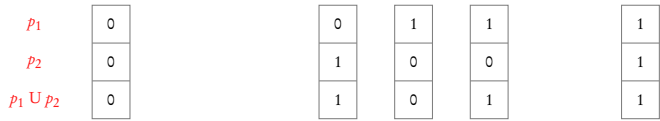






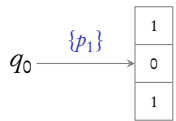
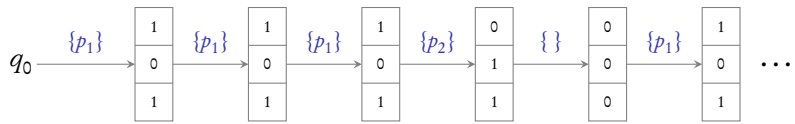
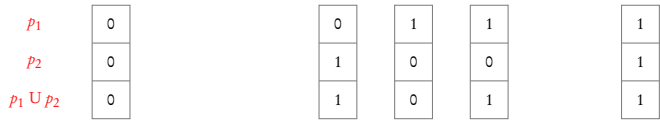


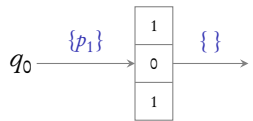
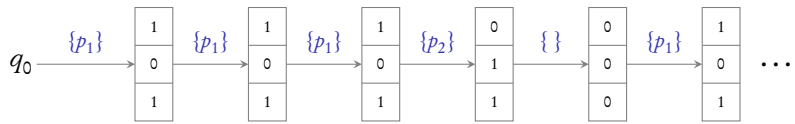
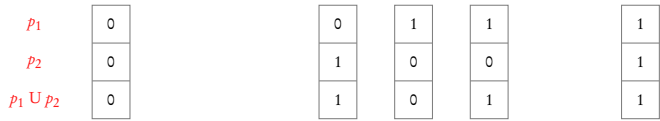


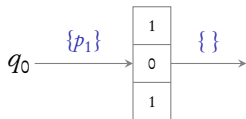
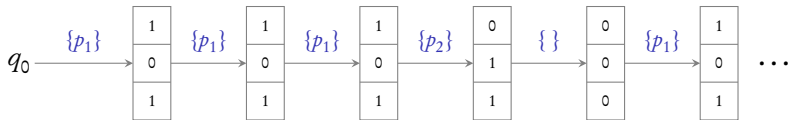
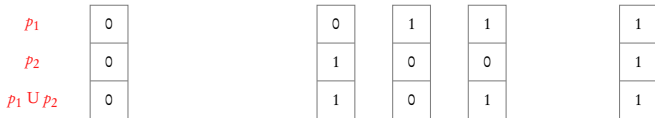


$q_0$









**No compatible transition**

$\phi_1$	*		1		0		1	
$\phi_2$	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0

1
0
1

$\longrightarrow q_0$

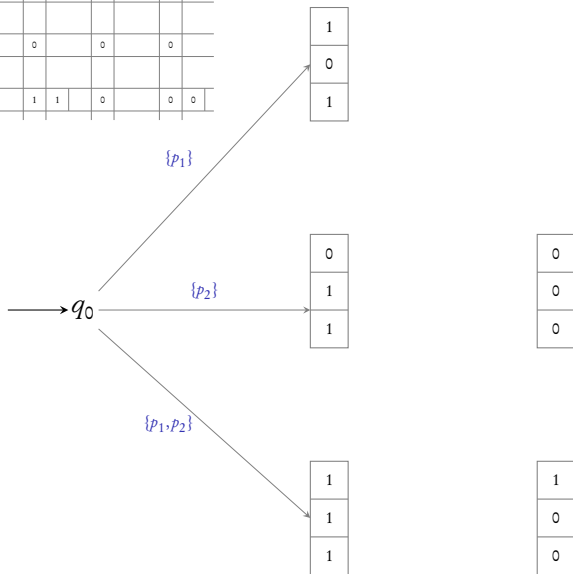
0
1
1

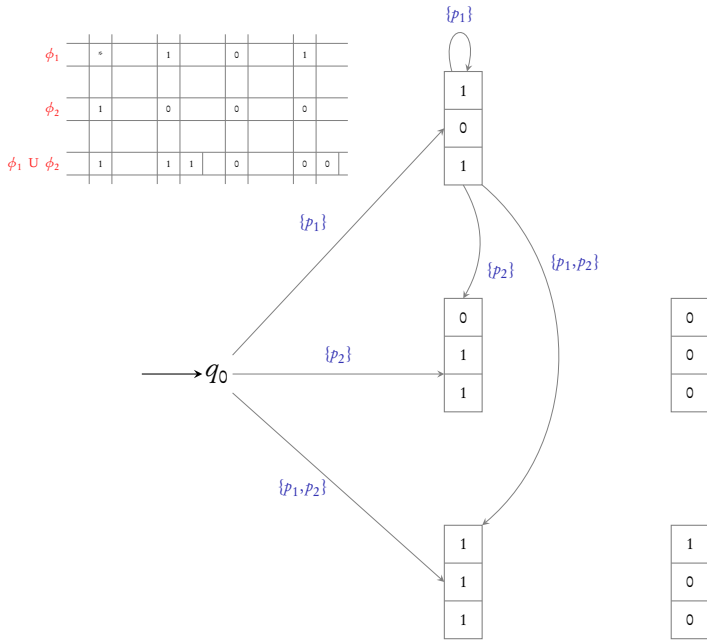
0
0
0

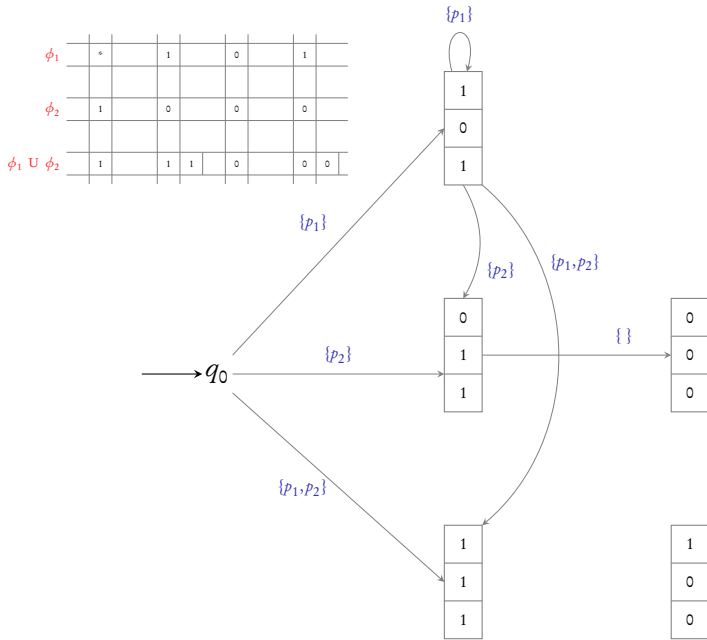
1
1
1

1
0
0

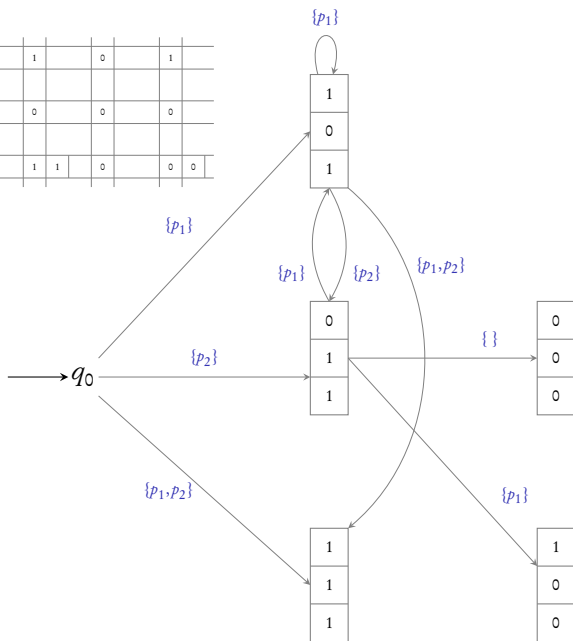
$\phi_1$	*		1		0		1	
$\phi_2$	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1		0		0 0





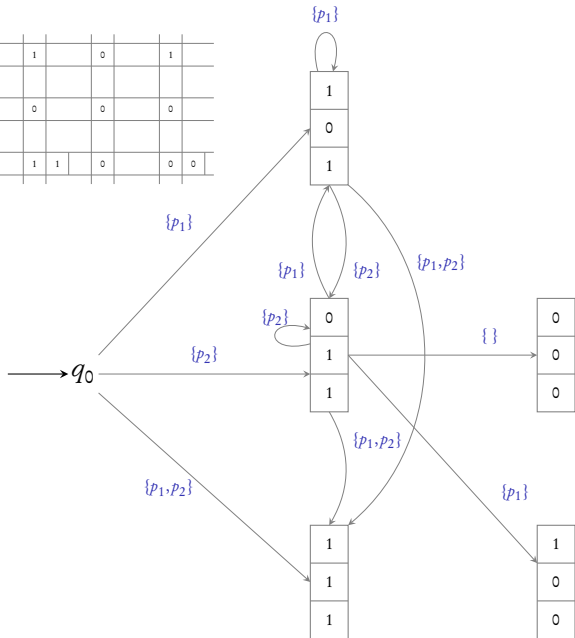


$\phi_1$	*		1		0		1	
$\phi_2$	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0

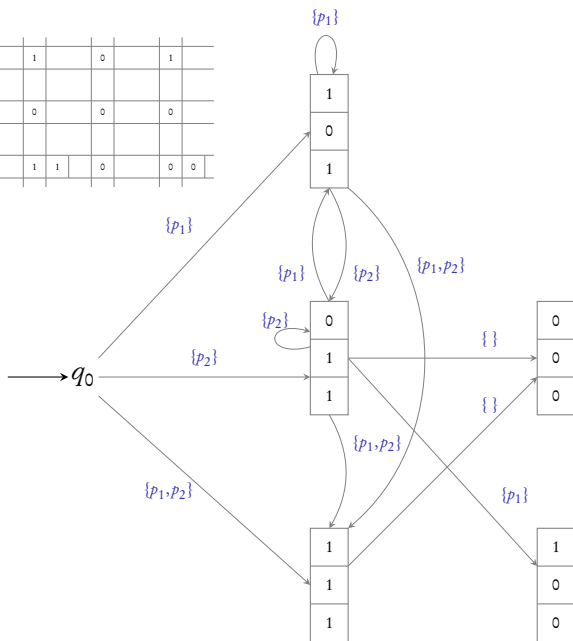




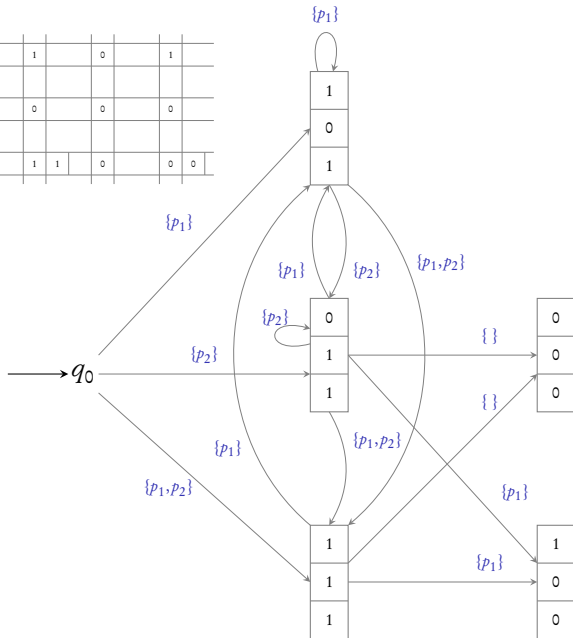
$\phi_1$	*		1		0		1	
$\phi_2$	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0



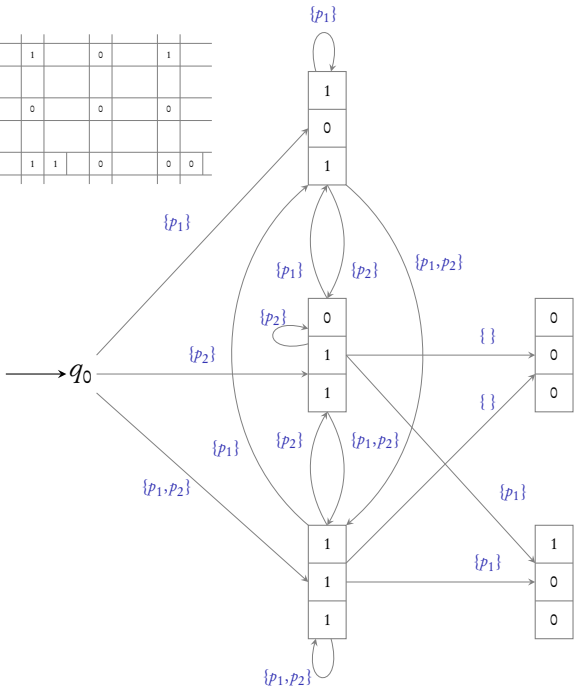
$\phi_1$	*		1		0		1	
$\phi_2$	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0



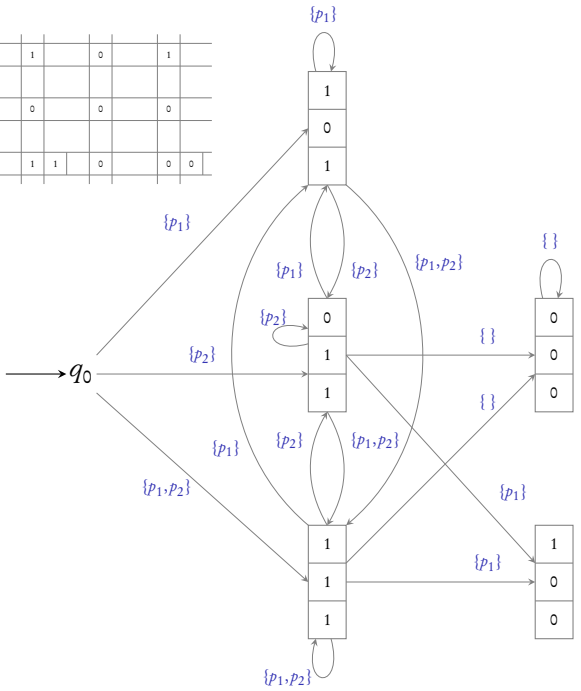
$\phi_1$	*		1		0		1	
$\phi_2$	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0



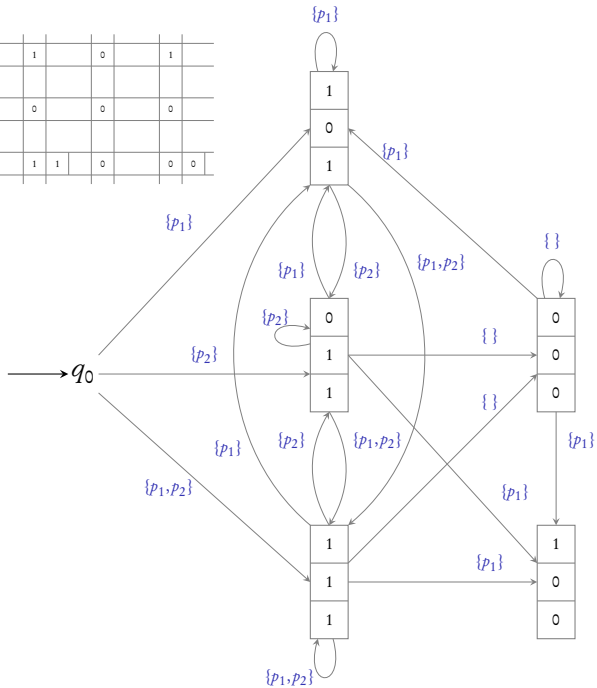
$\phi_1$	*		1		0		1	
$\phi_2$	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0



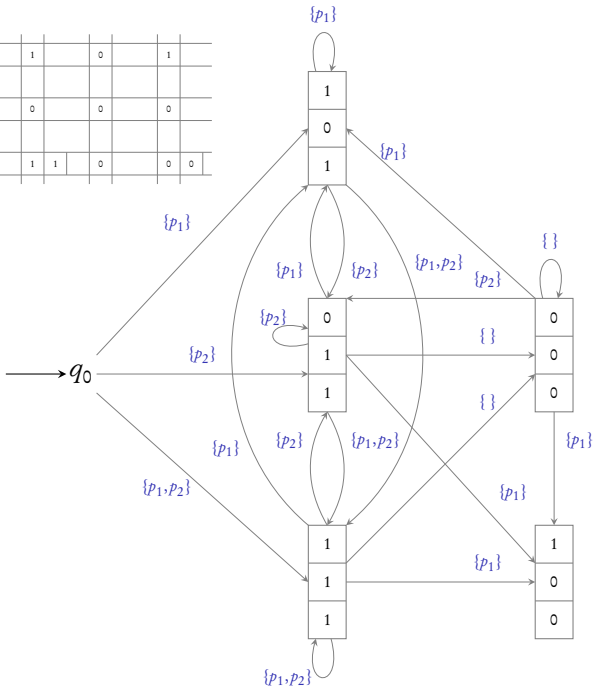
$\phi_1$	*		1		0		1	
$\phi_2$	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0



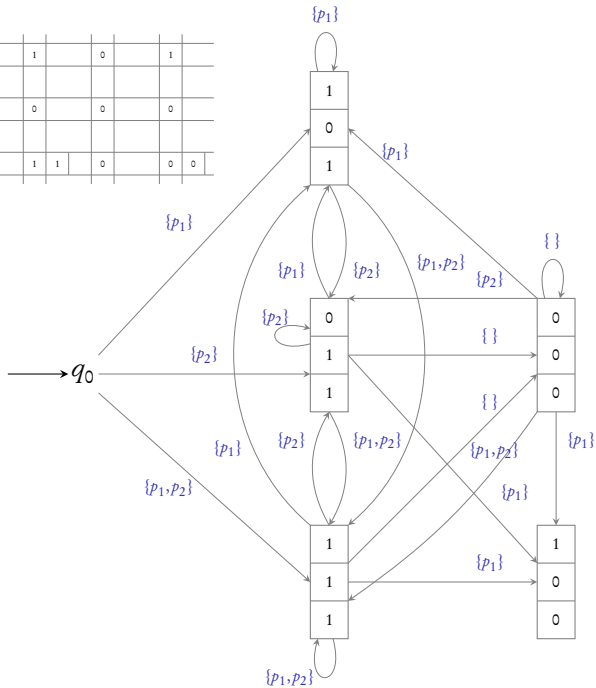
$\phi_1$	*		1		0		1	
$\phi_2$		1		0		0		0
$\phi_1 \cup \phi_2$		1		1	1		0	0



$\phi_1$	*		1		0		1	
$\phi_2$		1		0		0		0
$\phi_1 \cup \phi_2$		1		1	1		0	0

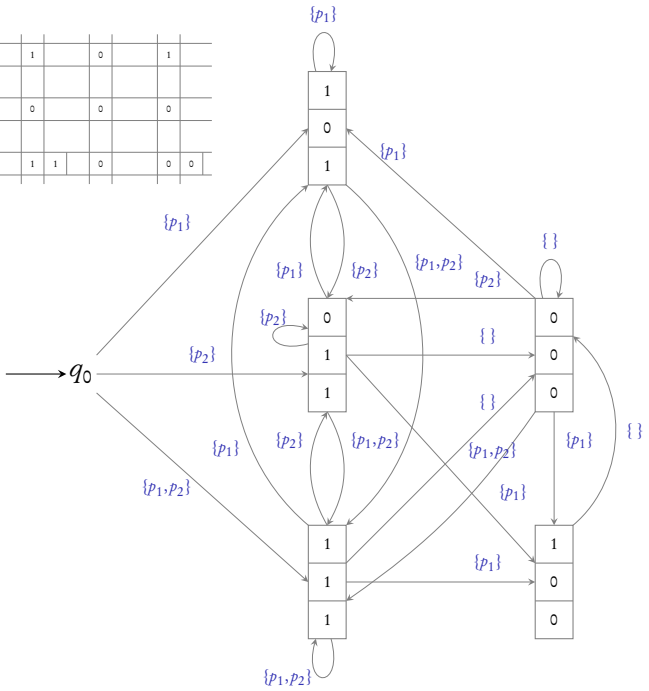


$\phi_1$	*		1		0		1	
$\phi_2$		1		0		0		0
$\phi_1 \cup \phi_2$		1		1	1		0	0

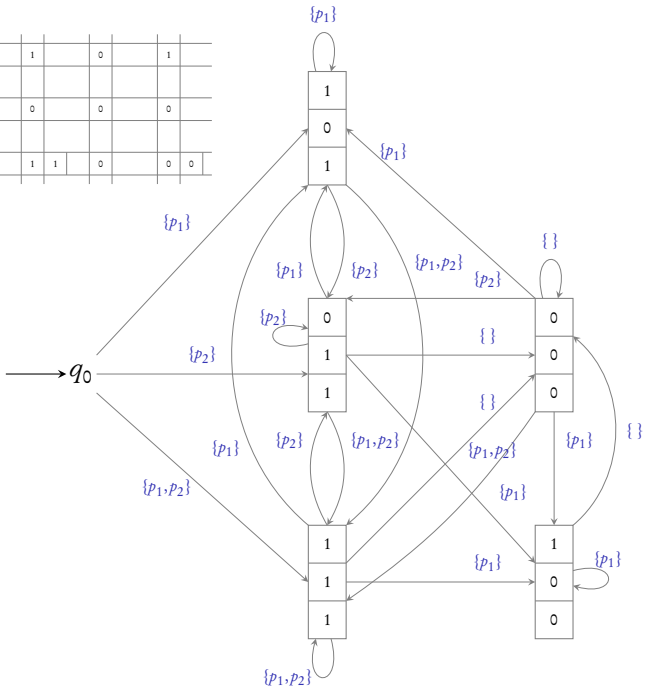




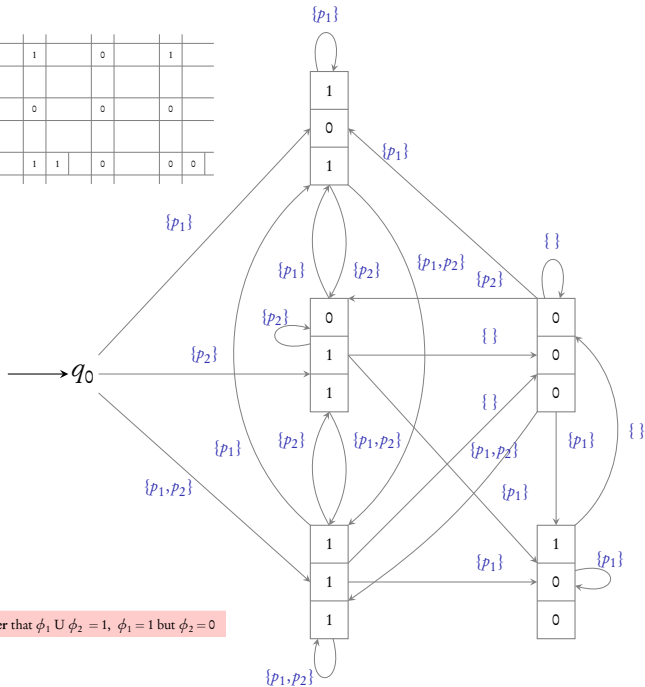
$\phi_1$	*		1		0		1	
$\phi_2$		1		0		0		0
$\phi_1 \cup \phi_2$		1		1	1		0	0



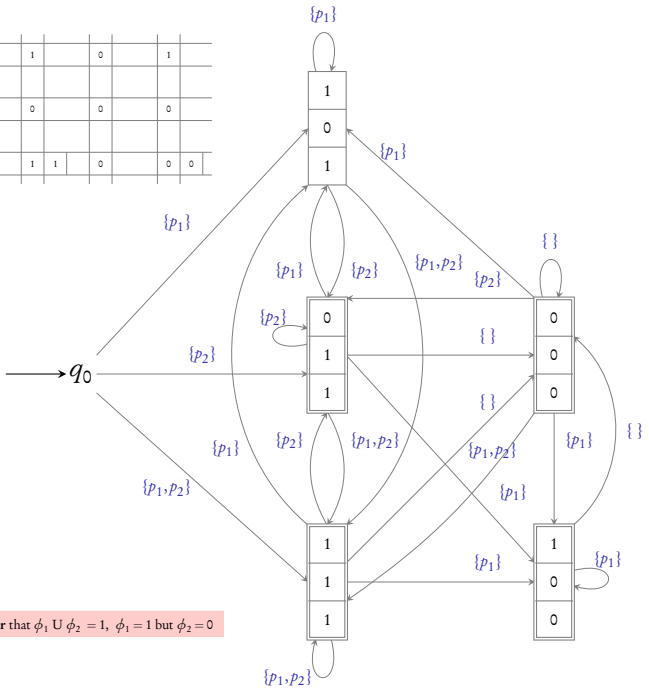
$\phi_1$	*		1		0		1	
$\phi_2$	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0



$\phi_1$	*		1		0		1	
$\phi_2$	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0



$\phi_1$	*		1		0		1	
$\phi_2$		1		0		0		0
$\phi_1 \cup \phi_2$		1		1	1		0	0



Cannot happen forever that  $\phi_1 \cup \phi_2 = 1$ ,  $\phi_1 = 1$  but  $\phi_2 = 0$

**Example 2:**  $(X p_1) \cup p_2$

$p_1$	*
$p_2$	*
$X p_1$	*
$(X p_1) \cup p_2$	*

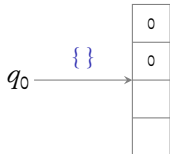
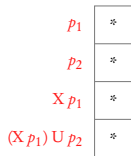
$p_1$	*
$p_2$	*
$X p_1$	*
$(X p_1) \cup p_2$	*

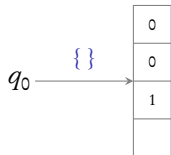
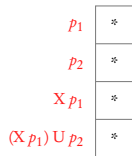
$q_0$

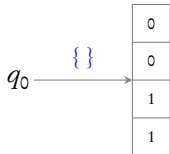
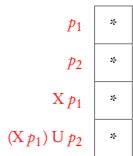
$p_1$	*
$p_2$	*
$X p_1$	*
$(X p_1) \cup p_2$	*

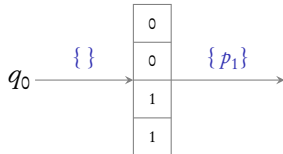
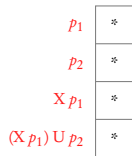
$q_0 \longrightarrow \{\}$

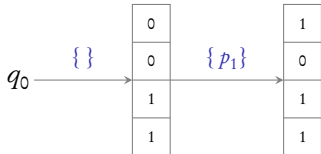
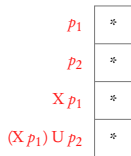


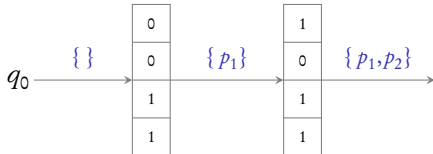
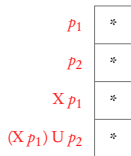


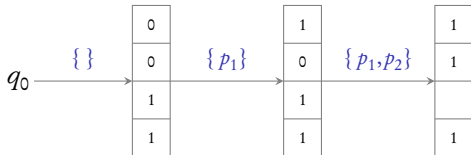
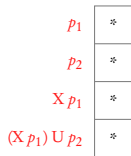


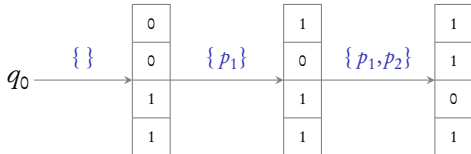
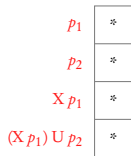




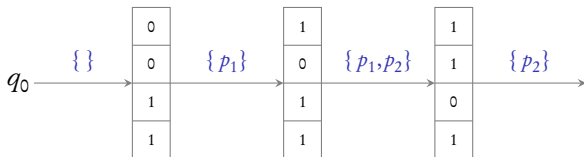
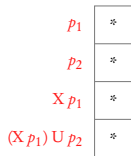


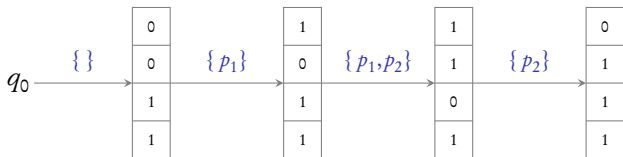
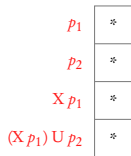


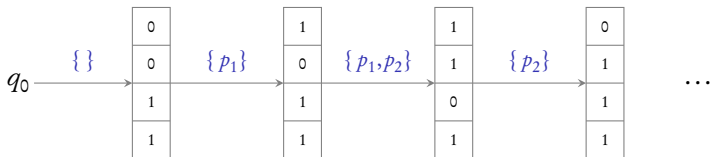
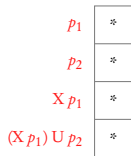












**Coming next:** Construction for an arbitrary LTL formula  $\phi$

**Step 1:** List down subformulae of  $\phi$

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$p_1$	*
$p_2$	*
$p_1 \cup p_2$	*

$p_1$	*
$p_2$	*
$\neg p_1$	*
$(\neg p_1) \cup p_2$	*

**Step 2:** Check **AND-NOT** and **Until** compatibility

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Incompatible states!



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Incompatible states!

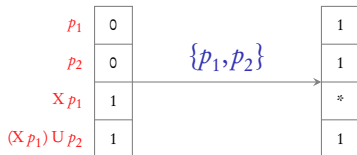
**Remove** incompatible states and **add** a new state  $\{q_0\}$

**Step 3:** Add transitions satisfying

**Word, X and Until** compatibility

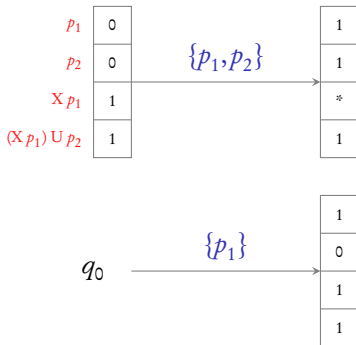
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From  $q_0$  add compatible transitions to states where last entry is 1

**Step 4:** Accepting states should ensure Until-eventuality condition

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For **every** Until subformula  $\phi_1 \text{ U } \phi_2$ , define

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0	1	0	1
0	0	1	1
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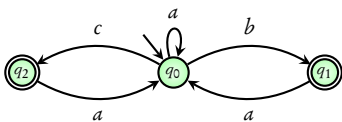
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In general, this algorithm gives NBA which is **exponential** in size of formula

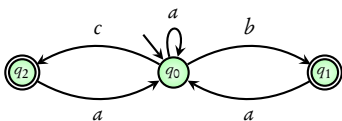
Module 4:  
**Generalized Büchi Automata**





$$(a^*(b+c)a)^\omega$$

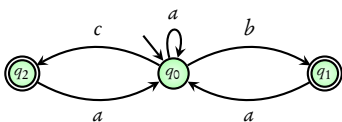
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Above NBA also accepts *ababababababab.....*

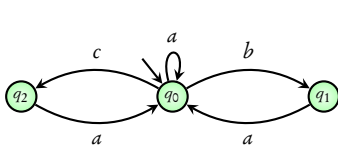


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Suppose we want NBA for **subset** of  $(a^*(b+c)a)^\omega$  where  
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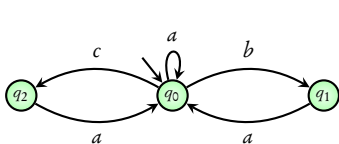
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# Generalized NBA



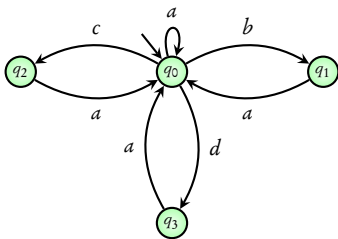
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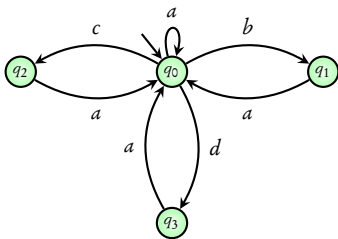
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Get GNBA for subset of  $(a^*(b + c + d)a)^\omega$  where:

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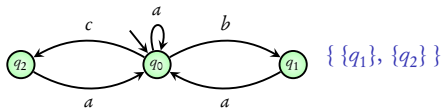
Accepting condition:  $\{ \{q_3\}, \{q_1, q_2\} \}$

# Generalized Büchi Automata

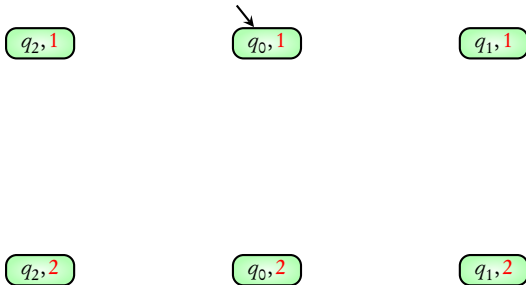
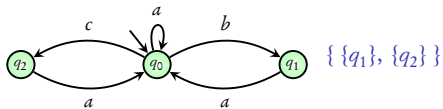
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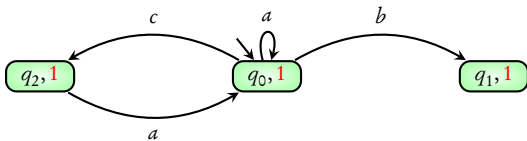
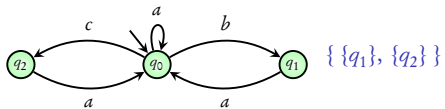
GNBA



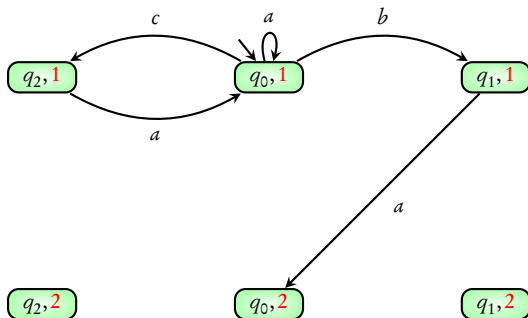
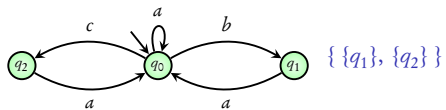
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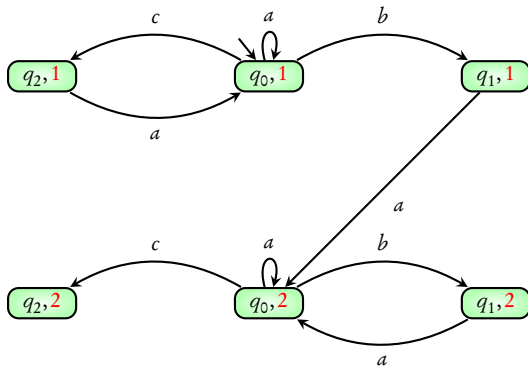
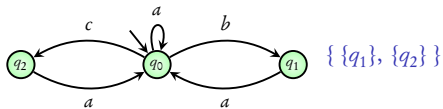
GNBA



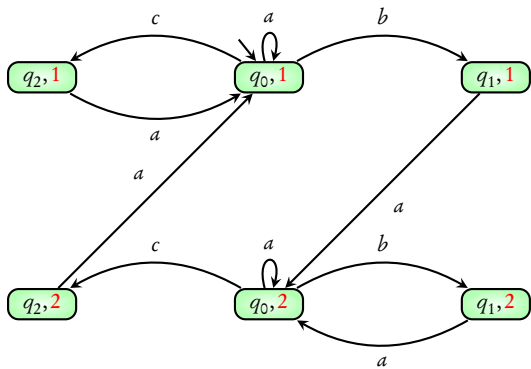
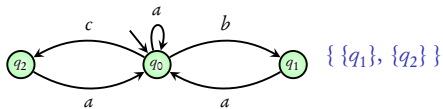
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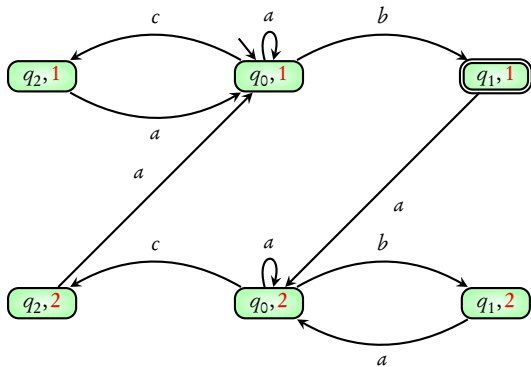
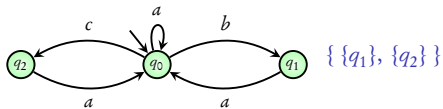
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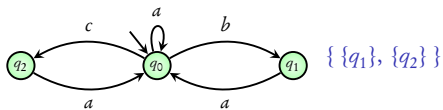
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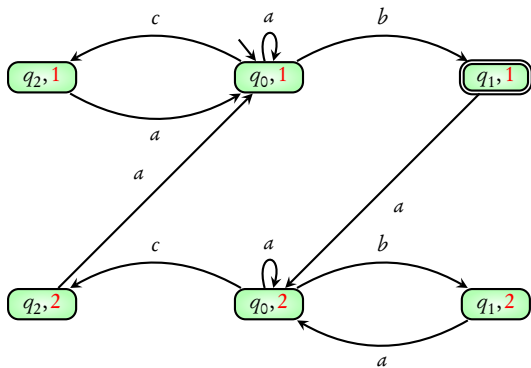
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GNBA

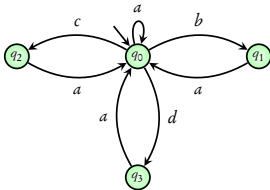


NBA



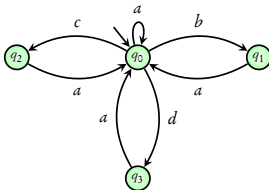


GNBA

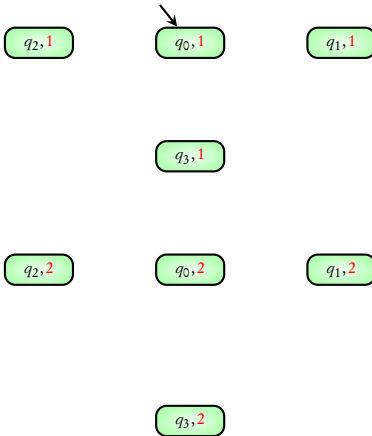


$\{q_3\}, \{q_1, q_2\}$

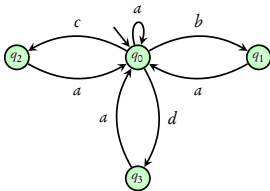
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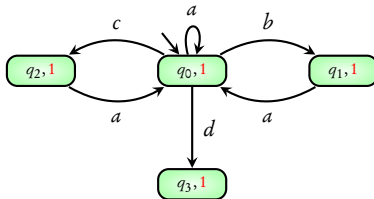
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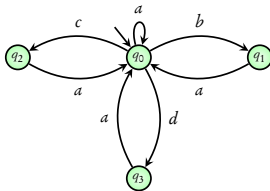
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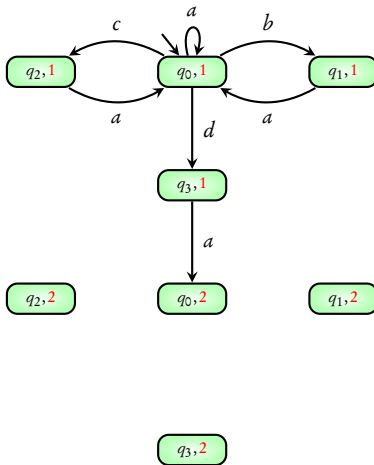
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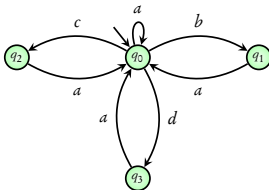
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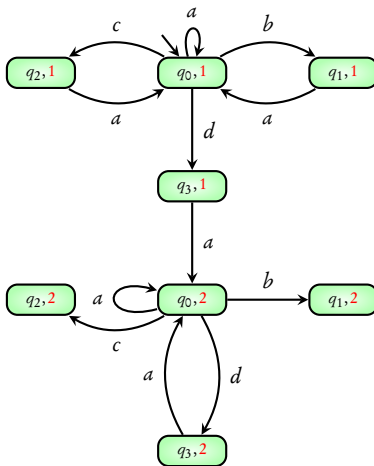
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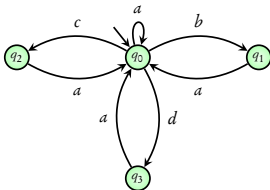
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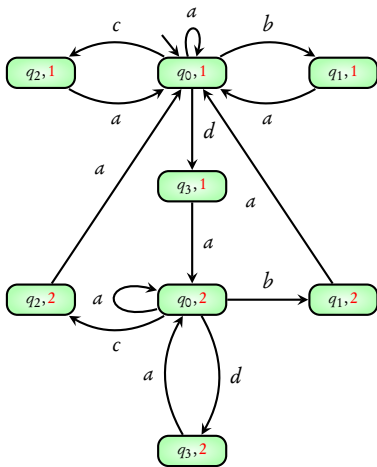
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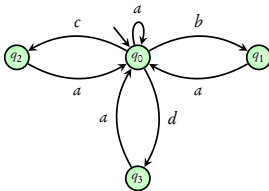
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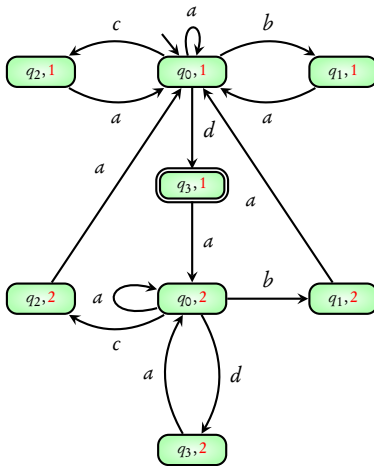
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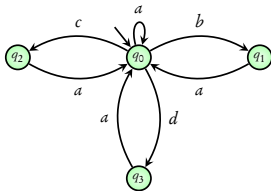
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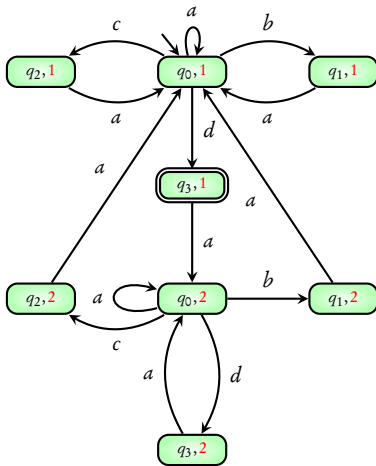


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